## TMA947/MAN280 <br> APPLIED OPTIMIZATION

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## Question 1

## (the simplex method)

$(2 \mathbf{p}) \quad$ a) To transform the problem to standard form, the free variable $x_{1}$ must be replaced by the non-negative variables $x_{1}^{+}$and $x_{1}^{-}$such that $x_{1}:=x_{1}^{+}-x_{1}^{-}$. A non-negative slack variable $s_{1}$ in the first constraint and a non-negative slack variable $s_{2}$ in the second constrained must be subtracted.
A BFS cannot be found directly, hence begin with phase 1 with artificial variables $a_{1} \geq 0$ added in the first constraint and $a_{2} \geq 0$ in the second constraint. The objective is to minimize $w=a_{1}+a_{2}$. Start with the BFS given by $\left(a_{1}, a_{2}\right)^{\mathrm{T}}$. In the first iteration of the simplex algorithm, $x_{1}^{-}$is the only variable with a negative reduced cost $(-3)$, and is therefore the only eligable incoming variable. The minimum ratio test shows that either $a_{1}$ or $a_{2}$ can be removed from the basis. We choose $a_{1}$ as the outgoing variable and update the basic variables to $\boldsymbol{x}_{B}=\left(x_{1}^{-}, a_{2}\right)^{\mathrm{T}}$. By computing the reduced costs, we see that $x_{3}$ is the only non-basic variable with negative reduced cost $(-2)$ and $x_{3}$ is chosen as incoming variable. The minimum ratio test shows that $a_{2}$ should leave the basis. By updating the basis and computing the reduced costs we see that we are now optimal with $w^{*}=0$ and we proceed to phase 2.
The BFS is given by $\boldsymbol{x}_{B}=\left(x_{1}^{-}, s_{1}\right)^{\mathrm{T}}, \boldsymbol{x}_{N}=\left(x_{1}^{+}, x_{2}, s_{2}\right)^{\mathrm{T}}$ and the reduced costs with the phase 2 cost vector $\boldsymbol{c}=(-1,1,1,-1,-1)^{\mathrm{T}}$ are

$$
\tilde{\boldsymbol{c}}_{\left(x_{1}^{+}, x_{2}, s_{2}\right)}^{\mathrm{T}}=(0,2,1 / 2) \geq \mathbf{0}^{3},
$$

and thus the optimality condition is fulfilled for the current basis. We have $\boldsymbol{x}_{B}^{*}=(1 / 2,0)^{\mathrm{T}}$, or in the original variables, $\boldsymbol{x}^{*}=\left(x_{1}, x_{2}\right)^{*}=(-1 / 2,0)^{\mathrm{T}}$, with the optimal value $z^{*}=1 / 2$.
$(\mathbf{1 p}) \quad$ b) Since there is an optimal solution to the problem, Strong duality guarantees the existence of a dual optimal solution. The expression for this is $\boldsymbol{y}^{* \mathrm{~T}}=\boldsymbol{c}_{B}^{\mathrm{T}} \boldsymbol{B}^{-1}=(0,1 / 2)$. However, the optimal basis is degenerate and it is possible to replace the zero-valued basic variable $s_{2}$ with a nonbasic variable as long as the basis matrix $\boldsymbol{B}$ still has linear independent columns. We see that it is possible to replace $s_{1}$ with $s_{2}$ which gives us $\boldsymbol{y}^{* \mathrm{~T}}=\boldsymbol{c}_{B}^{\mathrm{T}} \boldsymbol{B}^{-1}=(2 / 3,1 / 6)$ or to replace $s_{1}$ with $x_{2}$ which gives us $\boldsymbol{y}^{* \mathrm{~T}}=\boldsymbol{c}_{B}^{\mathrm{T}} \boldsymbol{B}^{-1}=(1,0)$. Also, any convex combination between these three dual optimal solution is also dual optimal.

EXAM SOLUTION
TMA947/MAN280 - APPLIED OPTIMIZATION

## (3p) Question 2

(convergence of an exterior penalty method)
See Theorem 13.3 in The Book.

## Question 3

(Weierstrass)
(1p) a) Yes it does. The function is continuous on a closed set. Also, we observe that the function is weakly coercice, i.e., when $\|\boldsymbol{x}\| \rightarrow \infty$, then $f_{1}(\boldsymbol{x}) \rightarrow \infty$. Weierstrass' theorem now guarantees that a global minimum exists.
$(1 \mathbf{p}) \quad$ b) No it does not. The function is not lower semi-continuous, so we cannot invoke Weierstrass' thereom. We observe that along the arc $\boldsymbol{x}=(t, 1,0)^{\mathrm{T}}$, where $t \rightarrow 0_{+}, f_{2}(\boldsymbol{x}) \rightarrow-\infty$.
$\mathbf{( 1 p )}$ c) No it does not. The function is not weakly coercive, so we cannot invoke Weierstrass theorem. We observe that along the $\operatorname{arc} \boldsymbol{x}=(t,-\sqrt{-t})^{\mathrm{T}}$, where $t \rightarrow-\infty, f_{3}(\boldsymbol{x}) \rightarrow-\infty$.

## Question 4

(modeling)
$(\mathbf{1 p})$ a) Introduce the variables:
$x_{i}$ : SEK aid given to country $i$.
The objective is to

$$
\max \frac{\sum_{i \in \mathcal{N}}\left(a_{i}+c_{i} x_{i}\right) p_{i}}{\sum_{i \in \mathcal{N}} p_{i}}
$$

and the constraints are

$$
\begin{align*}
& \sum_{i \in \mathcal{N}} x_{i} \leq 0.01 b  \tag{1}\\
& 0 \leq x_{i} \leq d_{i}, \quad \forall i \in \mathcal{N} \tag{2}
\end{align*}
$$

(1p) b) Introduce the binary variables
$y_{i}$ : with value one if country $i$ receives aid, zero otherwise.

Modify the constraints (2) into

$$
\begin{equation*}
0 \leq x_{i} \leq d_{i} y_{i}, \quad \forall i \in \mathcal{N} . \tag{3}
\end{equation*}
$$

Introduce an additional constraint

$$
\begin{equation*}
\sum_{i \in \mathcal{N}} y_{i} \leq M \tag{4}
\end{equation*}
$$

(1p) c) Introduce the variable
$w:$ minimal HDI of the countries considered.

Introduce the constraints

$$
\begin{equation*}
a_{i}+c_{i} x_{i} \geq w \quad \forall i \in \mathcal{N} . \tag{5}
\end{equation*}
$$

Change the objective function into
$\max w$.

## (3p) Question 5

(the Frank-Wolfe method)
Iteration 1: $\boldsymbol{x}_{0}=(0,0)^{\mathrm{T}}, f\left(\boldsymbol{x}_{0}\right)=0$. It is feasible, so we get an upper bound: $[L B D, U B D]=(-\inf , 0]$. We have that $\nabla f\left(\boldsymbol{x}_{0}\right)=(-10,-4)^{\mathrm{T}}$. Solve the LP $\min _{y \in X} \nabla f\left(\boldsymbol{x}_{0}\right)^{\mathrm{T}} \boldsymbol{y}$. The solution is obtained at $\boldsymbol{y}_{0}=(2,1)^{\mathrm{T}}$. The search direction is $\boldsymbol{p}_{0}=\boldsymbol{y}_{0}-\boldsymbol{x}_{0}=(2,1)^{\mathrm{T}}$. Since $f$ is convex, $g(\boldsymbol{y}):=f\left(\boldsymbol{x}_{0}\right)+\nabla f\left(\boldsymbol{x}_{0}\right)^{\mathrm{T}}(\boldsymbol{y}-$ $\left.\boldsymbol{x}_{0}\right) \leq f(\boldsymbol{y})$ for all $\boldsymbol{y} \in \mathbb{R}^{2}$. The LP problem $\min _{y \in X} g(\boldsymbol{y})$ is a relaxation of the original problem, hence an optimal objective value gives a lower bound. The objective value is $f\left(\boldsymbol{x}_{0}\right)+\nabla f\left(\boldsymbol{x}_{0}\right)^{\mathrm{T}}\left(\boldsymbol{y}_{0}-\boldsymbol{x}_{0}\right)=0+(-10,-4)(2,1)^{\mathrm{T}}=-24$. Hence $[L B D, U B D]=[-24,0]$. Line search: $\phi(\alpha)=f\left(x_{0}+\alpha \boldsymbol{p}_{0}\right)=f((2 \alpha, \alpha))=\ldots=$ $8 \alpha^{2}-24 \alpha . \phi^{\prime}(\alpha)=16 \alpha-24=0 . \quad \alpha=24 / 16>1$ Take a unit step: $\alpha=1$. Hence, the next point is $\boldsymbol{x}_{1}=(2,1)^{\mathrm{T}}$.

Iteration 2: $f\left(\boldsymbol{x}_{1}\right)=-10$. So $[L B D, U B D]=[-24,-10] . \nabla f\left(\boldsymbol{x}_{1}\right)=(-5,2)^{\mathrm{T}}$, the optimal solution to $\min _{y \in X} \nabla f\left(\boldsymbol{x}_{1}\right)^{\mathrm{T}} \boldsymbol{y}$ is obtained at $\boldsymbol{y}_{1}=(2,0)^{\mathrm{T}}$, hence the
search direction will be $\boldsymbol{p}_{1}=\left(\boldsymbol{y}_{1}-\boldsymbol{x}_{1}\right)=(0,-1)^{\mathrm{T}}$. Since $f\left(\boldsymbol{x}_{1}\right)+\nabla f\left(\boldsymbol{x}_{1}\right)^{\mathrm{T}}\left(\boldsymbol{y}_{1}-\right.$ $\left.\boldsymbol{x}_{1}\right)=-12$, we have $[L B D, U B D]=[-12,-10]$. Line search: $\phi(\alpha)=f\left(\boldsymbol{x}_{1}+\right.$ $\left.\alpha \boldsymbol{p}_{1}\right)=2^{2}+2(1-\alpha)+2(1-\alpha)^{2}-20-4(1-\alpha) . \phi^{\prime}(\alpha)=-2-4(1-\alpha)+4=0$. $\alpha=1 / 2$. $\boldsymbol{x}_{3}=(2,1 / 2)^{\mathrm{T}}$.

The point $\boldsymbol{x}_{3}$ is a KKT point since $\nabla f\left(\boldsymbol{x}_{3}\right)=(-11 / 2,0)$ and the active constraint is $g(\boldsymbol{x})=x_{1}-2$ with $\nabla g\left(\boldsymbol{x}_{3}\right)=(1,0)^{\mathrm{T}}$. The objective function is convex (eigenvalues of the Hessian are all non-negative) and the feasible region is a polyhedron, so the problem is convex. A KKT point is sufficient for optimality in convex problems, and $\boldsymbol{x}_{3}$ is therefore an optimal point.

## (3p) Question 6

(convex problem)
We first conclude that the feasible set is convex. The functions $g_{i}, i=1,2,3$, are affine, hence convex. The function $g_{4}$ is convex, since its Hessian matrix is constant and diagonal with diagonal entries 0,4 , and 8 , which all are nonnegative. In each of these four cases, the constraint is of the form $g_{i}(\boldsymbol{x}) \leq 0$; hence, by Proposition 3.44, each feasible set is convex, and moreover their intersection is convex by Proposition 3.3.

To establish that the objective function is convex on the convex feasible set of the problem at hand, we consider the function terms one by one. The function $\boldsymbol{x} \mapsto-\ln \left(x_{1}+x_{2}\right)$ is of the form $-\ln t$, where $t=x_{1}+x_{2}$. Introducing, for simplicity, $t$ as an additional variable, we notice that the equation just given is linear and therefore represent a further convex constraint. Due to constraints 1 and $2, t>0$ on the feasible set, whence $\ln$ is well defined there. Finally, $-\ln t$ is a (strictly) convex function on this domain. The second term of the objective is $x_{3} \ln x_{3}$. Again, we get from the third constraint that $x_{3}>0$, and hence the term is well-defined. Taking its derivative with respect to $x_{3}$ we get $1+\ln x_{3}$, and its derivative is, in turn, $1 / x_{3}$, which is positive. Hence, the objective function is a sum of two convex functions and therefore is convex. We are done.

## (3p) Question 7

(linear programming duality)
Suppose, for example, that $X$ is bounded. Then, there exists a bounded optimal solution for every value of the objective coefficient vector $\boldsymbol{c}$. Therefore, its dual must also have bounded optimal solutions for every value of $\boldsymbol{c}$. It follows that the dual problem must have feasible solutions for every $\boldsymbol{c}$. Consider the cone

$$
C:=\left\{\boldsymbol{y} \in \mathbb{R}^{m} \mid \boldsymbol{A}^{\mathrm{T}} \boldsymbol{y} \leq \mathbf{0}^{n}, \quad \boldsymbol{y} \geq \mathbf{0}^{m}\right\} .
$$

By the Representation Theorem, the set $Y$ is bounded if and only if $C$ contains only the zero vector. By the above, the set $\left\{\boldsymbol{y} \in \mathbb{R}^{m} \mid \boldsymbol{A}^{\mathrm{T}} \boldsymbol{y} \leq-\boldsymbol{e}, \quad \boldsymbol{y} \geq \mathbf{0}^{m}\right\}$, where $\boldsymbol{e}$ is the $m$-vector of ones, is non-empty. Clearly, any of its members are non-zero, and moreover they belong to the larger set $C$. Hence, $C$ does not only contain the zero vector, and so $Y$ is unbounded.

The case where one assumes that $Y$ is bounded is treated similarly.

