

## Extra material : Day 8

The extra stuff today consists of some brief remarks about planar graphs.

**DEFINITION :** A graph is said to be *planar* if it can be drawn in the plane without any two edges crossing. If the graph is actually drawn in such a way then it is called *plane*.

There is a nice simple characterisation of planarity due to Kuratowski. First, we need a definition :

**DEFINITION :** Let  $G, H$  be graphs.  $G$  is said to contain  $H$  as a *minor* if  $H$  can be obtained from  $G$  by a number of edge-contractions and deletions.

**Theorem (Kuratowski)** *A graph is planar if and only if it does not contain either  $K_5$  or  $K_{3,3}$  as a minor.*

**Theorem (4-colour theorem)** *If  $G$  is a planar graph, then  $\chi(G) \leq 4$ .*

This famous theorem was first proven in the mid '70s. The main reason it is so famous is that it was the first long-standing open problem which was solved with substantial help from computers. The idea of the proof was to exhibit a finite collection of 'minimal possible counterexamples', that is, a finite collection of planar graphs such that, if each of these were 4-colourable, then it would follow that all planar graphs were. Such a list, consisting of about 10,000 graphs, was produced. And then the computers were called in to check that each of these graphs was indeed 4-colourable. In the '90s a simplified version of the proof was obtained by Seymour, Robertson, ? and ??, which reduced the number of 'test graphs' to about 600.

**Theorem (Euler)** *Let  $G$  be a plane graph. Let  $V$  denote the number of vertices in  $G$ ,  $E$  the number of edges and  $R$  the number of regions into which  $G$  divides the plane, where we include the unbounded region outside  $G$  as one such. Then*

$$V - E + R = 2.$$

**PROOF :** I didn't go through it, but the standard method is to use induction on the number of edges in  $G$ .

REMARK : Let  $X$  be any compact, smooth surface in  $\mathbf{R}^3$ . Let  $G$  be any graph drawn on this surface in such a way that no two edges cross and such that  $G$  divides the surface into simply-connected regions. If we compute  $V - E + R$ , then the result depends only on  $X$ , but not on  $G$ . This number is called the *Euler characteristic* of the surface  $X$  and denoted  $\chi(X)$ . See a textbook on Differential Geometry for more.