

# Maintenance scheduling optimization

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# Maintenance optimization — a background

- Invitation 2000 from Volvo Aero Corporation (VAC, nowadays GKN Aerospace): maintenance of the RM12 jet engine
- Paired PhD project between applied math/optimization and math statistics/material fatigue and reliability
- Optimization student: a model for opportunistic maintenance; superior to simpler policies
- Math statistics student: models for the determination of life distributions based on crack growth
- Continuation projects: GKN; planning maintenance of components in wind power plants and scheduling of rail grinding

# A conversation with Bo Hägg, CEO Underhållsföretagen

- Maintenance = obtain reliability at the least cost
- Maint. costs/year: 14K Billion SEK (EU), 275 Billion SEK (S)
- Maintenance is often seen merely as a cost
- Maintenance is sometimes done too often—inspections and measurements may damage the systems
- Sometimes—like with road/rail infrastructure and “Miljonprogramhusen”—it is performed seldom
- Truth: well performed maintenance is an investment in availability and safety



# Maintenance principles

- Preventive maintenance: actions that prevent failure
- Corrective maintenance: actions after failure, repairs
- Condition based maintenance: measurements → predictions  
→ actions according to a maintenance principle
- Opportunistic maintenance: when maintenance must be performed, make also some (additional) preventive maintenance actions

# A simple example, I

- A system with  $n$  components
- Life of component  $i$  :  $T_i$  time units (intervals)
- Time horizon:  $T$  time units (e.g. contract period)
- Cost of a spare component of type  $i$  at time  $t$ :  $c_{it}$  monetary units
- Cost for performing any maintenance at time  $t$ :  $d_t$  monetary units

## A simple example, II

- Variables are logical – do something or not
- Model uses binary variables:

$$x_t = \begin{cases} 1, & \text{if "something" is done at time } t \\ 0, & \text{otherwise} \end{cases}$$

- A decision often implies other necessary decisions
- Example: if component  $i$  shall be replaced at time  $t$  maintenance must be performed
- Such logical relations are equivalent to linear constraints:

$$\text{if A then B} \iff x_A \leq x_B$$

# The basic replacement problem, I

- Goal: minimize the total cost for a working system during the contract period:

## Mathematical model

$$\underset{(x,z)}{\text{minimize}} \quad \sum_{t=1}^T \left( \sum_{i=1}^N c_{it} x_{it} + d_t z_t \right), \quad (1a)$$

$$\text{subject to} \quad \sum_{t=\ell+1}^{\ell+T_i} x_{it} \geq 1, \quad \ell = 0, \dots, T - T_i, \quad i = 1, \dots, N, \quad (1b)$$

$$x_{it} \leq z_t, \quad t = 1, \dots, T, \quad i = 1, \dots, N, \quad (1c)$$

$$x_{it} \geq 0, \quad t = 1, \dots, T, \quad i = 1, \dots, N, \quad (1d)$$

$$z_t \leq 1, \quad t = 1, \dots, T, \quad (1e)$$

$$x_{it}, z_t \in \{0, 1\}, \quad t = 1, \dots, T, \quad i = 1, \dots, N \quad (1f)$$

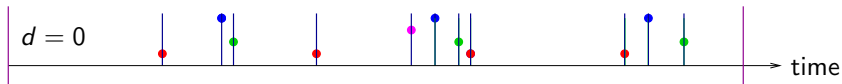
## The basic replacement problem, II

- Objective (1a): minimize the total cost of having a working system during the contract period
- Constraint (1b): for any given item  $i$  in the system, the component must be replaced at some point during *every* time interval of  $T_i$  time steps
- Constraint (1c): we cannot perform any replacement without paying the fixed cost  $d_t$  for performing a maintenance operation; once we do pay, any maintenance action becomes possible (at no extra fixed cost) at that time step
- Constraints (1d)–(1f) ensure that the variables take only meaningful values



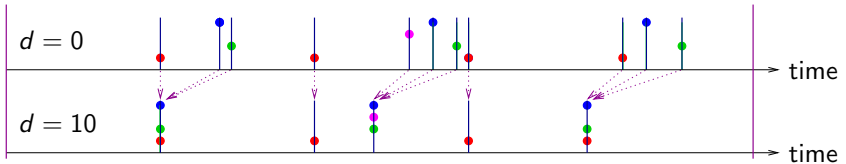
# Opportunistic maintenance or not?

- Example: four components with different prices and lives
- A replacement is marked with a dot; its colour represents the type of component replaced



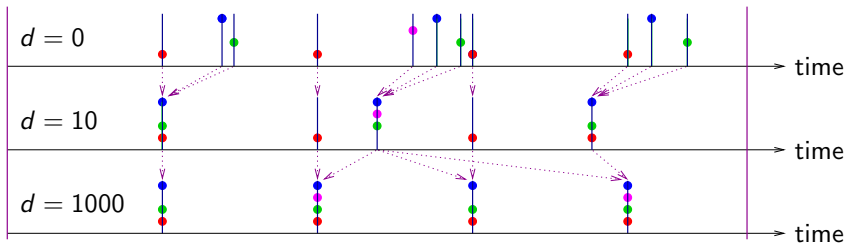
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# Opportunistic maintenance or not?

- Example: four components with different prices and lives
- A replacement is marked with a dot; its colour represents the type of component replaced
- The larger the fixed cost, the more beneficial is opportunistic maintenance becomes; also more items are replaced



# Constraint structure—example

- Component 3:  $\sum_{t=\ell+1}^{\ell+T_3} x_{3t} \geq 1, \quad \ell = 0, \dots, T - T_3$

- $T = 8, T_3 = 4 \implies \sum_{t=\ell+1}^{\ell+4} x_{3t} \geq 1, \quad \ell = 0, \dots, 4$

$$\begin{bmatrix} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} x_{31} \\ x_{32} \\ \vdots \\ x_{38} \end{bmatrix} \geq \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

# Property I: the replacement problem is NP-hard

## Theorem

*Set covering is polynomially reducible to the replacement problem*

- This essentially mean that we *cannot* expect to find an optimal solution in a time that is proportional to a polynomial function of the problem size ( $T(N + 1)$  variables and  $\approx 4NT$  constraints)
- Basic complexity theory: Chapter 2.6 in the course book

## Property II: with fixed $z$ the problem over $x$ is easy

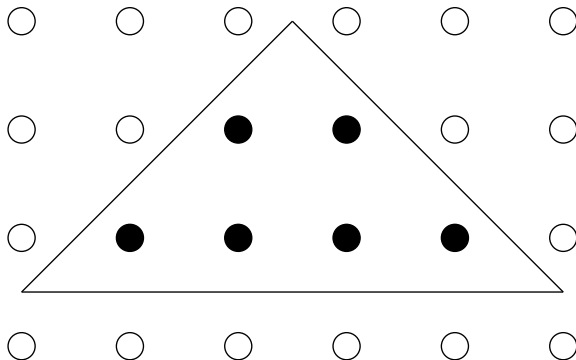
- The constraint matrix has the “consecutive ones” property
- ⇒ For fixed values of  $z$ , the problem over  $x$  can be solved as a linear program
- For each  $i$ , the linear programming dual problem can be solved by a “greedy” algorithm ⇒ primal solution by complementarity; see [a], Algorithm 1, page 297
- The latter is typically 5–40 times faster than solving as a general linear program, and 25–400 times faster when costs are monotone with time (i.e.,  $\forall t$  either  $c_{it} \leq c_{i,t+1}$  or  $c_{it} \geq c_{i,t+1}$ ); see [a], Algorithm 2, page 299

[a] T. Almgren, N. Andréasson, M. Patriksson, A.-B. Strömberg, A. Wojciechowski, M. Önnheim (2012): *The opportunistic replacement problem: theoretical analyses and numerical tests*, Mathematical Methods of Operations Research, 76(3) pp. 289–319.

<http://link.springer.com/article/10.1007%2Fs00186-012-0400-y>

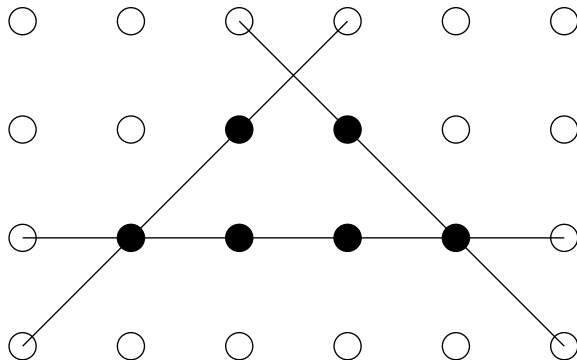
# Property III: all inequalities are facet defining

No inequalities are facet defining



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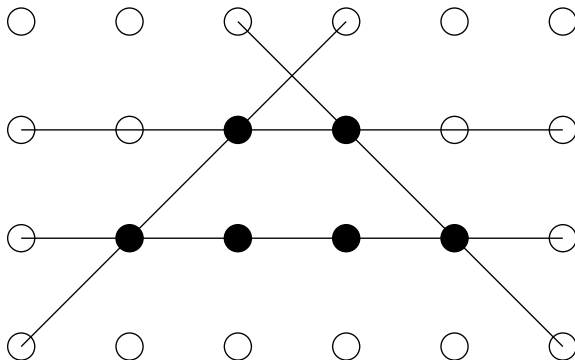
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# Property III: all inequalities are facet defining

Integral polyhedron



See [a], Section 5.1–5.2

# A generalized model

## New variable definition

Define the set

$$\mathcal{I} := \{ (s, t) \mid 0 \leq s < t \leq T + 1; s, t \in \mathbb{Z} \}$$

of replacement intervals and introduce the variables

$$x_{st}^i = \begin{cases} 1, & \text{if component } i \text{ receives PM at the} \\ & \text{times } s \text{ and } t, \text{ and not in-between,} \\ 0, & \text{otherwise,} \end{cases} \quad \begin{array}{l} i \in \mathcal{N}, \\ (s, t) \in \mathcal{I}, \end{array}$$

and

$$z_t = \begin{cases} 1, & \text{if maintenance occurs at time } t, \\ 0, & \text{otherwise,} \end{cases} \quad t \in \mathcal{T}.$$

# A generalized model

$$\text{minimize} \quad \sum_{t \in \mathcal{T}} d_t z_t + \sum_{i \in \mathcal{N}} \sum_{(s,t) \in \mathcal{I}} c_{st}^i x_{st}^i, \quad (2a)$$

$$\text{subject to} \quad \sum_{s=0}^{t-1} x_{st}^i \leq z_t, \quad i \in \mathcal{N}, t \in \mathcal{T}, \quad (2b)$$

$$\sum_{s=0}^{t-1} x_{st}^i = \sum_{r=t+1}^{T+1} x_{tr}^i, \quad i \in \mathcal{N}, t \in \mathcal{T}, \quad (2c)$$

$$\sum_{t=1}^{T+1} x_{0t}^i = 1, \quad i \in \mathcal{N}, \quad (2d)$$

$$x_{st}^i \in \{0, 1\}, \quad i \in \mathcal{N}, (s, t) \in \mathcal{I}, \quad (2e)$$

$$z_t \in \{0, 1\}, \quad t \in \mathcal{T}. \quad (2f)$$

# On the GKN project

- Aircraft engines are expensive:
  - Spare components cost up to 2 MSEK
  - Total cost of maintenance of one engine: 15–30 MSEK
  - Maximizing “time on wing” is important, both for civil and military aircraft
- The aircraft engine RM12 consists of 7 modules and 61 components in total
- A mathematical model has been constructed for the entire engine maintenance, including work costs for (dis)assembling the necessary modules and components for each maintenance occasion
- This model has slightly less than 6000 binary variables

## Results on the GKN problems

- An individual engine module with 10 components: cost reduction 35%; reduction of # maint. occasions 7% (compared with a simple policy similar to that used at GKN)
- Complete engine of 7 modules (61 components):
  - Cost reduction compared to maintaining (optimally) each individual module: 12%
  - Reduction of # maint. occasions: 60%
- Product development: found 5 components that can potentially reduce maintenance costs more than 5% through prolonged lives

# Maintenance of rails and wheels, I

- Paired PhD project in collaboration with CHARMEC
- Background: increased wear of rails and wheels due to an increase in speeds and loads
- Aim to develop decision support tools for the optimization of inspection and maintenance of rails and wheels wrt. LCC, safety and maintenance logistics
- The other PhD student (at CHARMEC) models the progressive degradation of rails and wheels
- Will result in advanced knowledge on how component condition indicators can be efficiently used in an optimization

# Maintenance of rails and wheels, II

- The picture below shows how a component of a rail is degraded over time (measured in portions of the rail sections having crack lengths in given intervals)
- Several levels of maintenance can be performed, with different effects (and different costs)
- The optimization will determine which maintenance action is the most appropriate at any given time

