# TMA947/MAN280 <br> OPTIMIZATION, BASIC COURSE 

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Examiner: Michael Patriksson

## Question 1

(the simplex method)
$(\mathbf{2 p}) \quad$ a) We first rewrite the problem on standard form. We introduce variables $x_{1}^{+}$ and $x_{1}^{-}$and let $x_{1}=x_{1}^{+}-x_{1}^{-}$. We also add slack variables $s_{1}$ and $s_{2}$.

$$
\begin{array}{lrrrrrr}
\operatorname{minimize} z= & 2 x_{1}^{+} & -2 x_{1}^{+} & -x_{2} & & & \\
\text { subject to } & -x_{1}^{+} & -x_{1}^{+} & +2 x_{2} & -s_{1} & & =2 \\
& -x_{1}^{+} & +x_{1}^{-} & +x_{2} & & +s_{2} & =3 \\
& x_{1}^{+}, & x_{1}^{-}, & x_{2}, & s_{1}, & s_{2} & \geq 0 .
\end{array}
$$

In phase I the artificial variable $a$ is added in the first constraint, $s_{2}$ is used as the second basic variable in order to obtain a unit matrix as the first basis. We obtain the phase I problem

$$
=2 .
$$

The starting BFS is thus $\left(a, s_{2}\right)^{\mathrm{T}}$. Calculating the vector of reduced costs for the non-basic variables $x_{1}^{+}, x_{1}^{-}, x_{2}, s_{1}$ yields $(1,-1,-2,1)^{\mathrm{T}}$. Least reduced cost implies that $x_{2}$ is the entering variable. The minimum ratio test shows that $a$ should leave the basis. We thus have a BFS without artificial variables, and may proceed with phase II.
We have the basic variables $\left(x_{2}, s_{2}\right)$. The vector of reduced costs for the non-basic variables $x_{1}^{+}, x_{1}^{-}$and $s_{1}$ is $(3 / 2,-3 / 2,-1 / 2)$. We let $x_{1}^{-}$enter the basis. The minimum ratio test implies that $x_{2}$ leaves the basis. We now have $x_{1}^{-}, s_{2}$ as basic variables. The vector of reduced costs for the non-basic variables $x_{1}^{+}, x_{2}$ and $s_{1}$ is $(0,3,-2)^{\mathrm{T}}$. Thus $s_{1}$ enters the basis. Minimum ratio implies that $s_{2}$ leaves basis. We now have $x_{1}^{-}, s_{1}$ as basic variables. The vector of reduced costs for the non-basic variables $x_{1}^{+}, x_{2}$ and $s_{2}$ is $(0,1,2)^{\mathrm{T}}$.

Since the reduced costs are all non-negative, the current BFS is optimal. Returning to the original variables, we obtain $\left(x_{1}, x_{2}\right)=(-3,0)$ as the optimal solution and -6 as the optimal value.
$\mathbf{( 1 p )}$ b) In the optimal BFS, the reduced cost corresponding to $x_{1}^{+}$is zero. Therefore, we can let $x_{1}^{+}$enter the basis without changing the objective. We do not obtain any leaving variable as minimum ratio implies that the problem is unbounded in that direction. This is simply the increasing $x_{1}^{+}$and

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increasing $x_{1}^{-}$by the same amount (which can be any positive number). So the problem in standard form does not have a unique optimal solution, but the problem formulated in the original variables does since all these solutions correspond to $(-3,0)$. Replacing one free variable with two positive variables always implies that each solution is non-unique in the sense described above.

## (3p) Question 2

(LP duality)
i) We interpret the variables $x_{i j}$ as the flow from node $i$ to node $j$ of the graph. The first set of constraints guarantee that the flow going into a node equals the flow leaving that node (i.e. flow balance) except at node 1 where $s$ units of flow enters and at node $n$ where $s$ unit of flow leaves. The second set of constraints imply a limitof flow on all arcs and that flow can not be negative. To summerize, we push $s$ units of flow into the graph at node 1 and take out $s$ units of flow at node $n$ and we maximize $s$. This can be interpreted as maximizing the flow through the graph from node 1 to node $n$.
ii) Let $v_{i}$ for $i \in\{1, \ldots, n\}$ denote the dual variables corresponding to the first set of constraints and $w_{i j}$ for $(i, j) \in A$ denote the dual variables corresponding to the second set of constraints.. The dual then becomes:

$$
\begin{aligned}
& \text { maximize } \sum_{(i, j) \in A} w_{i j} c_{i j} \\
& \text { subject to } w_{i j}-v_{i}+v_{j} \geq 0,(i, j) \in A, \\
& v_{1}-v_{n}=1, \\
& w_{i j} \geq 0,(i, j) \in A .
\end{aligned}
$$

iii) We are maximizing with positive costs on $w_{i j}$, therefore we would want to put them all to zero. Unfortunately, this is not feasible. First, the constraints imply that if $w_{i j}=0$ then $v_{i} \leq v_{j}$. Second $v_{1}=v_{n}+1$. But by following a path from node 1 to node $n$ throught the graph we obtain that $v_{1} \leq v_{n}$. Therefore, all paths that connect 1 with $n$ need to contain an $\operatorname{arc}(i, j)$ where $w_{i j}=1$ in order to break the chain of inequalities such that $v_{1}=v_{n}+1$ is possible. This will generate cost $c_{i j}$. If we let $c_{i j}$ correspond to the cost of including the arc $(i, j)$ in a cut, the dual is to find the minimal cost cut between nodes 1 and $n$.

## Question 3

(modeling)
$(\mathbf{1 p}) \quad$ a) Since $P^{k}$ is convex for $k=1, \ldots, K$, we have that

$$
\operatorname{conv}\left(\bigcap_{k=1}^{K} P^{k}\right)=\bigcap_{k=1}^{K} P^{k}=\left\{\boldsymbol{x} \in \mathbb{R}^{n} \mid A^{k} \boldsymbol{x} \leq \boldsymbol{b}^{k}, k=1, \ldots, K\right\} .
$$

Hence, we can write the optimization problem as that to

$$
\begin{array}{cl}
\operatorname{minimize} & \boldsymbol{c}^{\mathrm{T}} \boldsymbol{x}, \\
\text { subject to } & A^{k} \boldsymbol{x} \leq \boldsymbol{b}^{k}, \quad k=1, \ldots, K
\end{array}
$$

(2p)
b) Any point in $\operatorname{conv}\left(\bigcup_{k=1}^{K} P^{k}\right)$ can be represented as a convex combination of points in the individual polyhedra $P^{1}, \ldots, P^{K}$, i.e.,

$$
\boldsymbol{x}=\sum_{k=1}^{K} \lambda_{k} \boldsymbol{x}^{k}, \quad \text { with } \quad \sum_{k=1}^{K} \lambda_{k}=1, \quad \lambda_{k} \geq 0
$$

where $\boldsymbol{x}^{k} \in P^{k}, k=1, \ldots, K$. Hence, we can formulate the optimization problem as that to

$$
\begin{array}{rlrl}
\operatorname{minimize} & \boldsymbol{c}^{\mathrm{T}} \boldsymbol{x} & \\
\text { subject to } & & \boldsymbol{x} & =\sum_{k=1}^{K} \lambda_{k} \boldsymbol{x}^{k}, \\
A^{k} \boldsymbol{x}^{k} & \leq \boldsymbol{b}^{k}, \quad k=1, \ldots, K, \\
\sum_{k=1}^{K} \lambda_{k} & =1, \\
\lambda_{k} & \geq 0, \quad k=1, \ldots, K,
\end{array}
$$

where $\boldsymbol{x}, \boldsymbol{x}^{1}, \ldots, \boldsymbol{x}^{K}$ and $\lambda_{1}, \ldots, \lambda_{K}$ are the variables in the optimization model.
(This model can actually be extended to a linear model by the variable substitution $\overline{\boldsymbol{x}}^{k}=\lambda_{k} \boldsymbol{x}^{k}$.)

## Question 4

(exterior-penalty methods)
$\mathbf{( 1 p )}$ a) By applying the KKT conditions on the problem, we obtain the unique solution $\boldsymbol{x}^{*}=(1 / 2,1 / 2)^{\mathrm{T}}$ and $\lambda^{*}=-1$.
$(1 \mathbf{p}) \quad$ b) Applying the exterior quadratic penalty method, we get the unconstrained problem

$$
\min _{\boldsymbol{x} \in \mathbb{R}^{2}}\left(f(\boldsymbol{x})+\nu h(\boldsymbol{x})^{2}\right)=\min _{\boldsymbol{x} \in \mathbb{R}^{2}}\left(x_{1}^{2}+x_{2}^{2}+\nu\left(x_{1}+x_{2}-1\right)^{2}\right) .
$$

Setting the gradient to zero we obtain

$$
\begin{aligned}
& 2 x_{1}+2 \nu x_{1}+2 \nu x_{2}-2 \nu=0 \\
& 2 x_{2}+2 \nu x_{1}+2 \nu x_{2}-2 \nu=0
\end{aligned}
$$

with solution $\boldsymbol{x}_{\nu}=\frac{\nu}{1+2 \nu}(1,1)^{\mathrm{T}}$, for $\nu>0$. Letting $\nu \rightarrow \infty$ we see that the sequence $\boldsymbol{x}_{\nu}$ converges to $(1 / 2,1 / 2)^{\mathrm{T}}$ which is the unique optimal solution.
c) We note that the solution $\boldsymbol{x}_{\nu}$ fulfills $\nabla f\left(\boldsymbol{x}_{\nu}\right)+\left[2 \nu h\left(\boldsymbol{x}_{\nu}\right)\right] \nabla h\left(\boldsymbol{x}_{\nu}\right)=\mathbf{0}$. So a Lagrange multiplier estimate comes from $\lambda_{\nu}:=2 \nu h\left(\boldsymbol{x}_{\nu}\right)$. Insertion from b) yields $\lambda_{\nu}=\frac{-2 \nu}{1+2 \nu}$ which tends to $\lambda^{*}=-1$ when $\nu \rightarrow \infty$.

## Question 5

(linear programming: existence of optimal solutions)
$(2 \mathbf{p}) \quad$ a) This is the first part of Theorem 8.10.
$(1 \mathbf{p}) \quad$ b) This is the second part of Theorem 8.10.

## Question 6

(basic facts in optimization)
(1p) a) False. Any infeasible linear program will result in a phase I problem having a positive optimal value.
$(\mathbf{1 p}) \quad$ b) False. If $f$ is strictly concave, and if at some $\boldsymbol{x} \in \mathbb{R}^{n}$, the Hessian matrix $\nabla^{2} f(\boldsymbol{x})$ happens to be negative definite, then a search direction is welldefined, but it defines an ascent direction.
$\mathbf{( 1 p )} \quad$ c) False. Consider the linear program to minimize $x_{2}$ subject to the constraints $0 \leq x_{j} \leq 4, j=1,2$, and the additional constraint that $x_{1}+x_{2} \leq 2$. This problem has the optimal solution set $X^{*}=\left\{\boldsymbol{x} \in \mathbb{R}^{2} \mid x_{1} \in[0,2] ; x_{2}=0\right\}$. At the optimal solution $\boldsymbol{x}^{*}=(1,0)^{\mathrm{T}}, x_{1}+x_{2}<2$ holds. Believing that this means that the constraint $x_{1}+x_{2} \leq 2$ therefore is redundant results, however, in a grave mistake, as the new problem, having the constraints $0 \leq x_{j} \leq 4, j=1,2$, has the optimal set $X_{\text {new }}^{*}=\left\{\boldsymbol{x} \in \mathbb{R}^{2} \mid x_{1} \in[0,4] ; x_{2}=\right.$ $0\}$.

## Question 7

## (Lagrangian duality)

$(1 \mathbf{p}) \quad$ a) At $\boldsymbol{x}^{*}=(3,2)^{\mathrm{T}}$, the constraints $g_{1}(\boldsymbol{x}):=x_{1}+x_{2} \leq 5$ and $g_{2}(\boldsymbol{x}):=x_{1}-$ $3 \leq 0$ are active. Introducing Lagrange multipliers for them, we study the equation

$$
\nabla f\left(\boldsymbol{x}^{*}\right)+\mu_{1} \nabla g_{1}\left(\boldsymbol{x}^{*}\right)+\mu_{1} \nabla g_{1}\left(\boldsymbol{x}^{*}\right)=\mathbf{0}^{2}
$$

at $\boldsymbol{x}^{*}=(3,2)^{\mathrm{T}}$, that is,

$$
\binom{-2}{-1}+\mu_{1}\binom{1}{1}+\mu_{2}\binom{1}{0}=\binom{0}{0} .
$$

Hence, $\mu_{1}=\mu_{2}=1$.
As the KKT conditions are satisfied and the problem stated is convex, $\boldsymbol{x}^{*}=(3,2)^{\mathrm{T}}$ is indeed a globally optimal solution.
$(2 \mathbf{p}) \quad$ b) With the Lagrangian $L(\boldsymbol{x}, \mu):=f(\boldsymbol{x})+\mu\left(x_{1}+x_{2}-5\right)$, we obtain the explicit Lagrangian dual function as follows:
First, the subproblem solution is that

$$
\begin{aligned}
& x_{1}= \begin{cases}3, & \text { if } \mu \in[0,2], \\
5-\mu, & \text { if } \mu \in[2,5], \\
0, & \text { if } \mu \in[2, \infty),\end{cases} \\
& x_{2}= \begin{cases}3-\mu, & \text { if } \mu \in[0,3], \\
0, & \text { if } \mu \in[3, \infty)\end{cases}
\end{aligned}
$$

Hence,

$$
q(\mu)= \begin{cases}-\frac{1}{2} \mu^{2}+\mu+2, & \text { if } \mu \in[0,2] \\ -\mu^{2}+3 \mu, & \text { if } \mu \in[2,3] \\ -\frac{1}{2} \mu^{2}+\frac{9}{2}, & \text { if } \mu \in[3,5] \\ -5 \mu+17, & \text { if } \mu \in[5, \infty)\end{cases}
$$

which is to be maximized over non-negative values of $\mu$.
Second, the derivative of $q$ is

$$
q^{\prime}(\mu)= \begin{cases}-\mu+1, & \text { if } \mu \in[0,2] \\ -2 \mu+3, & \text { if } \mu \in[2,3] \\ -\mu, & \text { if } \mu \in[3,5] \\ -5, & \text { if } \mu \in[5, \infty)\end{cases}
$$

whence the optimal solution is found where $q^{\prime}$ changes sign, namely at $\mu^{*}=1$.
This multiplier value is indeed the same as the one found as $\mu_{1}$ in a). The corresponding primal solution then is $\boldsymbol{x}\left(\mu^{*}\right)=(3,2)^{\mathrm{T}}$. Further, $q\left(\mu^{*}\right)=5 / 2$ equals the value of $f\left(\boldsymbol{x}^{*}\right)$, whence strong duality is fulfilled.

