MVE165/MMG630, Applied Optimization
Lecture 2
Convexity; basic feasible solutions; the simplex method; degeneracy; unbounded solutions

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2011–03–24

Mathematical optimization models

$$\left[\begin{array}{lll} \text{minimize or maximize} & f(x_1,\ldots,x_n) & \\ \text{subject to} & g_i(x_1,\ldots,x_n) \left\{\begin{array}{c}\le\\=\\\ge\end{array}\right\} b_i, & i=1,\ldots,m \end{array}\right]$$

- $x_1,\ldots,x_n$ are the decision variables
- f and $g_1,\ldots,g_m$ are given functions of the decision variables
- $b_1,\ldots,b_m$ are specified constant parameters
- The functions can be nonlinear, e.g. quadratic, exponential, logarithmic, non-analytic, ...
- In general, linear forms are more tractable than non-linear

Linear optimization models (programs)

- The production inventory model is a linear program (LP), i.e., all relations are described by linear forms

- A general linear program:

$$\left[\begin{array}{lccccc} \text{min or max} & c_1x_1 + \ldots + c_nx_n & & & & \\ \text{subject to} & a_{i1}x_1 + \ldots + a_{in}x_n & \left\{ \begin{array}{c} \leq \\ = \\ \geq \end{array} \right\} & b_i, & i = 1, \ldots, m \\ & x_j & \geq & 0, & j = 1, \ldots, n \end{array}\right]$$

- The non-negativity constraints on x_j, $j = 1, \ldots, n$ are not necessary, but usually assumed (reformulation always possible)

Discrete/integer/binary modelling

- A variable is called *discrete* if it can take only a countable set of values, e.g.,
  - Continuous variable: $x \in [0, 8]$ or $0 \leq x \leq 8$
  - Discrete variable: $x \in \{0, 4.4, 5.2, 8.0\}$
  - *Integer* variable: $x \in \{0, 1, 4, 5, 8\}$
- A *binary* variable can only take the values 0 or 1, i.e., all or nothing

  E.g., a wind-mill can produce electricity only if it is built
  - Let $y = 1$ if the mill is built, otherwise $y = 0$
  - Capacity of a mill: C
  - Production $x \leq Cy$ (also limited by wind force etc.)
- In general, models with only continuous variables are more tractable than models with integrality/discrete requirements on the variables, but exceptions exist!

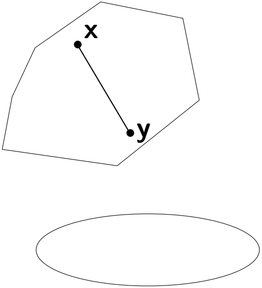
Convex sets

- A set S is convex if, for any elements $\mathbf{x}, \mathbf{y} \in S$ it holds that

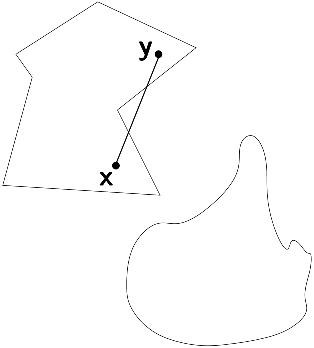
$$\alpha \mathbf{x} + (1-\alpha)\mathbf{y} \in S \text{ for all } 0 \leq \alpha \leq 1$$

- Examples:

Convex sets



Non-convex sets



$\Rightarrow$ Intersections of linear (in)equalities $\Rightarrow$ convex sets

Convex and concave functions

- A function f is convex on the set S if, for any elements $\mathbf{x}, \mathbf{y} \in S$ it holds that

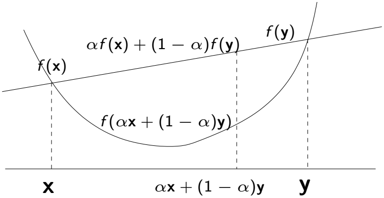
$$f(\alpha\mathbf{x} + (1-\alpha)\mathbf{y}) \leq \alpha f(\mathbf{x}) + (1-\alpha)f(\mathbf{y}) \text{ for all } 0 \leq \alpha \leq 1$$

- A function f is concave on the set S if, for any elements $\mathbf{x}, \mathbf{y} \in S$ it holds that

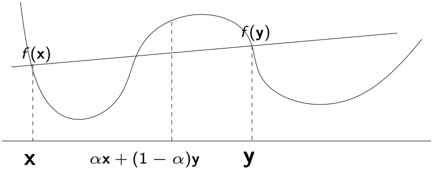
$$f(\alpha\mathbf{x} + (1-\alpha)\mathbf{y}) \geq \alpha f(\mathbf{x}) + (1-\alpha)f(\mathbf{y}) \text{ for all } 0 \leq \alpha \leq 1$$

$\Rightarrow$ Linear functions are convex (and concave)

Convex function



Non-convex function



Global solutions of convex optimization problems (Ch. 11)

- Let $\mathbf{x}^*$ be a *local* minimizer of a *convex function* over a *convex set*. Then $\mathbf{x}^*$ is also a *global* minimizer.

$\Rightarrow$ Every local optimum of a linear optimization problem is a global optimum

- If a linear optimization problem has any optimal solutions, at least one optimal solution is at an extreme point of the feasible set

$\Rightarrow$ Search for optimal extreme point(s)

- Next: Linear optimization problems and the simplex method

A general linear program – notation

$$\text{minimize or maximize} \quad c_1x_1 + \ldots + c_nx_n$$

$$\text{subject to} \quad a_{i1}x_1 + \ldots + a_{in}x_n \quad \left\{ \begin{array}{c} \leq \\ = \\ \geq \end{array} \right\} \quad b_i, \quad i = 1, \ldots, m$$

$$x_j \quad \left\{ \begin{array}{l} \leq 0 \\ \text{unrestricted in sign} \\ \geq 0 \end{array} \right\}, \quad j = 1, \ldots, n$$

- c_j, a_{ij}, and b_i are constant parameters for $i = 1, \ldots, m$ and $j = 1, \ldots, n$

The standard form and the simplex method for linear programs

- Every linear program can be reformulated such that:
  - all constraints are expressed as equalities with non-negative right hand sides
  - all variables are restricted to be non-negative
- Referred to as the *standard form*
- These requirements streamline the calculations of the *simplex method*
- Software solvers (e.g., Cplex, GLPK, Clp) can handle also inequality constraints and unrestricted variables – the reformulations are made automatically

The simplex method—reformulations

- The lego example:

$$\left[\begin{array}{rrrr} 2x_1 & +x_2 & \leq & 6 \\ 2x_1 & +2x_2 & \leq & 8 \\ & x_1, x_2 & \geq & 0 \end{array}\right] \Leftrightarrow \left[\begin{array}{rrrrrr} 2x_1 & +x_2 & +s_1 & & = & 6 \\ 2x_1 & +2x_2 & & +s_2 & = & 8 \\ & & & x_1, x_2, s_1, s_2 & \geq & 0 \end{array}\right]$$

- s_1 and s_2 are called *slack variables*—they "fill out" the (positive) distances between the left and right hand sides
- *Surplus variable* s_3 (a different example):

$$\left[\begin{array}{rrrrr} x_1 & + & x_2 & \geq & 800 \\ & & x_1, x_2 & \geq & 0 \end{array}\right] \Leftrightarrow \left[\begin{array}{rrrrrrr} x_1 & + & x_2 & - & s_3 & = & 800 \\ & & & & x_1, x_2, s_3 & \geq & 0 \end{array}\right]$$

The simplex method—reformulations, cont.

- Non-negative right hand side:

$$\begin{bmatrix} x_1 - x_2 & \leq -23 \\ x_1, x_2 & \geq 0 \end{bmatrix} \Leftrightarrow \begin{bmatrix} -x_1 + x_2 & \geq 23 \\ x_1, x_2 & \geq 0 \end{bmatrix} \Leftrightarrow \begin{bmatrix} -x_1 + x_2 - s_4 & = 23 \\ x_1, x_2, s_4 & \geq 0 \end{bmatrix}$$

- Sign-restricted (non-negative) variables:

$$\begin{bmatrix} x_1 + x_2 & \leq 10 \\ x_1 & \geq 0 \end{bmatrix} \Leftrightarrow \begin{bmatrix} x_1 + x_2^1 - x_2^2 & \leq 10 \\ x_1, x_2^1, x_2^2 & \geq 0 \end{bmatrix} \Leftrightarrow \begin{bmatrix} x_1 + x_2^1 - x_2^2 + s_5 & = 10 \\ x_1, x_2^1, x_2^2, s_5 & \geq 0 \end{bmatrix}$$

Basic feasible solutions

- Consider m equations of n variables, where $m \leq n$
- Set $n - m$ variables to zero and solve (if possible) the remaining $(m \times m)$ system of equations
- If the solution is *unique*, it is called a *basic* solution
- A basic solution corresponds to an *intersection* (feasible $(x \geq 0)$ or infeasible $(x \ngeq 0)$) of m hyperplanes in $\Re^m$
- Each *extreme point* of the feasible set is an intersection of m hyperplanes such that all variable values are ≥ 0
- Basic feasible solution $\Leftrightarrow$ extreme point of the feasible set

$$
\begin{aligned}
a_{11}x_1 + \ldots + a_{1n}x_n &= b_1 & \qquad x_1 &\geq 0 \\
a_{21}x_1 + \ldots + a_{2n}x_n &= b_2 & \qquad x_2 &\geq 0 \\
&\cdots & &\cdots \\
a_{m1}x_1 + \ldots + a_{mn}x_n &= b_m & \qquad x_n &\geq 0
\end{aligned}
$$

Basic feasible solutions, example

- Constraints:

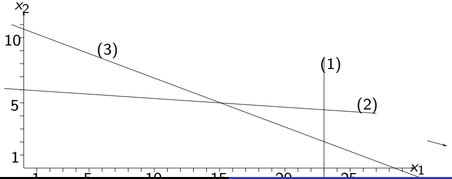
$$\begin{array}{rcrcll} x_1 & & & \leq & 23 & (1) \\ 0.067x_1 & + & x_2 & \leq & 6 & (2) \\ 3x_1 & + & 8x_2 & \leq & 85 & (3) \\ & & x_1, x_2 & \geq & 0 & \end{array}$$

- Add slack variables:

$$\begin{array}{rrrrrll} x_1 & & +s_1 & & & = 23 & (1) \\ 0.067x_1 & +x_2 & & +s_2 & & = 6 & (2) \\ 3x_1 & +8x_2 & & & +s_3 & = 85 & (3) \\ & & & & x_1, x_2, s_1, s_2, s_3 & \geq 0 & \end{array}$$

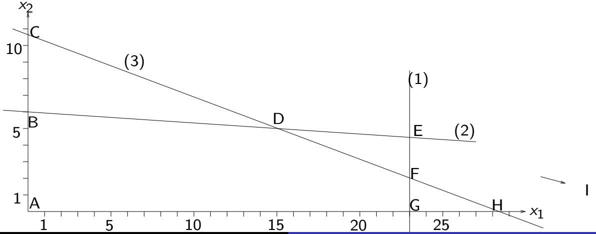
$m = 3$

$n = 5$

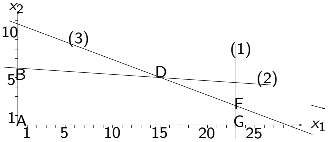


Basic and non-basic variables and solutions

basic variables	basic solution			non-basic variables $(0,0)$	point	feasible?
s_1, s_2, s_3	23	6	85	x_1, x_2	A	yes
s_1, s_2, x_1	$-5\frac{1}{3}$	$4\frac{1}{9}$	$28\frac{1}{3}$	s_3, x_2	H	no
s_1, s_2, x_2	23	$-4\frac{5}{8}$	$10\frac{5}{8}$	x_1, s_3	C	no
s_1, x_1, s_3	-67	90	-185	s_2, x_2	I	no
s_1, x_2, s_3	23	6	37	s_2, x_1	B	yes
x_1, s_2, s_3	23	$4\frac{7}{15}$	16	s_1, x_2	G	yes
x_2, s_2, s_3	-	-	-	s_1, x_1	-	-
x_1, x_2, s_1	15	5	8	s_2, s_3	D	yes
x_1, x_2, s_2	23	2	$2\frac{7}{15}$	s_1, s_3	F	yes
x_1, x_2, s_3	23	$4\frac{7}{15}$	$-19\frac{11}{15}$	s_1, s_2	E	no



Basic feasible solutions correspond to solutions to the system of equations that fulfil non-negativity



$$\left[\begin{array}{rrrrrl} x_1 & & +s_1 & & & = 23 \\ 0.067x_1 & +x_2 & & +s_2 & & = 6 \\ 3x_1 & +8x_2 & & & +s_3 & = 85 \end{array}\right]$$

A: $x_1 = x_2 = 0 \Rightarrow \left[\begin{array}{llll} s_1 & & & = 23 \\ & s_2 & & = 6 \\ & & s_3 & = 85 \end{array}\right]$

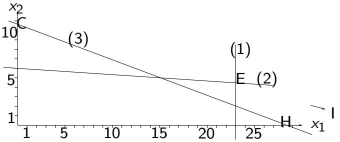
B: $x_1 = s_2 = 0 \Rightarrow \left[\begin{array}{rlll} & s_1 & & = 23 \\ x_2 & & & = 6 \\ 8x_2 & & +s_3 & = 85 \end{array}\right]$

D: $s_3 = s_2 = 0 \Rightarrow \left[\begin{array}{rrrl} x_1 & & +s_1 & = 23 \\ 0.067x_1 & +x_2 & & = 6 \\ 3x_1 & +8x_2 & & = 85 \end{array}\right]$

F: $s_3 = s_1 = 0 \Rightarrow \left[\begin{array}{rrrl} x_1 & & & = 23 \\ 0.067x_1 & +x_2 & +s_2 & = 6 \\ 3x_1 & +8x_2 & & = 85 \end{array}\right]$

G: $x_2 = s_1 = 0 \Rightarrow \left[\begin{array}{rrrl} x_1 & & & = 23 \\ 0.067x_1 & +s_2 & & = 6 \\ 3x_1 & & +s_3 & = 85 \end{array}\right]$

Basic infeasible solutions corresp. to solutions to the system of equations with one or more variables < 0



$$\left[\begin{array}{rrrrrl} x_1 & & +s_1 & & & = 23 \\ 0.067x_1 & +x_2 & & +s_2 & & = 6 \\ 3x_1 & +8x_2 & & & +s_3 & = 85 \end{array}\right]$$

H: $x_2 = s_3 = 0 \Rightarrow \left[\begin{array}{rrrl} x_1 & +s_1 & & = 23 \\ 0.067x_1 & & +s_2 & = 6 \\ 3x_1 & & & = 85 \end{array}\right]$

C: $x_1 = s_3 = 0 \Rightarrow \left[\begin{array}{rrrl} & s_1 & & = 23 \\ x_2 & & +s_2 & = 6 \\ 8x_2 & & & = 85 \end{array}\right]$

I: $s_2 = x_2 = 0 \Rightarrow \left[\begin{array}{rrrl} x_1 & +s_1 & & = 23 \\ 0.067x_1 & & & = 6 \\ 3x_1 & & +s_3 & = 85 \end{array}\right]$

-: $s_1 = x_1 = 0 \Rightarrow \left[\begin{array}{rrrl} & & 0 & = 23 \\ x_2 & +s_2 & & = 6 \\ 8x_2 & & +s_3 & = 85 \end{array}\right]$

E: $s_1 = s_2 = 0 \Rightarrow \left[\begin{array}{rrrl} x_1 & & & = 23 \\ 0.067x_1 & +x_2 & & = 6 \\ 3x_1 & +8x_2 & +s_3 & = 85 \end{array}\right]$

Basic feasible solutions and the simplex method

- Express the m *basic* variables in terms of the $n - m$ *non-basic* variables
- Example: Start at $x_1 = x_2 = 0 \Rightarrow s_1, s_2, s_3$ are *basic*

$$\begin{bmatrix} x_1 & & +s_1 & & & = 23 \\ \frac{1}{15}x_1 & +x_2 & & +s_2 & & = 6 \\ 3x_1 & +8x_2 & & & +s_3 & = 85 \end{bmatrix}$$

- Express s_1, s_2, and s_3 in terms of x_1 and x_2 (*non-basic*):

$$\begin{bmatrix} s_1 = & 23 & -x_1 & \\ s_2 = & 6 & -\frac{1}{15}x_1 & -x_2 \\ s_3 = & 85 & -3x_1 & -8x_2 \end{bmatrix}$$

- We wish to maximize the objective function $2x_1 + 3x_2$
- Express the objective in terms of the *non-basic* variables:

$$z = 2x_1 + 3x_2 \qquad \Leftrightarrow \qquad z - 2x_1 - 3x_2 = 0$$

Basic feasible solutions and the simplex method

- The *first basic solution* can be represented as

$$
\begin{array}{rrrrrrl|c}
-z & +2x_1 & +3x_2 & & & & = 0 & (0) \\
 & x_1 & & +s_1 & & & = 23 & (1) \\
 & \frac{1}{15}x_1 & +x_2 & & +s_2 & & = 6 & (2) \\
 & 3x_1 & +8x_2 & & & +s_3 & = 85 & (3)
\end{array}
$$

- Marginal values for increasing the non-basic variables x_1 and x_2 from zero: 2 and 3, resp.

$\Rightarrow$ Choose x_2 — let x_2 *enter the basis* DRAW GRAPH!!

- One basic variable (s_1, s_2, or s_3) must *leave the basis*. Which?
- The value of x_2 can increase until some basic variable reaches the value 0:

$$
\left.\begin{array}{ll}
(2) : s_2 = 6 - x_2 \geq 0 & \Rightarrow x_2 \leq 6 \\
(3) : s_3 = 85 - 8x_2 \geq 0 & \Rightarrow x_2 \leq 10\frac{5}{8}
\end{array}\right\} \Rightarrow
\begin{array}{c}
s_2 = 0 \text{ when} \\
x_2 = 6 \\
(\text{and } s_3 = 37)
\end{array}
$$

- s_2 will leave the basis

Change basis through row operations

- Eliminate s_2 from the basis, let x_2 enter the basis using row operations:

$-z$	$+2x_1$	$+3x_2$				$=$	0	(0)
	x_1		$+s_1$			$=$	23	(1)
	$\frac{1}{15}x_1$	$+x_2$		$+s_2$		$=$	6	(2)
	$3x_1$	$+8x_2$			$+s_3$	$=$	85	(3)
$-z$	$+\frac{9}{5}x_1$			$-3s_2$		$=$	-18	(0) $-3\cdot$(2)
	x_1		$+s_1$			$=$	23	(1)$-0\cdot$(2)
	$\frac{1}{15}x_1$	$+x_2$		$+s_2$		$=$	6	(2)
	$\frac{37}{15}x_1$			$-8s_2$	$+s_3$	$=$	37	(3)$-8\cdot$(2)

- Corresponding basic solution: $s_1 = 23$, $x_2 = 6$, $s_3 = 37$.
- Nonbasic variables: $x_1 = s_2 = 0$
- The marginal value of x_1 is $\frac{9}{5} > 0$. Let x_1 enter the basis
- Which one should leave? s_1, x_2, or s_3?

Change basis ...

$-z$	$+\frac{9}{5}x_1$			$-3s_2$		$=$	-18	(0)
	x_1		$+s_1$			$=$	23	(1)
	$\frac{1}{15}x_1$	$+x_2$		$+s_2$		$=$	6	(2)
	$\mathbf{\frac{37}{15}x_1}$			$-8s_2$	$+s_3$	$=$	37	(3)

- The value of x_1 can increase until some basic variable reaches the value 0:

$$\left.\begin{array}{ll} (1): s_1 = 23 - x_1 \geq 0 & \Rightarrow x_1 \leq 23 \\ (2): x_2 = 6 - \frac{1}{15}x_1 \geq 0 & \Rightarrow x_1 \leq 90 \\ (3): s_3 = 37 - \frac{37}{15}x_1 \geq 0 & \Rightarrow x_1 \leq 15 \end{array}\right\} \Rightarrow \begin{array}{c} s_3 = 0 \text{ when} \\ x_1 = 15 \end{array}$$

- x_1 enters the basis and s_3 leaves the basis
- Perform row operations:

$-z$				$+2.84s_2$	$-0.73s_3$	$=$	-45	$(0)-(3)\cdot\frac{15}{37}\cdot\frac{9}{5}$
			s_1	$+3.24s_2$	$-0.41s_3$	$=$	8	$(1)-(3)\cdot\frac{15}{37}$
		x_2		$+1.22s_2$	$-0.03s_3$	$=$	5	$(2)-(3)\cdot\frac{15}{37}\cdot\frac{1}{15}$
	x_1			$-3.24s_2$	$+0.41s_3$	$=$	15	$(3)\cdot\frac{15}{37}$

Change basis ...

$-z$			$+2.84s_2$	$-0.73s_3$	$=$	-45	(0)
		s_1	$+\mathbf{3.24s_2}$	$-0.41s_3$	$=$	8	(1)
	x_2		$+1.22s_2$	$-0.03s_3$	$=$	5	(2)
x_1			$-3.24s_2$	$+0.41s_3$	$=$	15	(3)

- Let s_2 enter the basis (marginal value > 0)
- The value of s_2 can increase until some basic variable $= 0$:

$$\left.\begin{array}{ll} (1): s_1 = 8 - 3.24s_2 \geq 0 & \Rightarrow s_2 \leq 2.47 \\ (2): x_2 = 5 - 1.22s_2 \geq 0 & \Rightarrow s_2 \leq 4.10 \\ (3): x_1 = 15 + 3.24s_2 \geq 0 & \Rightarrow s_2 \geq -4.63 \end{array}\right\} \Rightarrow \begin{array}{c} s_1 = 0 \text{ when} \\ s_2 = 2.47 \end{array}$$

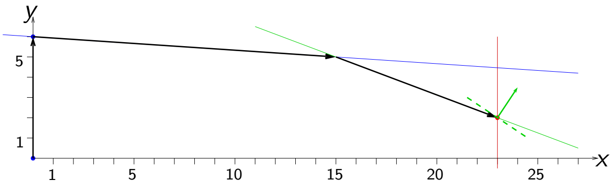
- s_2 enters the basis and s_1 will leave the basis
- Perform row operations:

$-z$			$-0.87s_1$		$-0.37s_3$	$=$	-52	$(0)-(1)\cdot\frac{2.84}{3.24}$
			$0.31s_1$	$+s_2$	$-0.12s_3$	$=$	2.47	$(1)\cdot\frac{1}{3.24}$
	x_2		$-0.37s_1$		$+0.12s_3$	$=$	2	$(2)-(1)\cdot\frac{1.22}{3.24}$
x_1			$+s_1$			$=$	23	$(3)+(1)$

Optimal basic solution

$-z$			$-0.87s_1$		$-0.37s_3$	$= -52$
			$0.31s_1$	$+s_2$	$-0.12s_3$	$= 2.47$
		x_2	$-0.37s_1$		$+0.12s_3$	$= 2$
	x_1		$+s_1$			$= 23$

- No marginal value is positive. No improvement can be made
- The optimal basis is given by $s_2 = 2.47$, $x_2 = 2$, and $x_1 = 23$
- Non-basic variables: $s_1 = s_3 = 0$
- Optimal value: $z = 52$



Summary of the solution course

basis	$-z$	x_1	x_2	s_1	s_2	s_3	RHS
$-z$	1	2	3	0	0	0	0
s_1	0	1	0	1	0	0	23
s_2	0	0.067	**1**	0	1	0	6
s_3	0	3	8	0	0	1	85
$-z$	1	1.80	0	0	-3	0	-18
s_1	0	1	0	1	0	0	23
x_2	0	0.07	1	0	1	0	6
s_3	0	**2.47**	0	0	-8	1	37
$-z$	1	0	0	0	2.84	-0.73	-45
s_1	0	0	0	1	**3.24**	-0.41	8
x_2	0	0	1	0	1.22	-0.03	5
x_1	0	1	0	0	-3.24	0.41	15
$-z$	1	0	0	-0.87	0	-0.37	-52
s_2	0	0	0	0.31	1	-0.12	2.47
x_2	0	0	1	-0.37	0	0.12	2
x_1	0	1	0	1	0	0	23

Summary of the simplex method

- **Optimality condition**: The *entering* variable in a maximization (minimization) problem should have the largest positive (negative) marginal value (reduced cost).

  The entering variable *determines a direction* in which the objective value increases (decreases).

  If all *reduced costs are negative* (positive), the current basis is *optimal*.

- **Feasibility condition**: The *leaving* variable is the one with smallest nonnegative quotient.

  Corresponds to the constraint that is "reached first"

Simplex search for linear (minimization) programs (Ch. 4.6)

1. **Initialization:** Choose any feasible basis, construct the corresponding basic solution $\mathbf{x}^0$, let $t = 0$
2. **Step direction:** Select a variable to enter the basis using the optimality condition (negative marginal value). Stop if no entering variable exists
3. **Step length:** Select a leaving variable using the feasibility condition (smallest non-negative quotient)
4. **New iterate:** Compute the new basic solution $\mathbf{x}^{t+1}$ by performing matrix operations.
5. Let $t := t + 1$ and repeat from 2

Solve the lego problem using the simplex method!

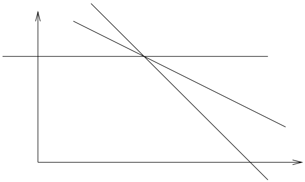
$$\begin{array}{lrcrcrcl} \text{maximize} & z & = & 1600x_1 & + & 1000x_2 & & \\ \text{subject to} & & & 2x_1 & + & x_2 & \leq & 6 \\ & & & 2x_1 & + & 2x_2 & \leq & 8 \\ & & & & & x_1,\ x_2 & \geq & 0 \end{array}$$

HOMEWORK!!

Degeneracy (Ch. 4.10)

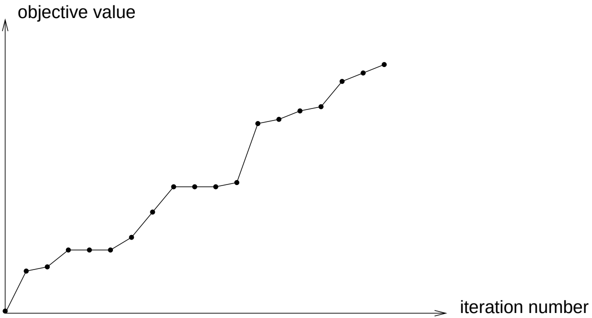
- If the smallest nonnegative quotient is zero, the value of a basic variable will become zero in the next iteration
- The solution is *degenerate*
- The objective value will *not* improve in this iteration
- Risk: *cycling* around (non-optimal) bases
- Reason: a *redundant* constraint "touches" the feasible set
- Example:

$$\begin{array}{rcrcl} x_1 & + & x_2 & \leq & 6 \\ & & x_2 & \leq & 3 \\ x_1 & + & 2x_2 & \leq & 9 \\ & & x_1,\ x_2 & \geq & 0 \end{array}$$



Degeneracy

- Typical objective function progress of the simplex method



- Computation rules to prevent from infinite cycling: careful choices of leaving and entering variables
- In modern software: perturb the right hand side $(b_i + \Delta b_i)$, solve, reduce the perturbation and resolve starting from the current basis. Repeat until $\Delta b_i = 0$.

Unbounded solutions (Ch. 4.4, 4.6)

- If all quotients are *negative*, the value of the variable entering the basis may increase *infinitely*
- The feasible set is *unbounded*
- In a real application this would probably be due to some incorrect assumption
- Example:

$$\begin{array}{lrrrl} \text{minimize} & z = & -x_1 & -2x_2 & \\ \text{subject to} & & -x_1 & +x_2 & \leq 2 \\ & & -2x_1 & +x_2 & \leq 1 \\ & & & x_1,\ x_2 & \geq 0 \end{array}$$

Draw graph!!

Unbounded solutions (Ch. 4.4, 4.6)

- A feasible basis is given by $x_1 = 1$, $x_2 = 3$, with corresponding tableau:

  Homework: Find this basis using the simplex method.

basis	$-z$	x_1	x_2	s_1	s_2	RHS
$-z$	1	0	0	5	-3	7
x_1	0	1	0	1	-1	1
x_2	0	0	1	2	-1	3

- Entering variable is s_2
- Row 1: $x_1 = 1 + s_2 \geq 0 \Rightarrow s_2 \geq -1$
- Row 2: $x_2 = 3 + s_2 \geq 0 \Rightarrow s_2 \geq -3$
- No leaving variable can be found, since no constraint will prevent s_2 from increasing infinitely