

### Solutions to Homework 1

**Q.1 (i)** Let  $u_1, \dots, u_n$  be any integers, positive or negative, satisfying

$$a_1u_1 + \dots + a_nu_n = 1.$$

We know from the lectures that such integers exist. Now let

$$m := \min\{a_1, \dots, a_n\}, \tag{1}$$

$$M := \max\{|u_1|, \dots, |u_n|\}, \tag{2}$$

$$s := \sum_{i=1}^n a_i.$$

Let  $r_0 := smM$ . I claim that  $G(a_1, \dots, a_n) < r_0$ . Clearly, there is a non-negative solution to  $\sum a_i x_i = r_0$ , namely  $x_i = mM$  for all  $i$ . Now let  $\delta \in \{0, 1, \dots, m-1\}$ . Then a solution to  $\sum a_i x_i = r_0 + \delta$  is given by

$$x_i = mM + \delta u_i, \quad i = 1, \dots, n. \tag{3}$$

By (1) and (2), we see that each  $x_i$  in (3) is non-negative. Thus we have shown that there is a non-negative solution to  $\sum a_i x_i = r$ , for each of a sequence of  $m$  integers starting at  $r_0 = smM$ . But from this it follows easily that there is a non-negative solution for any  $r \geq r_0$ . For suppose  $a_j = m$ . Then given a non-negative solution to  $\sum a_i x_i = r$ , for some  $r$ , a non-negative solution to  $\sum a_i x'_i = r + m$  is given by

$$x'_i = x_i, \quad \text{if } i \neq j, \quad x'_j = x_j + 1.$$

**(ii)** If  $a_1 = a_2 = 1$  then the theorem states that  $G(1, 1) = -1$ , in other words that every non-negative integer can be written as  $x + y$ , for some  $x, y \geq 0$ . This is clear. So we may henceforth assume, without loss of generality, that  $a_1 > a_2 \geq 1$ . We shall show two things :

**CLAIM 1 :** There is no non-negative integer solution to

$$a_1x + a_2y = (a_1 - 1)(a_2 - 1) - 1 = a_1a_2 - a_1 - a_2. \tag{4}$$

**CLAIM 2 :** There is a non-negative solution to

$$a_1x + a_2y = (a_1 - 1)(a_2 - 1) + \delta,$$

for every  $\delta \in \{0, 1, \dots, a_2 - 1\}$ . The fact that these claims in turn suffice to establish that  $G(a_1, a_2) = (a_1 - 1)(a_2 - 1) - 1$  is then argued in a manner similar to that given in part **(i)** above.

**PROOF OF CLAIM 1 :** Assume the contrary, and let  $(x, y)$  be a non-negative solution. First look at (4) modulo  $a_1$ . It reads

$$a_2(y + 1) \equiv 0 \pmod{a_1}.$$

Since  $\text{GCD}(a_1, a_2) = 1$  it follows from FTA that  $a_1|y + 1$ . Since  $y$  is non-negative, this means that  $y \geq a_1 - 1$ . Similarly, by looking at (4) modulo  $a_2$  we find that  $x \geq a_2 - 1$ . But then

$$a_1x + a_2y \geq a_1(a_2 - 1) + a_2(a_1 - 1) = 2a_1a_2 - a_1 - a_2 > a_1a_2 - a_1 - a_2,$$

a contradiction.

**PROOF OF CLAIM 2 :** The equation we want to solve can be rewritten as

$$a_1(x + 1) + a_2(y + 1) = a_1a_2 + \delta', \quad (5)$$

where  $\delta' \in \{1, 2, \dots, a_2\}$ . Since  $\text{GCD}(a_1, a_2) = 1$ , the integers  $a_1, 2a_1, \dots, a_2a_1$  represent, in some order, all congruence classes modulo  $a_2$  exactly once. In particular, there is some  $\xi \in \{1, 2, \dots, a_2\}$  such that

$$a_1\xi \equiv \delta' \pmod{a_2}.$$

In other words, there exists an integer  $\eta \geq 0$  such that

$$a_1\xi = \delta + a_2\eta \Rightarrow a_1\xi - a_2\eta = \delta.$$

Note that, since  $\xi \leq a_2$ , we also have  $\eta < a_1$ . Then a non-negative solution to (5) is given by  $x + 1 = \xi$ ,  $y + 1 = \eta + a_1$ .

**(iii)** Note that  $\text{GCD}(7, 19) = 1$ . Let

$$w := 7x + 19y. \quad (6)$$

Thus, in terms of  $w$  and  $z$ , the equation reads

$$w + 23z = 11. \quad (7)$$

One solution to (7) is clearly  $w = 34$ ,  $z = -1$ . Since  $\text{GCD}(1, 23) = 1$ , by Euclid the general solution to (7) is thus

$$w = 34 + 23n, \quad z = -1 - n, \quad n \in \mathbb{Z}.$$

Euclid also tells us that the general solution to (6) is

$$x = x_0w + 19m, \quad y = y_0w - 7m, \quad m \in \mathbb{Z},$$

where  $7x_0 + 19y_0 = 1$ . Clearly we can take, for example,  $x_0 = 11$ ,  $y_0 = -4$ . Substituting in everything, we find that the general solution to the original Diophantine equation is

$$\begin{aligned} x &= 11(34 + 23n) + 19m = 19m + 253n + 374, \\ y &= -4(34 + 23n) - 7m = -7m - 92n - 136, \\ z &= -1 - n, \quad m, n \in \mathbb{Z}. \end{aligned}$$

**(iv)** There seems to be a general agreement among those who did the homework that  $G(7, 19, 23) = 62$ . Note that  $(3, 1, 1)$  is a solution to the equation

$$7x + 19y + 23z = 63.$$

**(v)** There is a solution if and only if  $n \geq 12$  and  $n \notin \{13, 14, 15, 16, 18, 20, 21, 23, 25, 28, 30, 35\}$ . To see where the hell this problem came from, go to page 7 of the article [http://www.math.chalmers.se/~hegarty/HK\\_circle\\_oct26.pdf](http://www.math.chalmers.se/~hegarty/HK_circle_oct26.pdf)

**Q.2 (i)** Clearly, it suffices to prove the result for all primitive triples. Let  $(x, y, z)$  be any such triple. By Theorem 3.8 in the lecture notes we have, without loss of generality,

$$x = 2ab, \quad y = b^2 - a^2, \quad z = b^2 + a^2,$$

for some relatively prime integers  $a < b$  of opposite parity. Since one of  $a$  and  $b$  is thus even, it follows that  $4|x$ . If either  $a$  or  $b$  is divisible by 3, then so is  $x$ . Otherwise,  $b^2 \equiv a^2 \equiv 1 \pmod{3}$ , so  $3|y$ . In either case,  $3|xy$ . Since already  $4|x$ , FTA implies that  $12|xy$ . Finally, it suffices to show that at least one of  $x, y, z$  is divisible by 5. If either  $a$  or  $b$  is divisible by 5, then so is  $x$ . Otherwise, each of  $a^2$  and  $b^2$  must be congruent to either  $+1$  or  $-1 \pmod{5}$ . In particular,  $b^2 \equiv \pm a^2 \pmod{5}$ , hence  $5|(b^2 \mp a^2)$ , i.e.:  $5|y$  or  $5|z$ , as required.

**(ii)** Since  $t$  is odd, there exist positive integers  $i, j$  such that  $ti - 4j = 1$ . Now let  $m \in \mathbb{N}$  and take  $x = y = 2^{tm+j}$ ,  $z = 2^{4m+i}$ . Check that  $x^4 + y^4 = z^4$ . This gives an infinite family of solutions.

**(iii)** The exercise is easily reduced to proving that there are no 4-tuples  $(a, b, c, d) \in \mathbb{N}^4$  satisfying the pair of equations

$$a^2 + b^2 = c^2, \quad a^2 + d^2 = b^2. \quad (8)$$

We will use an infinite descent argument - assuming a solution exists to (8), we will construct a “smaller” one, in a sense which will become evident as the argument unfolds. This will yield a contradiction.

We can assume both Pythagorean triples are primitive, because if a prime  $p$  divides any two of  $a, b, c, d$ , then it is easily seen to divide all four, and hence  $(a/p, b/p, c/p, d/p)$  will be a smaller 4-tuple satisfying (8). Since both triples are primitive,  $b$  must be odd, hence  $a$  must be even. By Theorem 3.8 in the lecture notes, there exist integers positive integers  $x_1, y_1, x_2, y_2$ , with  $\text{GCD}(x_i, y_i) = 1$ ,  $x_i$  and  $y_i$  of opposite parity, such that

$$a = 2x_1y_1 = 2x_2y_2, \quad b = y_1^2 - x_1^2 = y_2^2 + x_2^2, \quad c = y_1^2 + x_1^2, \quad d = y_2^2 - x_2^2. \quad (9)$$

Note that the equations for  $b$  lead to

$$y_1^2 = x_1^2 + x_2^2 + y_2^2, \quad (10)$$

which forces  $y_1$  to be odd and  $x_1$  even, as vice versa would lead to the contradiction  $0 \equiv 2 \pmod{4}$ . Let

$$g_1 := \text{GCD}(x_1, x_2), \quad g_2 := \text{GCD}(y_1, y_2), \quad g_3 := \text{GCD}(x_1, y_2), \quad g_4 := \text{GCD}(x_2, y_1).$$

Since  $\text{GCD}(x_1, y_1) = \text{GCD}(x_2, y_2) = 1$ , it is easily seen that the  $g_i$  are pairwise relatively prime. Moreover, the equations for  $a$  in (9) imply that  $x_1 y_1 = x_2 y_2$ , from which it then follows that

$$x_1 = g_1 g_3, \quad x_2 = g_1 g_4, \quad y_1 = g_2 g_4, \quad y_2 = g_2 g_3. \quad (11)$$

Since  $y_1$  is odd and  $x_1$  even and the  $g_i$  are pairwise relatively prime, we deduce that  $g_2$  and  $g_4$  are both odd, and exactly one of  $g_1$  and  $g_3$  is even. If  $g_1$  is even and  $g_3$  odd, then substituting (11) into (10) leads to

$$g_2^2 \left( \frac{g_4^2 - g_3^2}{2} \right) = g_1^2 \left( \frac{g_4^2 + g_3^2}{2} \right).$$

The terms on the left and right are relatively prime in pairs, so it must be the case (by FTA) that

$$g_2^2 = \frac{g_4^2 + g_3^2}{2}, \quad g_1^2 = \frac{g_4^2 - g_3^2}{2},$$

which implies that the 4-tuple  $(g_1, g_2, g_4, g_3)$  also satisfies (8). But  $g_1 \leq g_1 g_2 g_3 g_4 = x_1 y_1 < 2x_1 y_1 = a$ , so this new 4-tuple is indeed “smaller”.

If instead  $g_3$  is even and  $g_1$  odd, a similar argument leads to the contradiction that  $(g_3, g_1, g_2, g_4)$  is a “smaller” 4-tuple. This completes the proof.

**(iv)** Let  $d = \text{GCD}(p^2 - q^2, p^2 + q^2)$ . Then  $d$  divides  $(p^2 - q^2) \pm (p^2 + q^2)$ , i.e.:  $d$  divides both  $2p^2$  and  $2q^2$ . Hence  $d$  divides  $\text{GCD}(2p^2, 2q^2) = 2$ , since  $\text{GCD}(p, q) = 1$ . Hence if  $p^2 - q^2$  were to divide  $p^2 + q^2$ , then  $p^2 - q^2$  would have to divide 2, in other words, since  $p > q$ , either  $p^2 - q^2 = 1$  or  $p^2 - q^2 = 2$ . It is easily checked that the latter equation has no integer solution, and the only solutions to the former are  $p = \pm 1, q = 0$ . But  $p$  and  $q$  are assumed positive, contradiction.

**(v)** Suppose by way of contradiction that we have a solution  $x^4 - y^4 = z^2$ . If  $d \mid x$  and  $d \mid y$ , then  $d^2 \mid z$  and so  $(x/d, y/d, z/d^2)$  is another solution. We may thus suppose  $x, y, z$  are pairwise relatively prime. The equation factorises as  $z^2 = (x^2 - y^2)(x^2 + y^2)$ . Since  $\text{GCD}(x, y) = 1$  it follows easily (see part **(iv)**) that  $\text{GCD}(x^2 - y^2, x^2 + y^2) \in \{1, 2\}$ . If the GCD is 1, then FTA implies that each of  $x^2 - y^2$  and  $x^2 + y^2$  is a square, i.e.: there exists integers  $a, b$  such that  $x^2 - y^2 = a^2$  and  $x^2 + y^2 = b^2$ . But then the

4-tuple  $(y, x, b, a)$  would satisfy (8), contradicting the fact that the latter system has no integer solutions. Similarly, if the GCD is 2, then we can write  $\left(\frac{z}{2}\right)^2 = \left(\frac{x^2-y^2}{2}\right)\left(\frac{x^2+y^2}{2}\right)$ , and now each factor on the right must be a square. So in this case, there exists integers  $a, b$  such that  $x^2 - y^2 = 2a^2$  and  $x^2 + y^2 = 2b^2$ . But then the 4-tuple  $(a, b, x, y)$  satisfies (8), again a contradiction.

**Q.3 (i)** Let  $p_1, p_2, \dots, p_k$  be the set of all primes appearing in the factorisations of either  $a$  or  $b$ . Let

$$a = \prod_{i=1}^k p_i^{\alpha_i}, \quad b = \prod_{i=1}^k p_i^{\beta_i}, \quad \alpha_i, \beta_i \geq 0.$$

We seek a solution of the form

$$x = \prod_{i=1}^k p_i^{\gamma_i}, \quad y = \prod_{i=1}^k p_i^{\delta_i}, \quad \gamma_i, \delta_i \geq 0.$$

One readily checks that  $ax^n = by^m$  is satisfied if and only if, for each  $i = 1, \dots, k$ ,

$$n\gamma_i - m\delta_i = \beta_i - \alpha_i. \quad (12)$$

But since  $\text{GCD}(m, n) = 1$ , (12) will always have a solution in non-negative integers.

**(ii)** Write the equation as  $x/y = \sqrt[n]{b/a}$ . First, compute  $d = \text{GCD}(a, b)$  using Euclid's algorithm. Then divide out  $d$  from both  $a$  and  $b$ , to get new numbers  $a_1, b_1$ . So far, this runs in polynomial time. Since  $\text{GCD}(b_1, a_1) = 1$ , the quotient  $b_1/a_1$  is the  $n$ :th power of a rational number if and only if each of  $a_1$  and  $b_1$  is the  $n$ :th power of an integer (this follows easily from FTA). Thus we just need to approximate  $\sqrt[n]{b_1}$  and  $\sqrt[n]{a_1}$ . Clearly, an approximation to the nearest integer can be obtained in polynomial time (by some kind of divide-and-conquer procedure, in other words, something which uses the Intermediate Value Theorem), and hence we can also determine in polynomial time if the  $n$ :th roots are in fact integers and, if so, which ones.

**Q.4** The first two properties are easy to verify, so I will concentrate on the third. To simplify notation, put  $\xi := \sqrt{-2}$ . Let  $a, b$  be two non-zero elements of  $R$ . First suppose  $b$  is a positive integer and let  $a = x + y\xi$ , where  $x, y \in \mathbb{Z}$ . Let  $r_1, r_2$  be the numerically least residues of  $x$  resp.  $y$  modulo  $b$ , i.e.: the unique numbers  $r_1, r_2 \in (-b/2, b/2]$  such that  $x = q_1b + r_1$  and  $y = q_2b + r_2$  for some  $q_1, q_2 \in \mathbb{Z}$ . Let  $q := q_1 + q_2\xi$  and  $r := r_1 + r_2\xi$ . Then

$a = qb + r$  and

$$d(r) = r_1^2 + 2r_2^2 \leq 3 \left( \frac{b}{2} \right)^2 < b^2 = d(b),$$

which proves property (iii) of a Euclidean domain in the case where  $b \in \mathbb{N}$ . Now consider a general  $b \in R$ . Set  $u := a\bar{b}$ ,  $v := b\bar{b} = d(b) \in \mathbb{N}$ . By the above, there exist  $q, r \in R$  such that

$$u = qv + r \quad \text{and} \quad 0 < d(r) < d(v).$$

Dividing across by  $\bar{b}$  we have that  $a = qb + r_0$ , where  $r_0 = r/\bar{b}$ . But the function  $d$  is clearly multiplicative, so  $d(v) = d(b)d(\bar{b})$  and  $d(r) = d(r_0)d(\bar{b})$  and hence  $0 < d(r) < d(v)$  implies that  $0 < d(r_0) < d(b)$ , so we are done.

**REMARK :** One may give a geometric description of the above argument, which is probably more enlightening than the ‘algebraic’ presentation given above.  $R = \mathbb{Z}[\sqrt{-2}]$  is a rectangular lattice in the complex plane generated by 1 and  $\sqrt{-2}$ . A rectangle with vertices at  $(0, 0)$ ,  $(1, 0)$ ,  $(1, 2)$  and  $(0, 2)$  forms a fundamental domain for this lattice (once one removes an appropriate part of the boundary). The set  $\{qb : q \in R\}$  is an ideal in the ring and corresponds geometrically to a sublattice generated by the orthogonal vectors  $b$  and  $\sqrt{-2}b$ . These two vectors span a rectangular fundamental domain  $\mathcal{F}$  for the sublattice. The point  $a \in \mathcal{C}$  must lie in some translate of  $\mathcal{F}$ , and hence there must be a point in the lattice whose distance to  $a$  is certainly no more than half the diameter of  $\mathcal{F}$ , i.e.: there exists  $q \in R$  such that

$$|a - qb| \leq \frac{1}{2} \text{diam}(\mathcal{F}) = \frac{1}{2} \sqrt{d(b)^2 + d(\sqrt{-2}b)^2} = \frac{\sqrt{3}}{2} d(b) < d(b), \quad \text{v.s.v.}$$

**Q.5 (i)** Since every  $a_n \geq 0$ , the limit, if it exists, must be non-negative. Subadditivity implies that  $a_n \leq na_1$ , hence  $a_n/n \leq a_1$  for any  $n$ . Hence the sequence  $(a_n/n)$  is bounded and must have a convergent subsequence. Suppose there is a subsequence  $(a_{n_i}/n_i)$  converging to a limit  $L$ , say. If  $L$  is not the limit of the entire sequence then we can find  $\epsilon > 0$  and another subsequence  $(a_{m_i}/m_i)$ , such that, for all  $i \geq i_0$ ,

$$\left| \frac{a_{n_i}}{n_i} - L \right| < \frac{\epsilon}{2}, \quad \left| \frac{a_{m_i}}{m_i} - L \right| > \epsilon. \quad (13)$$

Let’s suppose that, for all  $i \geq i_0$ ,  $a_{m_i}/m_i > a_{n_i}/n_i$  - a similar argument works in the case of the reverse inequality holding. Thus, by (13), for all  $i_1, i_2 \geq i_0$  we have that

$$\frac{a_{m_{i_1}}}{m_{i_1}} - \frac{a_{n_{i_2}}}{n_{i_2}} > \frac{\epsilon}{2}.$$

In particular, for all  $i \geq i_0$  one has that

$$\frac{a_{m_i}}{m_i} - \frac{a_{n_{i_0}}}{n_{i_0}} > \frac{\epsilon}{2}. \quad (14)$$

Let  $S := \max\{a_s : 1 \leq s < n_{i_0}\}$ . For each  $i$ , there exist integers  $q_i, r_i$  such that  $m_i = q_i n_{i_0} + r_i$ , where  $0 \leq r_i < n_{i_0}$ . By subadditivity, we have that  $a_{m_i} \leq q_i a_{n_{i_0}} + a_{r_i}$  and hence that

$$\frac{a_{m_i}}{m_i} \leq \frac{q_i a_{n_{i_0}} + a_{r_i}}{m_i} \leq \frac{a_{n_{i_0}}}{n_{i_0}} + \frac{S}{m_i}. \quad (15)$$

Now  $S$  is a fixed number and thus  $S/m_i \rightarrow 0$  as  $i \rightarrow \infty$ . Hence (15) will contradict (14) for all sufficiently large  $i$ . This completes the proof.

(REMARK : The result proven above is known in the literature as *Fakete's Lemma*).

(ii) Let  $m, n \in \mathbb{N}$ . If  $A$  is a subset of  $\{1, 2, \dots, m+n\}$  which is free of 3-term arithmetic progressions, then both  $A_1$  and  $A_2$  are also free of 3-APs, where  $A_1 := A \cap \{1, \dots, m\}$  and  $A_2 := A \cap \{m+1, \dots, m+n\}$ . This and the fact that a translate of an AP is still an AP easily lead to the conclusion that the function  $f(n)$  in this exercise is sub-additive. Then the result follows immediately from part (i).

(iii) First suppose  $a_0 \neq 0$ . Let  $p$  be any prime which does not divide  $a_0$  and let the set  $A$  consist of all multiples of  $p$ . Then  $A$  contains no solutions to the equation  $\mathcal{L}$ , as it doesn't even contain any solutions modulo  $p$ . Hence, in this case,

$$\liminf_{n \rightarrow \infty} \frac{f(n)}{n} \geq \liminf_{n \rightarrow \infty} \frac{|A \cap \{1, \dots, n\}|}{n} = \frac{1}{p} > 0.$$

Next, suppose  $\sum_{i=1}^n a_i \neq 0$ . Let  $s := \sum_{i=1}^n a_i$ , let  $k$  be the smallest positive integer such that  $ks \neq a_0$  and let  $p$  be any prime satisfying  $p > \max\{k|s|, |a_0|\}$ . Now let  $A = \{n \in \mathbb{Z} : n \equiv k \pmod{p}\}$ . It is easy to see that  $A$  contains no solutions to the equation  $\mathcal{L}$ , simply because it doesn't even contain any solutions modulo  $p$ . Hence, as above,  $\liminf f(n)/n \geq 1/p > 0$ .

**Q.6 (i)** In general, one can construct a subset of  $\mathbb{N}$  containing no solutions to some fixed equation by means of a *greedy choice procedure*, i.e.: go through the natural numbers in increasing order and, at each step, add the number to the set  $A$  if it doesn't create any solutions. For some equations (though by no means all !), this turns out to give an optimal construction. An example is just the equation  $2x = y$ . It is easy to check that the set  $A$

obtained by this greedy choice is the set

$$A = \{n \in \mathbb{N} : 2^{2k} \mid n \text{ for some } k \geq 0\}. \quad (16)$$

In other words,  $A$  consists of all those number which are divisible by an even power of 2. This set has an asymptotic density given by

$$d(A) = \frac{1}{2} \left( \sum_{k=0}^{\infty} \frac{1}{4^k} \right) = \frac{2}{3}.$$

No set avoiding solutions to  $2x = y$  can have a higher asymptotic upper density. For if  $B$  is such a set, then  $B \cap 2B = \phi$ , where  $2B = \{2b : b \in B\}$ . Thus  $1 \geq \bar{d}(B \cup 2B) = \bar{d}(B) + \bar{d}(2B) = \frac{3}{2}\bar{d}(B)$ , which implies that  $\bar{d}(B) \leq 2/3$ .

From the above reasoning, it is easy to deduce that  $\lim_{n \rightarrow \infty} f(n)/n = 2/3$ . In fact, if one argues more carefully, one can in fact show that, for every  $n \in \mathbb{N}$ , the set  $A \cap \{1, \dots, n\}$ , where  $A$  is given by (16), is a subset of  $\{1, \dots, n\}$  of largest possible size containing no solutions to  $2x = y$ .

(ii) The limit is at least  $1/2$ , because the set of odd numbers avoids solutions to  $3x = y + z$ . In fact, the limit is  $1/2$ . See the 3-page supplementary document on the homepage for a proof of a more precise result. For generalisations of the result, see Papers No. 25 and 29 on my research page.

**Q.7** Factorise an integer  $n$  as

$$n = \prod_{i=1}^k p_i^{\alpha_i}.$$

Then

$$d(n) = \prod_{i=1}^k (\alpha_i + 1),$$

and hence

$$\frac{d(n)}{n^\epsilon} = \prod_{i=1}^k \frac{\alpha_i + 1}{p_i^{\epsilon \alpha_i}}.$$

Fix  $\epsilon > 0$ . For a prime  $p$  and a positive integer  $\alpha$ , let  $f_\epsilon(p, \alpha) := \frac{\alpha+1}{p^{\epsilon \alpha}}$ . For fixed  $\epsilon$ , it is clear that  $f_\epsilon(p, \alpha) \rightarrow 0$  as long as any one of  $p$  and  $\alpha$  goes to infinity, even if the other is held fixed. In particular, there will be only a finite number (depending on  $\epsilon$ ) of pairs  $(p, \alpha)$  such that  $f_\epsilon(p, \alpha) > 1/2$ . It follows easily that  $d(n)/n^\epsilon \rightarrow 0$  as  $n \rightarrow \infty$ .

(NOTE : You can replace  $1/2$  by any number strictly less than one and the argument will work).



**Q.8** It suffices to show that, for any  $r > 1$  and  $\varepsilon > 0$  sufficiently small, there exist primes  $p, q$  such that  $(1 - 2\varepsilon)r < p/q < (1 + 3\varepsilon)r$ . By the prime number theorem,

$$\begin{aligned} \pi((1 - \varepsilon)x) &\sim \frac{(1 - \varepsilon)x}{\log(1 - \varepsilon)x}, & \pi((1 + \varepsilon)x) &\sim \frac{(1 + \varepsilon)x}{\log(1 + \varepsilon)x}, \\ \pi((1 - \varepsilon)rx) &\sim \frac{(1 - \varepsilon)rx}{\log(1 - \varepsilon)rx}, & \pi((1 + \varepsilon)rx) &\sim \frac{(1 + \varepsilon)rx}{\log(1 + \varepsilon)rx}. \end{aligned}$$

It follows that, for  $x$  sufficiently large and  $\varepsilon$  sufficiently small, there exist distinct primes  $p, q$  such that  $(1 - \varepsilon)x < q < (1 + \varepsilon)x$  and  $(1 - \varepsilon)rx < p < (1 + \varepsilon)rx$ . Thus  $r \left(\frac{1 - \varepsilon}{1 + \varepsilon}\right) < p/q < r \left(\frac{1 + \varepsilon}{1 - \varepsilon}\right)$ , which in turn implies that  $r(1 - 2\varepsilon) < p/q < (1 + 3\varepsilon)r$ , for sufficiently small  $\varepsilon$ .