Maintenance scheduling optimization

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Maintenance optimization — a background

- Invitation 2000 from Volvo Aero Corporation (VAC, nowadays GKN Aerospace): maintenance of the RM12 jet engine
- Paired PhD project between applied math/optimization and math statistics/material fatigue and reliability
- Optimization PhD student: a model for opportunistic maintenance; superior to simpler policies
- Mathematical statistics student: models for the determination of life distributions based on crack growth
- Continuation projects: GKN; planning maintenance of components in wind power plants and scheduling of rail grinding

A conversation with Bo Hägg, CEO Underhållsföretagen

- Optimal *maintenance* = obtain reliability at the least cost
- *Maintenance* costs/year:
  14000 billion SEK (in EU), 275 billion SEK (in Sweden)
- *Maintenance* is often seen merely as a cost
- *Maintenance* is sometimes done too often—inspections and measurements may damage the systems
- Sometimes—like with road/rail infrastructure and "Miljonprogramhusen'—it is performed seldom
- Truth: well performed *maintenance* is an investment in availability and safety

Maintenance principles

- Preventive maintenance (PM): actions that prevent failure

- Corrective maintenance (CM): actions after failure, repairs

- Condition based maintenance (CBM): measurements $\rightarrow$ predictions $\rightarrow$ actions according to a maintenance principle

- Opportunistic maintenance (OM): when maintenance must be performed, make also some (additional) preventive maintenance actions

A simple example, I

A system with n components

- Life of component i: T_i time units (intervals)

- Time horizon: T time units (e.g. contract period)

- Cost of a spare component of type i at time t: c_{it} monetary units

- Cost for performing any maintenance at time t: d_t monetary units

A simple example, II

Variables are logical – do something or not

The model uses binary variables:

$$x_t = \begin{cases} 1, & \text{if "something" is done at time } t \\ 0, & \text{otherwise} \end{cases}$$

A decision often implies other necessary decisions

- Example: if component i shall be replaced at time t maintenance must be performed
- Such logical relations are equivalent to linear constraints:

$$\text{if A then B} \qquad \Longleftrightarrow \qquad x_A \leq x_B$$

The basic replacement problem, I

Goal: minimize the total cost for keeping the system working during the contract period:

Mathematical model

$$
\operatorname*{minimize}_{(x,z)} \quad \sum_{t=1}^{T} \left(\sum_{i=1}^{N} c_{it} x_{it} + d_t z_t \right), \tag{1a}
$$

$$
\text{subject to} \quad \sum_{t=\ell+1}^{\ell+T_i} x_{it} \geq 1, \qquad \ell = 0, \ldots, T-T_i,\ i = 1, \ldots, N, \tag{1b}
$$

$$
x_{it} \leq z_t, \qquad t = 1, \ldots, T,\ i = 1, \ldots, N, \tag{1c}
$$

$$
x_{it} \geq 0, \qquad t = 1, \ldots, T,\ i = 1, \ldots, N, \tag{1d}
$$

$$
z_t \leq 1, \qquad t = 1, \ldots, T, \tag{1e}
$$

$$
x_{it}, z_t \in \{0,1\}, \quad t = 1, \ldots, T,\ i = 1, \ldots, N \tag{1f}
$$

The basic replacement problem, II

Objective (1a)

Minimize the total cost of having a working system during the contract period

Constraint (1b)

For any given item i in the system, the component must be replaced at some point during *every* time interval of T_i time steps

Constraint (1c)

No replacement can be performed at time t without paying the fixed cost d_t (for a maint. operation); once we pay, any maint. action becomes possible (at no extra fixed cost) at time t
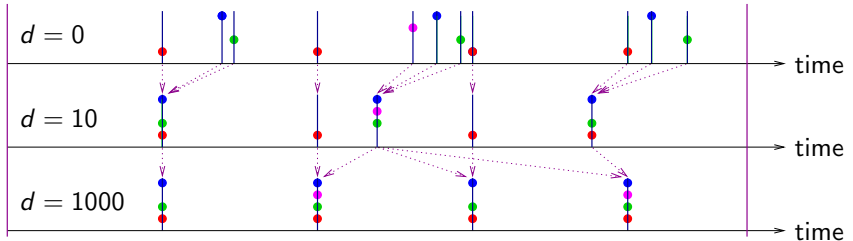
Constraints (1d)–(1f)

Ensure that the variables take only meaningful values

Opportunistic maintenance or not?

Example: four components with different prices and lives

- A *replacement* is marked with a dot; its colour represents the type of component replaced
- The larger the fixed cost, the more beneficial *opportunistic* maintenance becomes; also more items are replaced

Constraint structure—example

Time horizon: $T = 8$. Component #3: $T_3 = 4$

$$\sum_{t=\ell+1}^{\ell+T_3} x_{3t} \geq 1, \qquad \ell = 0, \ldots, T - T_3$$

$$\iff \sum_{t=\ell+1}^{\ell+4} x_{3t} \geq 1, \qquad \ell = 0, \ldots, 4$$

$$\begin{bmatrix} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} x_{31} \\ x_{32} \\ \vdots \\ x_{38} \end{bmatrix} \geq \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

Property I: the replacement problem is NP-hard

Theorem

Set covering is polynomially reducible to the replacement problem

- This essentially means that we *cannot* expect to find an optimal solution in a time that is proportional to a polynomial function of the problem size (i.e., $T(N+1)$ variables and $\approx 4NT$ constraints)

- Basic complexity theory: Chapter 2.6 in the course book

Property II: with fixed z the problem over x is easy

- The constraint matrix has the "consecutive ones" property
- $\Rightarrow$ For fixed values of z, the problem over x can be solved as a linear program
- For each i, the linear programming dual problem can be solved by a "greedy" algorithm $\Rightarrow$ primal solution by complementarity; see [a], Algorithm 1, page 297
- The latter is typically 5–40 times faster than solving as a general linear program, and 25–400 times faster when costs are monotone with time (i.e., $\forall t$ either $c_{it} \leq c_{i,t+1}$ or $c_{it} \geq c_{i,t+1}$); see [a], Algorithm 2, page 299
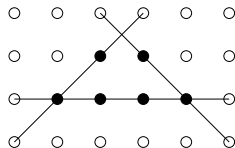
[a] T. Almgren, N. Andréasson, M. Patriksson, A.-B. Strömberg, A. Wojciechowski, M. Önnheim (2012): *The opportunistic replacement problem: theoretical analyses and numerical tests*, Mathematical Methods of Operations Research, 76(3) pp. 289–319.

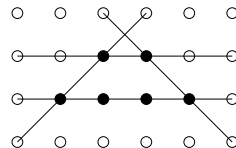# Property III: all inequalities are facet defining[*]

No inequalities are
facet defining
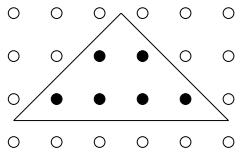
All inequalities are
facet defining

An integral
polyhederon

[*] See [a], Section 5.1–5.2

A generalized model

New variable definition

Define the set

$$\mathcal{I} := \{ (s, t) \mid 0 \leq s < t \leq T + 1; \, s, t \in Z \}$$

of *replacement intervals* and introduce the variables

$$x_{st}^i = \begin{cases} 1, & \text{if component } i \text{ receives PM at the} \\ & \text{times } s \text{ and } t, \text{ and not in-between,} \\ 0, & \text{otherwise,} \end{cases} \quad \begin{array}{l} i \in \mathcal{N}, \\ (s, t) \in \mathcal{I}, \end{array}$$

and

$$z_t = \begin{cases} 1, & \text{if maintenance occurs at time } t, \\ 0, & \text{otherwise,} \end{cases} \quad t \in \mathcal{T}.$$

A generalized model

$$
\begin{align}
\text{minimize} \quad & \sum_{t \in \mathcal{T}} d_t z_t + \sum_{i \in \mathcal{N}} \sum_{(s,t) \in \mathcal{I}} c_{st}^i x_{st}^i, && \text{(2a)} \\
\text{subject to} \quad & \sum_{s=0}^{t-1} x_{st}^i \leq z_t, && i \in \mathcal{N}, t \in \mathcal{T}, && \text{(2b)} \\
& \sum_{s=0}^{t-1} x_{st}^i = \sum_{r=t+1}^{T+1} x_{tr}^i, && i \in \mathcal{N}, t \in \mathcal{T}, && \text{(2c)} \\
& \sum_{t=1}^{T+1} x_{0t}^i = 1, && i \in \mathcal{N}, && \text{(2d)} \\
& x_{st}^i \in \{0,1\}, && i \in \mathcal{N}, (s,t) \in \mathcal{I}, && \text{(2e)} \\
& z_t \in \{0,1\}, && t \in \mathcal{T}. && \text{(2f)}
\end{align}
$$

On the GKN project

Aircraft engines are expensive

- Spare components cost up to 2 MSEK
- Total cost of maintenance of one engine: 15–30 MSEK
- Maximizing "time on wing" is important, both for civil and military aircraft

- The aircraft engine RM12 consists of 7 modules and 61 components in total
- A mathematical model has been constructed for the entire engine maintenance, including work costs for (dis)assembling the necessary modules and components for each maintenance occasion
- This model has slightly less than 6000 binary variables

Results on the GKN problems

An individual engine module with 10 components

- Cost reduction: 35%
- Reduction of # maintenance occasions: 7%

as compared with a simple policy similar to that used at GKN

A complete engine with 7 modules (61 components)

- Cost reduction compared to maintaining (optimally) each individual module: 12%
- Reduction of # maintenance occasions: 60%

Product development

Found 5 components which can potentially reduce maintenance costs more than 5% through prolonged lives

Assignment 2: Maintenance Scheduling

Assignment tasks in summary

- Study and compare the two ILP models for maintenance planning
- Elaborate with the models - wrt facets, integrality property, time horizon
- Heuristic solutions—local search method
- Add side constraints—modelling additional properties
- Students aiming at grade 4, 5, or VG must answer ALL the questions
- Deadline for handing in report: May 6
- Written opposition on another group's report.
  Deadline: May 11