# MVE165/MMG630, Applied Optimization Lecture 1

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ecture 1

Applied Optimization

### Staff and homepage

### Staff

- ► Examiner/lecturer: Ann-Brith Strömberg (anstr@chalmers.se, room L2087)
- ▶ Substitute lecturers: Peter Lindroth & Adam Wojciechowski
- ► Guest lecturers: Fredrik Hedenus (Energy and Environment), Michael Patriksson (Mathematical Sciences), Elin Svensson (Energy and Environment), and Caroline Olsson (Radiation Physics, Clinical Sciences)

### Course homepage

http://www.math.chalmers.se/Math/Grundutb/CTH/mve165/0809/

- ► Contains details, information on assignments and exercises, deadlines, lecture notes, etc
- Will be updated every week during the course



#### 107107

### Course contents and organization

#### Contents

- ► Applications of optimization
- ▶ Mathematical modelling
- ► Solution techniques algorithms
- ► Software solvers

### Organization

- ► Lectures mathematical optimization theory
- ► Exercises use solvers, oral examination (or report)
- ► Guest lectures applications of optimization
- ► Assignments modelling, use solvers, written reports, opposition & oral presentations
- ► Assignment work should be done in groups of two persons



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### Literature and examination

#### Literature

- ▶ R.L. Rardin: "Optimization in Operations Research", published by Prentice-Hall, New Jersey (1998)
- ► Found at Cremona (by the end of this week)

### Examination

- ► Two correctly solved computer exercises (oral examination or written reports)
- ▶ Written reports of three assignments (1, 2, and 3a or 3b)
- ▶ A written opposition to Assignment 2
- ▶ An oral presentation of Assignment 3a or 3b
- ▶ To be able to receive a grade higher than 3 or G, the written reports and opposition as well as the oral presentation must be of high quality. Students aiming at grade 4, 5, or VG must also pass an oral exam



### Overview of the lectures

- ▶ Linear programming, modelling, theory, solution methods, sensitivity analysis
- ▶ Optimization models that can be described as flows in networks, solution methods
- ▶ Discrete optimization models and solution methods
- ▶ Non-linear programming models, with(out) constraints, solution methods
- ► Multiple objective optimization
- ► Optimization under uncertainty
- Mixtures of the above

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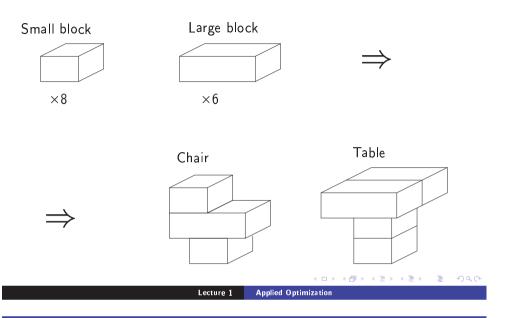
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### Optimization

"Do something as good as possible"

- **Something:** Which are the decision alternatives?
- **Possible:** What restrictions are there?
- ▶ **Good:** What is a relevant optimization criterion?

# A manufacturing example: Produce tables and chairs from two types of blocks



# A manufacturing example, continued

- ▶ A chair is assembled from one large and two small blocks
- ▶ A table is assembled from two blocks of each size
- ▶ Only 6 large and 8 small blocks are avaliable
- ▶ A table is sold at a revenue of 1600:-
- ▶ A chair is sold at a revenue of 1000:-
- Assume that all items produced can be sold and determine an optimal production plan



### A mathematical optimization model

Something: Which are the decision alternatives? ⇒ Variables

 $x_1$  = number of tables produced and sold  $x_2$  = number of chairs produced and sold

Possible: What restrictions are there? ⇒ Constraints

$$2x_1 + x_2 \le 6$$
 (6 large blocks)  
 $2x_1 + 2x_2 \le 8$  (8 small blocks)  
 $x_1, x_2 \ge 0$  (physical restrictions)  
 $(x_1, x_2 \text{ integral})$  (physical restrictions)

Good: What is a relevant optimization criterion? ⇒ Objective function

maximize 
$$z = 1600x_1 + 1000x_2$$
 ( $z = \text{total revenue}$ )

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# Solve the model using LEGO!

- Start at no production:  $x_1 = x_2 = 0$ Use the "best marginal profit" to choose the item to produce
  - ▶  $x_1$  has the highest marginal profit (1600:-/table) ⇒ produce as many tables as possible
  - At  $x_1 = 3$ : no more large blocks left
- ► The marginal value of  $x_2$  is now 200:- since taking apart one table (-1600:-) yields two chairs  $(+2000:-) \Rightarrow 400:-/2$  chairs
  - ▶ Increase  $x_2$  maximally  $\Rightarrow$  decrease  $x_1$
  - At  $x_1 = x_2 = 2$ : no more small blocks
- ▶ The marginal value of  $x_1$  is negative (to build one more table one has to take apart two chairs  $\Rightarrow -400$ :-) The marginal value of  $x_2$  is -600:- (to build one more chair

The marginal value of  $x_2$  is -600:- (to build one more chair one table must be taken apart)

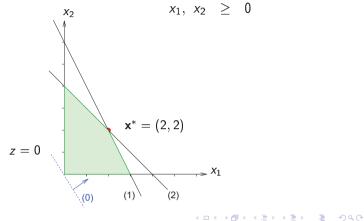
 $\implies$  Optimal solution:  $x_1 = x_2 = 2$ 

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### Geometrical solution of the model

maximize 
$$z = 1600x_1 + 1000x_2$$
 (0) subject to  $2x_1 + x_2 \le 6$  (1)

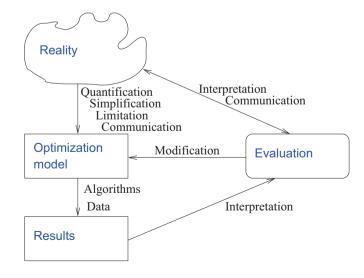
$$2x_1 + 2x_2 \leq 8 \qquad (2)$$



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# Operations research—more than just mathematics



### Concepts

- ▶ Feasible solution: satisfies all constraints
- ▶ **Optimal solution:** feasible AND objective function value as good as for every feasible solution
- ► **Sensitivity analysis:** how the solution depend on input parameters
- ▶ Tractability: Can the the model be solved in reasonable time?
- ► Validity: Does the conclusions drawn from the solution hold for the REAL problem
- ▶ **Operations research:** Always a tradeoff between validity of the model and its tractability to analysis

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### More concepts

- ▶ **Optimal solution:** proven to be as good as any other feasible solution
- ► Heuristic or approximate solution: feasible, not guaranteed to be exactly optimal, quality measures can be computed
- ▶ **Deterministic optimization model:** All parameter values assumed known with certainty
- ► **Stochastic optimization model:** involves quantities known only in probability (optimization under uncertainty)
- ► Multiple objective optimization: typically no feasible solution exist that is optimal in ALL objectives, search for Pareto optimal solutions

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- ▶ Deliver windows over a six-month period
- ▶ Demand for each month: 100, 250, 190, 140, 220 & 110 units
- ▶ Production cost/window: 50 €,45 €, 55 €, 48 €, 52 € & 50 €
- ▶ Store a produced window from one month to the next at 8 €
- Meet the demands and minimize costs
- ▶ Find an optimal production schedule

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### Define the decision variables

 $x_i$  = number of units produced in month i = 1, ..., 6

 $y_i$  = units left in the inventory at the end of month  $i=1,\ldots,6$ 

▶ The "flow" of windows can be illustrated as:

# Define the limitations/constraints

#### ► Each month:

initial inventory + production - ending inventory = demand

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### Objective function: minimize the costs

- ► Production cost ( $\in$ ):  $50 \ x_1 + 45 \ x_2 + 55 \ x_3 + 48 \ x_4 + 52 \ x_5 + 50 \ x_6$
- ► Inventory cost ( $\in$ ): 8 ( $y_1 + y_2 + y_3 + y_4 + y_5 + y_6$ )
- Objective:

minimize 
$$50x_1 + 45x_2 + 55x_3 + 48x_4 + 52x_5 + 50x_6 + 8(y_1 + y_2 + y_3 + y_4 + y_5 + y_6)$$

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minimize 
$$\sum_{i=1}^{6} c_i x_i + 8 \sum_{i=1}^{6} y_i$$
,  
subject to  $y_{i-1} + x_i - y_i = d_i$ ,  $i = 1, \dots, 6$ ,  
 $y_0 = 0$ ,  
 $x_i, y_i \geq 0$ ,  $i = 1, \dots, 6$ ,

The vector of demand:

$$d = (d_i)_{i=1}^6 = (100, 250, 190, 140, 220, 110)$$

The vector of production costs:

$$c = (c_i)_{i=1}^6 = (50, 45, 55, 48, 52, 50)$$

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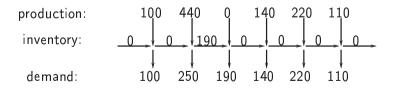
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# An optimal solution—optimal production schedule

Optimal production:  $x = (x_i)_{i=1}^6 = (100, 440, 0, 140, 220, 110)$ 

Optimal inventory:  $y = (y_i)_{i=0}^6 = (0, 0, 190, 0, 0, 0, 0)$ 



The minimal total cost is 49980 €

# Mathematical optimization models

- $\triangleright$   $x_1, \ldots, x_n$  are the decision variables
- f and  $g_1, \ldots, g_m$  are given functions of the decision variables
- $ightharpoonup b_1, \ldots, b_m$  are specified constant parameters
- ► The functions can be nonlinear, e.g. quadratic, exponential, logarithmic, non-analytic, ...
- ▶ In general, linear forms are more tractable than non-linear



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# Linear programming models

- ► The production inventory model is a linear program (LP) all relations are described by linear forms
- ► A general linear program:

$$\left[\begin{array}{cccc} \text{min or max} & c_1x_1+\ldots+c_nx_n \\ \\ \text{subject to} & a_{i1}x_1+\ldots+a_{in}x_n & \left\{\begin{array}{c} \leq \\ \geq \end{array}\right\} & b_i, \quad i=1,\ldots,m \\ \\ x_j & \geq & 0, \quad j=1,\ldots,n \end{array}\right]$$

▶ The non-negativity constraints on  $x_j$ , j = 1, ..., n are not necessary, but usually assumed (reformulation always possible)



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### Discrete/integer/binary modelling

- ► A variable is discrete if it can take only a countable set of values, e.g.,
  - ▶ Discrete variable:  $x \in \{0, 4.4, 5.2, 8.0\}$
  - ▶ Integer variable:  $x \in \{0, 1, 4, 5, 8\}$
- ► A binary variable can only take values 0 or 1 all or nothing E.g., a wind-mill can produce electricity only if it is built
  - ▶ Let y = 1 if the mill is built, else y = 0
  - ► Capacity of a mill: C
  - ▶ Production  $x \le C \cdot y$  (also limited by wind force etc.)
- ▶ In general, models with only continuous variables are more tractable than models with integrality/discrete requirements on variables, but exceptions exist!



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