

A few facts you need to know to be able to read the extra material  
on Gödel's incompleteness theorem

Last change: 10 December

- The set of natural numbers is denoted by  $D_N$ .
- Peano's axioms are denoted by  $\mathcal{N}$  and  $\vdash_{\mathcal{N}}A$  means that the formula  $A$  is derivable from  $\mathcal{N}$ .
- The function symbol for the successor function is  $'$ , van Dalen uses  $S$ .
- A  $k$ -place relation  $R$  on the natural numbers is *representable* in  $\mathcal{N}$  by the formula  $A(x_1, \dots, x_k)$  if the following two conditions hold:
  - if  $R(n_1, \dots, n_k)$  then  $\vdash_{\mathcal{N}}A(0^{(n_1)}, \dots, 0^{(n_k)})$ ,
  - if  $R(n_1, \dots, n_k)$  is false then  $\vdash_{\mathcal{N}}\neg A(0^{(n_1)}, \dots, 0^{(n_k)})$ .
- A  $k$ -place relation  $R$  is *decidable* if there exists an algorithm  $Decide_R$  (i. e. a program in your favourite programming language), such that
  - if  $R(n_1, \dots, n_k)$  then  $Decide_R(n_1, \dots, n_k) = T$ ,
  - if  $R(n_1, \dots, n_k)$  is false then  $Decide_R(n_1, \dots, n_k) = F$
- We will use the following fact without proof:  
If the relation  $R$  is decidable, then it is representable in  $\mathcal{N}$ .