

MVE165/MMG631  
Linear and integer optimization with applications  
Lecture 13  
Multiobjective optimization

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## Applied optimization — multiple objectives

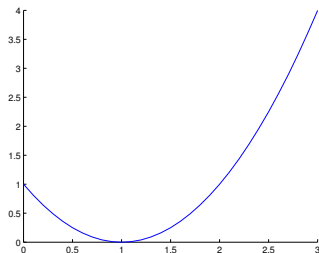
- Many practical optimization problems have several objectives which may be in conflict
- Some goals cannot be reduced to a common scale of cost/profit  $\Rightarrow$  trade-offs must be addressed
- Examples
  - Financial investments — risk vs. return
  - Engine design — efficiency vs.  $\text{NO}_x$  vs. soot
  - Wind power production — investment vs. operation (Ass 3a)
  - Electricity generation — costs vs. emissions (Ass 3b)

### Literature on multiple objectives' optimization

Copies from the book *Optimization in Operations Research* by R.L. Rardin (1998) pp. 373–387, handed out (on paper, copies kept outside Ann-Brith's office, room MV:L2087)

# Optimization of multiple objectives

- Consider the minimization of  $f(x) := (x - 1)^2$  subject to  $0 \leq x \leq 3$
- Optimal solution:  $x^* = 1$   
(since the function  $f$  is convex)



# Optimization of multiple objectives

Consider then two objectives

$$\text{minimize } [f_1(x); f_2(x)]$$

$$\text{subject to } 0 \leq x \leq 3$$

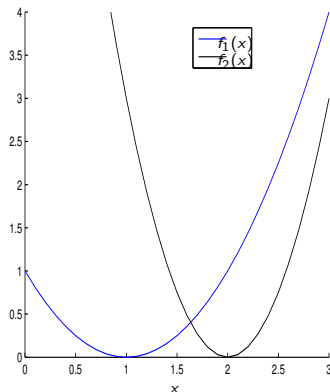
where

$$f_1(x) := (x - 1)^2, \quad f_2(x) := 3(x - 2)^2$$

- How can an optimal solution be defined?

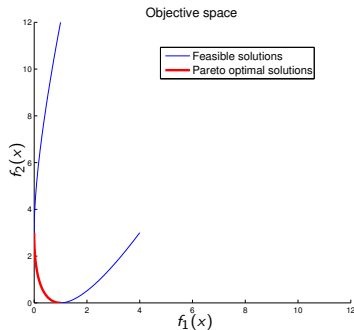
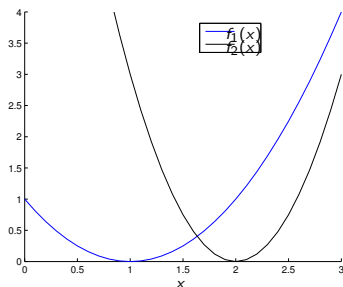
A solution is *Pareto optimal* if *no other* feasible solution has a better value in *all* objectives

- All points  $x \in [1, 2]$  are Pareto optimal



# Pareto optimal solutions in the objective space

- minimize  $[f_1(x); f_2(x)]$  subject to  $0 \leq x \leq 3$   
where  $f_1(x) := (x - 1)^2$  and  $f_2(x) := 3(x - 2)^2$
- A solution is *Pareto optimal* if *no other* feasible solution has a better value in *all* objectives

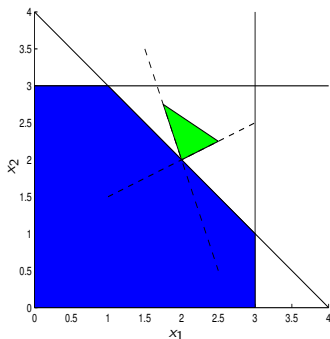


- Pareto optima  $\Leftrightarrow$  nondominated points  $\Leftrightarrow$  efficient frontier

# Efficient points

- Consider a bi-objective linear program:

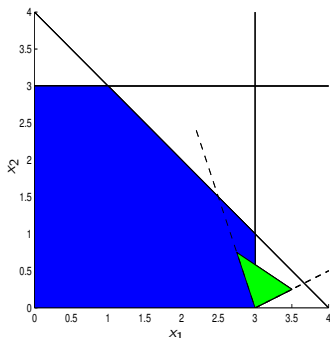
$$\begin{array}{ll} \text{maximize} & 3x_1 + x_2 \\ \text{maximize} & -x_1 + 2x_2 \\ \text{subject to} & x_1 + x_2 \leq 4 \\ & 0 \leq x_1 \leq 3 \\ & 0 \leq x_2 \leq 3 \end{array}$$



- The solutions in the green cone, defined by  $\left\{ \mathbf{x} \in \mathbb{R}^2 \mid \mathbf{x} = \begin{pmatrix} 2 \\ 2 \end{pmatrix} + \lambda_1 \begin{pmatrix} -1 \\ 3 \end{pmatrix} + \lambda_2 \begin{pmatrix} 2 \\ 1 \end{pmatrix}; \lambda_1, \lambda_2 > 0 \right\}$ , are better than the solution  $(2, 2)$  w.r.t. both objectives
- The point  $\mathbf{x} = (2, 2)$  is an *efficient*, or *non-dominated*, solution

# Dominated points

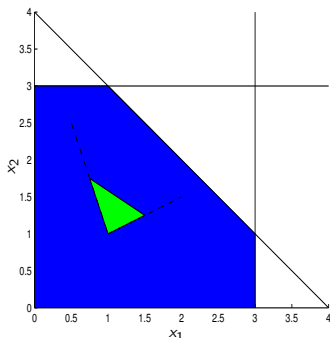
$$\begin{array}{ll} \text{maximize} & 3x_1 + x_2 \\ \text{maximize} & -x_1 + 2x_2 \\ \text{subject to} & x_1 + x_2 \leq 4 \\ & 0 \leq x_1 \leq 3 \\ & 0 \leq x_2 \leq 3 \end{array}$$



- The point  $x = (3, 0)$  is *dominated* by the solutions in the green cone
- Feasible solutions exist that are better w.r.t. both objectives

# Dominated points

$$\begin{array}{ll} \text{maximize} & 3x_1 + x_2 \\ \text{maximize} & -x_1 + 2x_2 \\ \text{subject to} & x_1 + x_2 \leq 4 \\ & 0 \leq x_1 \leq 3 \\ & 0 \leq x_2 \leq 3 \end{array}$$

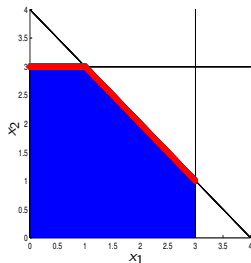


- The point  $x = (1, 1)$  is dominated by the solutions in the green cone
- Feasible solutions exist that are better w.r.t. both objectives



# The efficient frontier—the set of Pareto optimal solutions

$$\begin{array}{ll} \text{maximize} & 3x_1 + x_2 \\ \text{maximize} & -x_1 + 2x_2 \\ \text{subject to} & x_1 + x_2 \leq 4 \\ & 0 \leq x_1 \leq 3 \\ & 0 \leq x_2 \leq 3 \end{array}$$



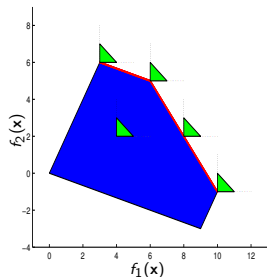
- The set of efficient solutions is given by

$$\left\{ \mathbf{x} \in \mathbb{R}^2 \mid \mathbf{x} = \alpha \begin{pmatrix} 3 \\ 1 \end{pmatrix} + (1 - \alpha) \begin{pmatrix} 1 \\ 3 \end{pmatrix}, 0 \leq \alpha \leq 1 \right\} \cup \left\{ \mathbf{x} \in \mathbb{R}^2 \mid \mathbf{x} = \alpha \begin{pmatrix} 1 \\ 3 \end{pmatrix} + (1 - \alpha) \begin{pmatrix} 0 \\ 3 \end{pmatrix}, 0 \leq \alpha \leq 1 \right\}$$

Note that this is *not* a convex set!

# The Pareto optimal set in the objective space

$$\begin{array}{ll} \text{maximize} & f_1(\mathbf{x}) := 3x_1 + x_2 \\ \text{maximize} & f_2(\mathbf{x}) := -x_1 + 2x_2 \\ \text{subject to} & x_1 + x_2 \leq 4 \\ & 0 \leq x_1 \leq 3 \\ & 0 \leq x_2 \leq 3 \end{array}$$



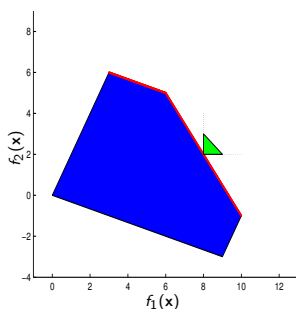
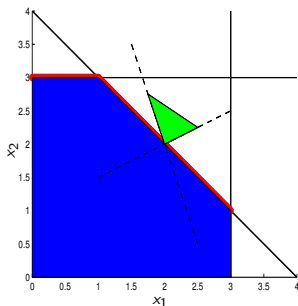
- The set of Pareto optimal objective values is given by

$$\left\{ (f_1, f_2) \in \mathbb{R}^2 \mid \mathbf{f} = \alpha \begin{pmatrix} 10 \\ -1 \end{pmatrix} + (1 - \alpha) \begin{pmatrix} 6 \\ 5 \end{pmatrix}, 0 \leq \alpha \leq 1 \right\} \cup \left\{ (f_1, f_2) \in \mathbb{R}^2 \mid \mathbf{f} = \alpha \begin{pmatrix} 6 \\ 5 \end{pmatrix} + (1 - \alpha) \begin{pmatrix} 3 \\ 6 \end{pmatrix}, 0 \leq \alpha \leq 1 \right\}$$

# Mapping from the decision space to the objective space

maximize  $[3x_1 + x_2; -x_1 + 2x_2]$

subject to  $x_1 + x_2 \leq 4, \quad 0 \leq x_1 \leq 3, \quad 0 \leq x_2 \leq 3$



# Solutions methods for multiobjective optimization

Construct the efficient frontier by treating one objective as a constraint and optimizing for the other

$$\begin{array}{ll} \text{maximize} & 3x_1 + x_2 \\ \text{subject to} & -x_1 + 2x_2 \geq \varepsilon \\ & x_1 + x_2 \leq 4 \\ & 0 \leq x_1 \leq 3 \\ & 0 \leq x_2 \leq 3 \end{array}$$

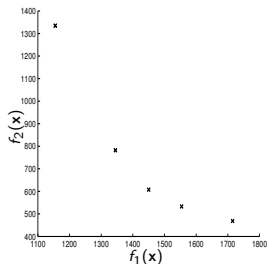
- Here, let  $\varepsilon \in [-1, 6]$ . Why?
- What if the number of objectives is  $\geq 3$ ?
- How many single objective linear programs do we have to solve for seven objectives and ten values of  $\varepsilon_k$  for each objective  $f_k$ ,  $k = 1, \dots, 7$ ?
- It is called the  $\varepsilon$ -constraints method

## Solution methods: weighted sums of objectives

- Give each maximization (minimization) objective a positive (negative) weight
- Solve a single objective maximization problem
- ⇒ Yields an efficient solution
- Drawback 1: Well spread weights do not necessarily produce solutions that are well spread on the efficient frontier

$$\text{Ex: } \left\{ \frac{1}{10}, \frac{1}{2}, 1, 2, 10 \right\}$$

- Drawback 2: If the objectives are *non-concave* (maximization) or if the feasible set is *non-convex*, as, e.g., integrality constrained, then *not all points on the efficient frontier may be possible to detect using weighted sums of objectives*



# The efficient frontier in the case of non-convexity

## A bi-objective binary linear program

$$\text{maximize } f_1(\mathbf{x}) := 3x_1 + x_2 - x_3$$

$$\text{maximize } f_2(\mathbf{x}) := x_1 - x_2 + 3x_3$$

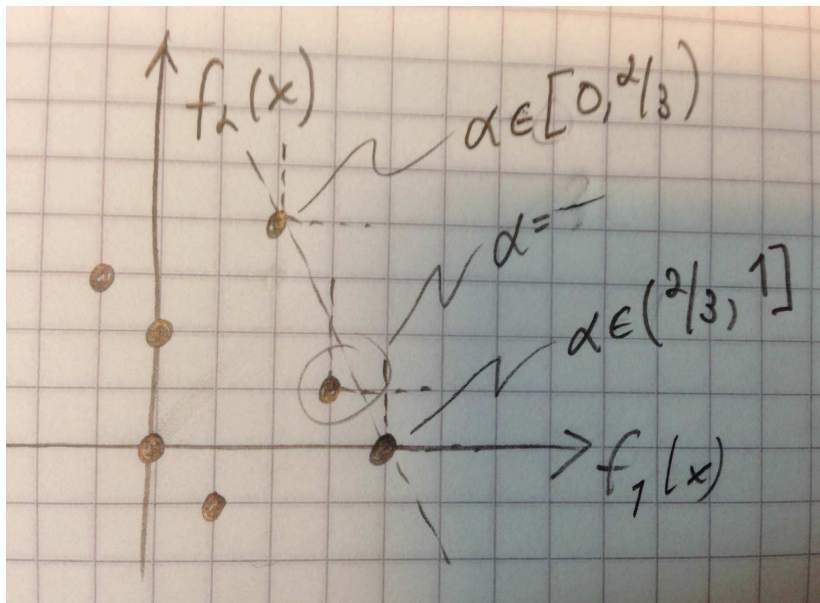
$$\text{subject to } \mathbf{x} \in X := \{ \mathbf{x} \in \mathbb{B}^3 \mid x_1 + x_2 + x_3 \leq 2 \}$$

Then,

$$X := \left\{ \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \\ 1 \end{pmatrix}, \begin{pmatrix} 1 \\ 0 \\ 1 \end{pmatrix}, \begin{pmatrix} 1 \\ 1 \\ 0 \end{pmatrix} \right\},$$

$$f_1(X) = \{0, -1, 1, 3, 0, 2, 4\} \quad \text{and} \quad f_2(X) = \{0, 3, -1, 1, 2, 4, 0\}$$

# The efficient frontier in the case of non-convexity



# The efficient frontier in the case of non-convexity

Solution by weighted maximization: Let  $\alpha \in [0, 1]$

$$\begin{aligned}\alpha f_1(\mathbf{x}) + (1 - \alpha)f_2(\mathbf{x}) &= \alpha(3x_1 + x_2 - x_3) + (1 - \alpha)(x_1 - x_2 + 3x_3) \\ &= (2\alpha + 1)x_1 + (2\alpha - 1)x_2 + (3 - 4\alpha)x_3\end{aligned}$$

Resulting binary linear program:

$$\begin{array}{ll}\text{maximize} & (2\alpha + 1)x_1 + (2\alpha - 1)x_2 + (3 - 4\alpha)x_3 \\ \text{subject to} & \mathbf{x} \in X\end{array}$$

- $\alpha \in [0, \frac{2}{3}) \implies \mathbf{x}^* = (1, 0, 1)^T$  &  $\mathbf{f}^* = (2, 4)^T$
  - $\alpha = \frac{2}{3} \implies \mathbf{x}^* \in \{(1, 0, 1)^T, (1, 1, 0)^T\}$  &  $\mathbf{f}^* \in \{(2, 4)^T, (4, 0)^T\}$
  - $\alpha \in (\frac{2}{3}, 1] \implies \mathbf{x}^* = (1, 1, 0)^T$  &  $\mathbf{f}^* = (4, 0)^T$
- 
- But the Pareto-optimal solution  $\mathbf{x}^* = (1, 0, 0)^T$  with  $\mathbf{f}^* = (3, 1)^T$  cannot be found using the weighted sums method
  - What would the  $\varepsilon$ -constraints method yield?



- Consider solving the previous example using the  $\varepsilon$ -constraint method

### The resulting one-objective binary linear program

$$\begin{array}{ll} \text{maximize}_{\mathbf{x}} & f_1(\mathbf{x}) := 3x_1 + x_2 - x_3 \\ \text{subject to} & f_2(\mathbf{x}) := x_1 - x_2 + 3x_3 \geq \varepsilon \\ & \mathbf{x} \in X := \{ \mathbf{x} \in \mathbb{B}^3 \mid x_1 + x_2 + x_3 \leq 2 \} \end{array}$$

- Then vary  $\varepsilon$  within relevant bounds (which are these?)

## Solution methods: soft constraints

Consider the multiobjective optimization problem to

$$\text{maximize}_{\mathbf{x}} [f_1(\mathbf{x}); \dots; f_K(\mathbf{x})] \quad \text{subject to } \mathbf{x} \in X$$

- Define a *target value*  $t_k$  and a *deficiency variable*  $d_k \geq 0$  for each objective  $f_k$
- Construct a *soft constraint* for each objective:

$$\text{maximize } f_k(\mathbf{x}) \quad \Rightarrow \quad f_k(\mathbf{x}) + d_k \geq t_k, \quad k = 1, \dots, K$$

Minimize the sum of deficiencies: (\*\*)

$$\begin{aligned} \text{minimize}_{\mathbf{x}, \mathbf{d}} \quad & \sum_{k \in K} d_k \\ \text{subject to} \quad & f_k(\mathbf{x}) + d_k \geq t_k, \quad k = 1, \dots, K \\ & d_k \geq 0, \quad k = 1, \dots, K \\ & \mathbf{x} \in X \end{aligned}$$

- When is an optimum of (\*\*) an efficient solution? [Draw!!]
- Important: Find first a common scale for  $f_k$ ,  $k = 1, \dots, K$

# Normalizing the objectives

Find a common scale for  $f_k$ ,  $k = 1, \dots, K$

- Consider the multiobjective optimization problem to

maximize $_{\mathbf{x}}$   $[f_1(\mathbf{x}); \dots; f_K(\mathbf{x})]$  subject to  $\mathbf{x} \in X$

- Define

$$\tilde{f}_k(\mathbf{x}) := \frac{f_k(\mathbf{x}) - f_k^{\min}}{f_k^{\max} - f_k^{\min}}, \quad k = 1, \dots, K,$$

where  $f_k^{\max} := \max_{\mathbf{x} \in X} \{f_k(\mathbf{x})\}$  and  $f_k^{\min} := \min_{\mathbf{x} \in X} \{f_k(\mathbf{x})\}$

- Then,  $\tilde{f}_k(\mathbf{x}) \in [0, 1]$  for all  $\mathbf{x} \in X$ , so that the functions  $\tilde{f}_k$  can be compared on a common scale