

### Extra material : Day 7

**Theorem (Dirac)** *Let  $G$  be a simple graph on  $n$  vertices in which every vertex has degree at least  $n/2$ . Then  $G$  has a Hamilton cycle.*

PROOF : Let  $G$  be a counterexample with the largest possible number of edges - we shall derive the contradiction that  $G$  contains a Hamilton cycle after all. By maximality of  $|E(G)|$ , if  $x$  and  $y$  are any two non-adjacent vertices in  $G$ , then addition of the edge  $\{x, y\}$  would create a Hamilton cycle in  $G$ . In other words, there must exist a path from  $x$  to  $y$  in  $G$  passing through all the vertices. Denote the path as

$$y = x_1 \rightarrow x_2 \rightarrow \cdots \rightarrow x_{n-1} \rightarrow x_n = z.$$

We consider two sets  $S$  and  $T$ , defined as follows :

$$\begin{aligned} S &= \{i : 1 \leq i \leq n-1 \text{ and } \{y, x_{i+1}\} \in E(G)\}, \\ T &= \{i : 1 \leq i \leq n-1 \text{ and } \{x_i, z\} \in E(G)\}. \end{aligned}$$

On the one hand, both  $S$  and  $T$  are subsets of  $\{1, \dots, n-1\}$ , so

$$|S \cup T| \leq n-1. \tag{1}$$

On the other hand,

$$\begin{aligned} |S| &= \deg(y) \geq n/2, \\ |T| &= \deg(z) \geq n/2, \end{aligned}$$

which implies that

$$|S| + |T| \geq n. \tag{2}$$

From (1) and (2), we can deduce that  $S \cap T$  is non-empty. Let  $i \in S \cap T$ . Then the following is a Hamilton cycle in  $G$  :

$$y \rightarrow x_2 \rightarrow \cdots \rightarrow x_i \rightarrow z \rightarrow x_{n-1} \rightarrow \cdots \rightarrow x_{i+1} \rightarrow y.$$

With that, the proof of the theorem is complete.