

MVE165/MMG631  
Linear and integer optimization with applications  
Lecture 14  
Overview of nonlinear programming  
and summary of the course

Ann-Brith Strömberg

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### Structural optimization

- Design of aircraft, ships, bridges, etc
- Decide on the material and the topology and thickness of a mechanical structure
- Minimize weight, maximize stiffness, constraints on deformation at certain loads, strength, fatigue limit, etc

### Analysis and design of traffic networks

- Estimate traffic flows and discharges
- Detect bottlenecks
- Analyze effects of traffic signals, tolls, etc

## Least squares

Adaptation of data

## Engine development, design of antennas or tyres, etc.

For each function evaluation a computationally expensive (time consuming) simulation may be needed

## Maximize the volume of a cylinder

While keeping the surface area constant

## Wind power generation

The energy content in the wind is  $\propto v^3$  (in Ass 3a it is discretized and measured data is used)

# An overview of nonlinear optimization

## General notation for nonlinear programs

$$\begin{array}{ll} \text{minimize } \mathbf{x} \in \mathbb{R}^n & f(\mathbf{x}) \\ \text{subject to} & g_i(\mathbf{x}) \leq 0, \quad i \in \mathcal{L}, \\ & h_i(\mathbf{x}) = 0, \quad i \in \mathcal{E}. \end{array}$$

## Some special cases

- Unconstrained problems ( $\mathcal{L} = \mathcal{E} = \emptyset$ ):

$$\boxed{\text{minimize } f(\mathbf{x}) \text{ subject to } \mathbf{x} \in \mathbb{R}^n}$$

- Convex programming:  $f$  convex,  $g_i$  convex,  $i \in \mathcal{L}$ ,  $h_i$  linear,  $i \in \mathcal{E}$ .
- Linear constraints:  $g_i$ ,  $i \in \mathcal{L}$ , and  $h_i$ ,  $i \in \mathcal{E}$

- Quadratic programming:

$$\boxed{f(\mathbf{x}) = \mathbf{c}^T \mathbf{x} + \frac{1}{2} \mathbf{x}^T \mathbf{Q} \mathbf{x}}$$

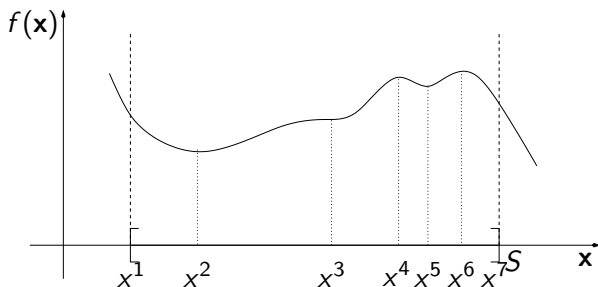
- Linear programming:

$$\boxed{f(\mathbf{x}) = \mathbf{c}^T \mathbf{x}}$$

# Properties of nonlinear programs

- The mathematical properties of nonlinear optimization problems can be very different
- *No* algorithm exists that solves *all* nonlinear optimization problems
- An optimal solution does *not* have to be located at an extreme point
- Nonlinear programs can be unconstrained
  - What if a *linear program* has no constraints?
- $f$  may be differentiable or non-differentiable  
E.g., the Lagrangean dual objective function
- *For convex problems*: Algorithms (typically) converge to an optimal solution
- Nonlinear problems can have *local* optima that are *not global* optima

Consider the problem to minimize  $f(\mathbf{x})$  subject to  $\mathbf{x} \in S$



Possible extremal points are

- boundary points of  $S = [x^1, x^7]$  (i.e.,  $\{x^1, x^7\}$ )
- stationary points, where  $f'(\mathbf{x}) = 0$  (i.e.,  $\{x^2, \dots, x^6\}$ )
- discontinuities in  $f$  or  $f'$

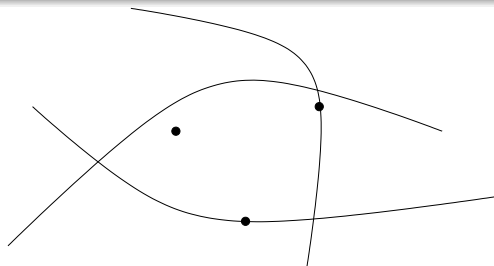
DRAW!

## Boundary points

$\bar{\mathbf{x}}$  is a *boundary* point to the feasible set

$$S = \{\mathbf{x} \in \mathbb{R}^n \mid g_i(\mathbf{x}) \leq 0, i \in \mathcal{L}\}$$

if  $g_i(\bar{\mathbf{x}}) \leq 0, i \in \mathcal{L}$ , and  $g_i(\bar{\mathbf{x}}) = 0$  for at least one index  $i \in \mathcal{L}$



## Stationary points

$\bar{\mathbf{x}}$  is a *stationary* point to  $f$  if  $\nabla f(\bar{\mathbf{x}}) = \mathbf{0}^n$  (for  $n = 1$ : if  $f'(\bar{\mathbf{x}}) = 0$ )

Consider the nonlinear optimization problem to

minimize  $f(\mathbf{x})$  subject to  $\mathbf{x} \in S$

### Local minimum

- *In words:* A solution is a *local* minimum if it is *feasible* and no other feasible solution in a sufficiently *small neighbourhood* of the solution at hand has a lower objective value
- *Formally:*  $\bar{\mathbf{x}}$  is a local minimum if  $\bar{\mathbf{x}} \in S$  and  $\exists \epsilon > 0$  such that  $f(\bar{\mathbf{x}}) \leq f(\mathbf{x})$  for all  $\mathbf{x} \in \{\mathbf{y} \in S : \|\mathbf{y} - \bar{\mathbf{x}}\| \leq \epsilon\}$  DRAW!!

### Global minimum

- *In words:* A solution is a *global* minimum if it is *feasible* and no other feasible solution has a lower objective value
- *Formally:*  $\bar{\mathbf{x}}$  is a global minimum if  $\bar{\mathbf{x}} \in S$  and  $f(\bar{\mathbf{x}}) \leq f(\mathbf{x})$  for all  $\mathbf{x} \in S$



# When is a local optimum also a global optimum? (Ch. 9.3)

The concept of convexity is essential

- Functions: convex (minimization), concave (maximization)
- Sets: convex (minimization and maximization)
- The minimization (maximization) of a convex (concave) function over a convex set is referred to as a convex optimization problem

Definition 9.5: Convex optimization problem

If  $f$  and  $g_i$ ,  $i \in \mathcal{L}$ , are convex functions, then

$$\text{minimize } f(\mathbf{x}) \text{ subject to } g_i(\mathbf{x}) \leq 0, i \in \mathcal{L}$$

is said to be a *convex* optimization problem

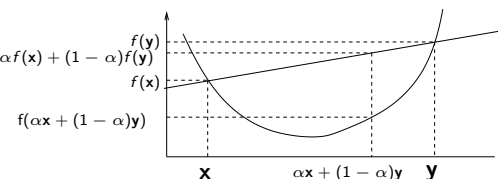
Theorem 9.1: Global optimum

Let  $\mathbf{x}^*$  be a *local* optimum of a *convex* optimization problem. Then  $\mathbf{x}^*$  is also a *global* optimum

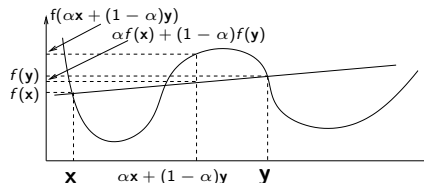
# Convex functions

A function  $f$  is *convex* on  $S$  if, for any  $\mathbf{x}, \mathbf{y} \in S$  it holds that  $f(\alpha\mathbf{x} + (1 - \alpha)\mathbf{y}) \leq \alpha f(\mathbf{x}) + (1 - \alpha)f(\mathbf{y})$  for all  $0 \leq \alpha \leq 1$

A CONVEX FUNCTION



A NON-CONVEX FUNCTION



The function  $f$  is *strictly convex* on  $S$  if, for any  $\mathbf{x}, \mathbf{y} \in S$  such that  $\mathbf{x} \neq \mathbf{y}$  it holds that

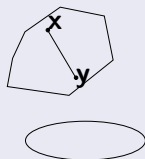
$$f(\alpha\mathbf{x} + [1 - \alpha]\mathbf{y}) < \alpha f(\mathbf{x}) + (1 - \alpha)f(\mathbf{y}) \text{ for all } 0 < \alpha < 1$$

# Convex sets

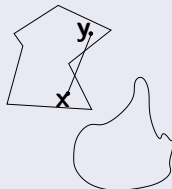
A set  $S$  is convex if, for any  $\mathbf{x}, \mathbf{y} \in S$  it holds that  $\alpha\mathbf{x} + (1 - \alpha)\mathbf{y} \in S$  for all  $0 \leq \alpha \leq 1$

## Examples

Convex sets



Non-convex sets



Consider a set  $S$  defined by the intersection of  $m = |\mathcal{L}|$  inequalities, where the functions  $g_i : \mathbb{R}^n \mapsto \mathbb{R}$ ,  $i \in \mathcal{L}$ , as

$$S = \{ \mathbf{x} \in \mathbb{R}^n \mid g_i(\mathbf{x}) \leq 0, i \in \mathcal{L} \}$$

## Theorems 9.2 & 9.3

If all the functions  $g_i$ ,  $i \in \mathcal{L}$ , are convex on  $\mathbb{R}^n$ , then the set  $S$  is convex

# The Karush-Kuhn-Tucker conditions: necessary conditions for optimality

Let  $S := \{ \mathbf{x} \in \mathbb{R}^n \mid g_i(\mathbf{x}) \leq 0, i \in \mathcal{L} \}$

- Assume that the following hold
  - the function  $f : \mathbb{R}^n \mapsto \mathbb{R}$  is differentiable;
  - the functions  $g_i : \mathbb{R}^n \mapsto \mathbb{R}, i \in \mathcal{L}$ , are convex and differentiable;
  - there exists a point  $\bar{\mathbf{x}} \in S$  such that  $g_i(\bar{\mathbf{x}}) < 0, i \in \mathcal{L}$
- If  $\mathbf{x}^* \in S$  is a local minimum of  $f$  over  $S$ , then there exists a vector  $\boldsymbol{\mu} \in \mathbb{R}^m$  (where  $m = |\mathcal{L}|$ ) such that

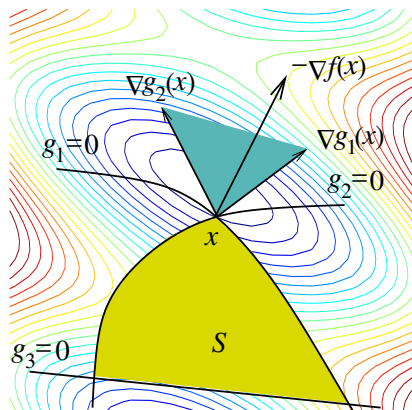
$$\nabla f(\mathbf{x}^*) + \sum_{i \in \mathcal{L}} \mu_i \nabla g_i(\mathbf{x}^*) = \mathbf{0}^n,$$

$$\mu_i g_i(\mathbf{x}^*) = 0, \quad i \in \mathcal{L},$$

$$\boldsymbol{\mu} \geq \mathbf{0}^m.$$

□

# Geometry of the Karush-Kuhn-Tucker conditions



**Figur:** Geometric interpretation of the Karush-Kuhn-Tucker conditions. At a local minimum, the negative gradient of the objective function can be expressed as a non-negative linear combination of the gradients of the active constraints at this point

# The Karush-Kuhn-Tucker conditions: sufficient for optimality under convexity

Assume that the functions  $f, g_i : \mathbb{R}^n \mapsto \mathbb{R}, i \in \mathcal{L}$ , are convex and differentiable, and let  $S = \{ \mathbf{x} \in \mathbb{R}^n \mid g_i(\mathbf{x}) \leq 0, i \in \mathcal{L} \}$

If the conditions (where  $m = |\mathcal{L}|$ )

$$\nabla f(\mathbf{x}^*) + \sum_{i \in \mathcal{L}} \mu_i \nabla g_i(\mathbf{x}^*) = \mathbf{0}^n,$$

$$\mu_i g_i(\mathbf{x}^*) = 0, \quad i \in \mathcal{L},$$

$$\boldsymbol{\mu} \geq \mathbf{0}^m$$

hold, then  $\mathbf{x}^* \in S$  is a global minimum of  $f$  over  $S$ . □

- The Karush-Kuhn-Tucker conditions can also be stated for optimization problems with equality constraints
- For unconstrained optimization KKT reads:  $\nabla f(\mathbf{x}^*) = \mathbf{0}$
- For a quadratic program KKT forms a system of linear (in)equalities plus the complementarity constraints

## The optimality conditions can be used to..

- verify an (local) optimal solution
- solve certain special cases of nonlinear programs (e.g. quadratic programs)
- algorithm construction
- derive properties of a solution to a non-linear program

## Example

$$\begin{aligned} \text{minimize} \quad & f(\mathbf{x}) := 2x_1^2 + 2x_1x_2 + x_2^2 - 10x_1 - 10x_2 \\ \text{subject to} \quad & x_1^2 + x_2^2 \leq 5 \\ & 3x_1 + x_2 \leq 6 \end{aligned}$$

Is  $\mathbf{x}^0 = (1, 2)^T$  a Karush-Kuhn-Tucker point?

- Is it an optimal solution?
- Derive:  $\nabla f(\mathbf{x}) = (4x_1 + 2x_2 - 10, 2x_1 + 2x_2 - 10)^T$ ,  
 $\nabla g_1(\mathbf{x}) = (2x_1, 2x_2)^T$ , and  $\nabla g_2(\mathbf{x}) = (3, 1)^T$

$$\begin{aligned} 4x_1^0 + 2x_2^0 - 10 + 2x_1^0\mu_1 + 3\mu_2 &= 0 \\ 2x_1^0 + 2x_2^0 - 10 + 2x_2^0\mu_1 + \mu_2 &= 0 \\ \mu_1[(x_1^0)^2 + (x_2^0)^2 - 5] + \mu_2(3x_1^0 + x_2^0 - 6) &= 0 \\ \mu_1, \mu_2 &\geq 0 \end{aligned}$$

$\iff$

$$\begin{aligned} 2\mu_1 + 3\mu_2 &= 2 \\ 4\mu_1 + \mu_2 &= 4 \\ 0\mu_1 - \mu_2 &= 0 \\ \mu_1, \mu_2 &\geq 0 \end{aligned}$$

$$\Rightarrow \mu_2 = 0 \quad \Rightarrow \quad \mu_1 = 1 \geq 0$$



## Example, continued

OK, the Karush-Kuhn-Tucker conditions hold

Is the solution optimal? Check convexity!

- $\nabla^2 f(\mathbf{x}) = \begin{pmatrix} 4 & 2 \\ 2 & 2 \end{pmatrix}$ ,  $\nabla^2 g_1(\mathbf{x}) = \begin{pmatrix} 2 & 0 \\ 0 & 2 \end{pmatrix}$ ,  $\nabla^2 g_2(\mathbf{x}) = \mathbf{0}^{2 \times 2}$

$\Rightarrow f$ ,  $g_1$ , and  $g_2$  are convex

$\Rightarrow \mathbf{x}^0 = (1, 2)^T$  is an optimal solution and  $f(\mathbf{x}^0) = -20$

- 1 Choose a starting solution,  $\mathbf{x}^0 \in \mathfrak{R}^n$ . Let  $k = 0$
- 2 Determine a **search direction**  $\mathbf{d}^k$
- 3 If a termination criterion is fulfilled  $\Rightarrow$  Stop!
- 4 Determine a step length,  $t_k$ , by solving:

$$\text{minimize}_{t \geq 0} \varphi(t) := f(\mathbf{x}^k + t \cdot \mathbf{d}^k)$$

- 5 New iteration point,  $\mathbf{x}^{k+1} = \mathbf{x}^k + t_k \cdot \mathbf{d}^k$
- 6 Let  $k := k + 1$  and return to step 2

How choose **search directions**  $\mathbf{d}^k$ , **step lengths**  $t_k$ , and **termination criteria**?

Goal:  $f(\mathbf{x}^{k+1}) < f(\mathbf{x}^k)$  (minimization)

- How does  $f$  change locally in a direction  $\mathbf{d}^k$  at  $\mathbf{x}^k$ ?
- Taylor expansion (Ch. 9.2):  
$$f(\mathbf{x}^k + t\mathbf{d}^k) = f(\mathbf{x}^k) + t\nabla f(\mathbf{x}^k)^\top \mathbf{d}^k + \mathcal{O}(t^2)$$
- For sufficiently small  $t > 0$ :  
$$f(\mathbf{x}^k + t\mathbf{d}^k) < f(\mathbf{x}^k) \quad \Rightarrow \quad \nabla f(\mathbf{x}^k)^\top \mathbf{d}^k < 0$$

$\Rightarrow$

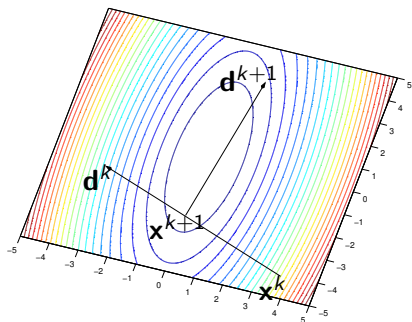
### Definition

If  $\nabla f(\mathbf{x}^k)^\top \mathbf{d}^k < 0$  then  $\mathbf{d}^k$  is a descent direction for  $f$  at  $\mathbf{x}^k$   
If  $\nabla f(\mathbf{x}^k)^\top \mathbf{d}^k > 0$  then  $\mathbf{d}^k$  is an ascent direction for  $f$  at  $\mathbf{x}^k$

We wish to minimize (maximize)  $f$  over  $\mathbb{R}^n$

$\Rightarrow$  Choose  $\mathbf{d}^k$  as a descent (an ascent) direction from  $\mathbf{x}^k$

# An improving step



**Figur:** At  $\mathbf{x}^k$ , the descent direction  $\mathbf{d}^k$  is generated. A step  $t_k$  is taken in this direction, producing  $\mathbf{x}^{k+1}$ . At this point, a new descent direction  $\mathbf{d}^{k+1}$  is generated, etc

- 1 Choose a starting solution,  $\mathbf{x}^0 \in \mathbb{R}^n$ . Let  $k = 0$
- 2 Determine a search direction  $\mathbf{d}^k$
- 3 If a termination criterion is fulfilled  $\Rightarrow$  Stop!
- 4 Determine a step length,  $t_k$ , by solving:

$$\text{minimize}_{t \geq 0} \varphi(t) := f(\mathbf{x}^k + t \cdot \mathbf{d}^k)$$

- 5 New iteration point,  $\mathbf{x}^{k+1} = \mathbf{x}^k + t_k \cdot \mathbf{d}^k$
- 6 Let  $k := k + 1$  and return to step 2

- Solve  $\min_{t \geq 0} \varphi(t) := f(\mathbf{x}^k + t \cdot \mathbf{d}^k)$  where  $\mathbf{d}^k$  is a descent direction from  $\mathbf{x}^k$
- A minimization problem in one variable  $\Rightarrow$  Solution  $t_k$
- Analytic solution:  $\varphi'(t_k) = 0$  (seldom possible to derive)

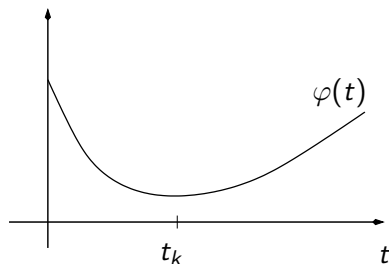
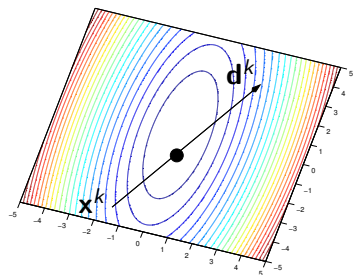
### Numerical solution methods

- The golden section method (reduce the interval of uncertainty)
- The bi-section method (reduce the interval of uncertainty)
- Newton-Raphson's method
- Armijo's method

### In practice

Do not solve exactly, but to a sufficient improvement of the function value:  $f(\mathbf{x}^k + t_k \mathbf{d}^k) \leq f(\mathbf{x}^k) - \varepsilon$  for some  $\varepsilon > 0$

# Line search



**Figur:** A line search in a descent direction.  
 $t_k$  solves  $\min_{t \geq 0} \varphi(t) := f(\mathbf{x}^k + t \cdot \mathbf{d}^k)$

# General iterative search method for unconstrained optimization

- 1 Choose a starting solution,  $\mathbf{x}^0 \in \mathbb{R}^n$ . Let  $k = 0$
- 2 Determine a search direction  $\mathbf{d}^k$
- 3 If a **termination criterion** is fulfilled  $\Rightarrow$  Stop!
- 4 Determine a step length,  $t_k$ , by solving:

$$\text{minimize}_{t \geq 0} \varphi(t) := f(\mathbf{x}^k + t \cdot \mathbf{d}^k)$$

- 5 New iteration point,  $\mathbf{x}^{k+1} = \mathbf{x}^k + t_k \cdot \mathbf{d}^k$
- 6 Let  $k := k + 1$  and return to step 2



## Termination criteria

Needed since  $\nabla f(\mathbf{x}^k) = \mathbf{0}$  will never be fulfilled exactly

Typical choices ( $\varepsilon_j > 0, j = 1, \dots, 4$ )

- (a)  $\|\nabla f(\mathbf{x}^k)\| < \varepsilon_1$
- (b)  $|f(\mathbf{x}^{k+1}) - f(\mathbf{x}^k)| < \varepsilon_2$
- (c)  $\|\mathbf{x}^{k+1} - \mathbf{x}^k\| < \varepsilon_3$
- (d)  $t_k < \varepsilon_4$

The criteria (a)–(d) are often combined

The search method only guarantees a stationary solution, whose properties are determined by the properties of  $f$  (convexity, ...)

# Constrained optimization: Penalty methods

Consider both inequality and equality constraints

$$\begin{aligned} & \text{minimize}_{\mathbf{x} \in \mathbb{R}^n} && f(\mathbf{x}) \\ & \text{subject to} && g_i(\mathbf{x}) \leq 0, \quad i \in \mathcal{L}, \\ & && h_i(\mathbf{x}) = 0, \quad i \in \mathcal{E}. \end{aligned} \tag{1}$$

Drop the constraints and add terms in the objective that *penalize infeasible solutions*

$$\text{minimize}_{\mathbf{x} \in \mathbb{R}^n} F_\mu(\mathbf{x}) := f(\mathbf{x}) + \mu \sum_{i \in \mathcal{L} \cup \mathcal{E}} \alpha_i(\mathbf{x}) \tag{2}$$

where  $\mu > 0$  and  $\alpha_i(\mathbf{x}) = \begin{cases} = 0 & \text{if } \mathbf{x} \text{ satisfies constraint } i \\ > 0 & \text{otherwise} \end{cases}$

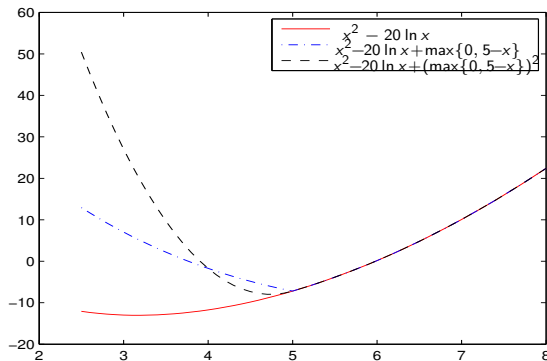
Common penalty functions (which of these are differentiable?)

$$i \in \mathcal{L}: \quad \alpha_i(\mathbf{x}) = \max\{0, g_i(\mathbf{x})\} \quad \text{or} \quad \alpha_i(\mathbf{x}) = (\max\{0, g_i(\mathbf{x})\})^2$$

$$i \in \mathcal{E}: \quad \alpha_i(\mathbf{x}) = |h_i(\mathbf{x})| \quad \text{or} \quad \alpha_i(\mathbf{x}) = |h_i(\mathbf{x})|^2$$

# Squared and non-squared penalty functions

minimize  $(x^2 - 20 \ln x)$  subject to  $x \geq 5$



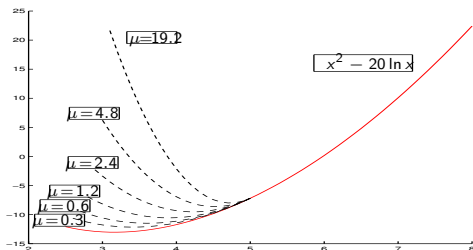
Figur: Squared and non-squared penalty function.  $g_i$  differentiable  $\implies$  squared penalty function differentiable

# Squared penalty functions

- In practice: Start with a low value of  $\mu > 0$  and increase the value as the computations proceed

- **Example:**  $\text{minimize } (x^2 - 20 \ln x) \text{ subject to } x \geq 5$  (\*)

$\Rightarrow$   $\text{minimize } (x^2 - 20 \ln x + \mu(\max\{0, 5 - x\})^2)$  (\*\*)



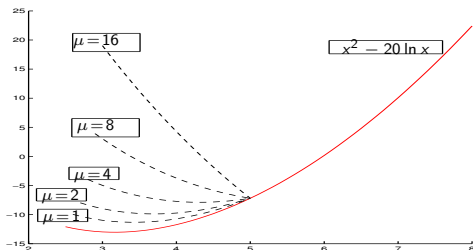
**Figur:** Squared penalty function:  $\exists \mu < \infty$  such that an optimal solution for (\*\*) is optimal (feasible) for (\*)

# Non-squared penalty functions

- In practice: Start with a low value of  $\mu > 0$  and increase the value as the computations proceed

- **Example:**  $\text{minimize } (x^2 - 20 \ln x) \text{ subject to } x \geq 5$  (+)

$\Rightarrow \text{minimize } (x^2 - 20 \ln x + \mu \max\{0, 5 - x\})$  (++)



**Figur:** Non-squared penalty function: For  $\mu \geq 6$  the optimal solution for (++) is optimal (and feasible) for (+)

# Constrained optimization: Barrier methods

Consider only inequality constraints

$$\begin{aligned} & \text{minimize}_{\mathbf{x} \in \mathbb{R}^n} && f(\mathbf{x}) \\ & \text{subject to} && g_i(\mathbf{x}) \leq 0, \quad i \in \mathcal{L} \end{aligned} \quad (3)$$

- Drop the constraints and add terms in the objective that *prevents from approaching the boundary* of the feasible set

$$\text{minimize}_{\mathbf{x} \in \mathbb{R}^n} F_\mu(\mathbf{x}) := f(\mathbf{x}) + \mu \sum_{i \in \mathcal{L}} \alpha_i(\mathbf{x}) \quad (4)$$

where  $\mu > 0$  and  $\alpha_i(\mathbf{x}) \rightarrow +\infty$  as  $g_i(\mathbf{x}) \rightarrow 0$  (as constraint  $i$  approaches being active)

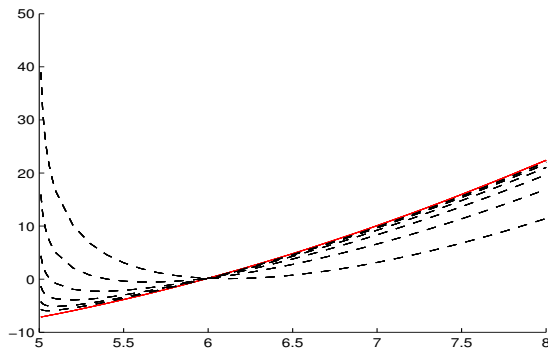
Common barrier functions

$$\alpha_i(\mathbf{x}) = -\ln[-g_i(\mathbf{x})] \quad \text{or} \quad \alpha_i(\mathbf{x}) = \frac{-1}{g_i(\mathbf{x})}$$

# Logarithmic barrier functions

- Choose  $\mu > 0$  and decrease it as the computations proceed
- **Example:** minimize  $(x^2 - 20 \ln x)$  subject to  $x \geq 5$

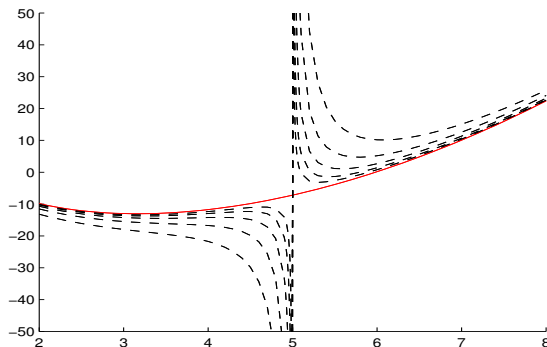
$\Rightarrow$  minimize  $_{x>5} (x^2 - 20 \ln x - \mu \ln(x - 5))$



# Fractional barrier functions

- Choose  $\mu > 0$  and decrease it as the computations proceed
- **Example:** minimize  $(x^2 - 20 \ln x)$  subject to  $x \geq 5$

$\Rightarrow$  minimize  $x > 5 \left( x^2 - 20 \ln x + \frac{\mu}{x-5} \right)$





## Summary of the theoretical content of the course ...

... which may appear at the oral exam:

- Mathematical modelling of optimization problems; graphic solution
- *Linear programming*: BFSs; the simplex method; degeneracy; multiple optima; unbounded solution; infeasibility; starting solutions; LP duality; post-optimal and sensitivity analysis
- *Discrete and combinatorial optimization*: models of specific ILP problems; mathematical properties; complexity; algorithms; local/global optima; neighbourhoods; heuristics
- *Network flows*: Shortest paths; dynamic programming; LP models of network flows; maximum flows; minimum cost network flows; unimodularity; integrality property
- *Multi-objective optimization*: Pareto optimality; (non-)convexity; solution methods
- *Non-linear optimization*: convexity; local/global optimality; mathematical properties; solution methods