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problem

algorithms for the Lagrangian dual
Lecture 3: Lagrangian duality and

$$S^R \subset S.$$

where $f_R : \mathbb{R}^n \rightarrow \mathbb{R}$ is a function with $f \leq f_R$ on S , and

(2b) subject to $x \in S^R$,

(2a) $f_* = \inf_x^x f_R(x)$,

- A relaxation to (1) has the following form: find

where $f : \mathbb{R}^n \rightarrow \mathbb{R}$ given function, $S \subset \mathbb{R}^n$.

(1b) subject to $x \in S$,

(1a) $f_*(x) = \inf_x^x f(x)$

- Problem: find

The Relaxation Theorem

• Relaxation Theorem: (a) [relaxation] $f_*^R \leq f_*$.
 bound on f_* .

bounds on f_* : Lagrangian relaxation yields lower
 • Applications: exterior penalty methods yield lower

$$(3) \quad S \ni x \quad f_R(x^R) \geq f(x)$$

• Proof portioon. For (c), note that
 then x^R_* is an optimal solution to (1) as well.

$$(3) \quad x^R_* \in S \quad f_R(x^R_*) = f(x^R_*)$$

optimal solution, x^R_* , for which it holds that
 (c) [optimal relaxation] If the problem (2) has an
 (b) [infeasibility] If (2) is infeasible, then so is (1).
 • Relaxation Theorem: (a) [relaxation] $f_*^R \leq f_*$.

problem has at least one feasible solution.
that is, that f is bounded from below and that the

$$(5) \quad -\infty < f_* < \infty,$$

- For this problem, we assume that

given functions, and $X \subseteq \mathbb{R}^n$.

where $f : \mathbb{R}^n \rightarrow \mathbb{R}$ and $g_i : \mathbb{R}^n \rightarrow \mathbb{R}$ ($i = 1, 2, \dots, m$) are

$$g_i(x) \geq 0, \quad i = 1, \dots, m, \quad (4c)$$

(4b) subject to $x \in X$

$$(4a) \quad f_*(x) := \inf_x f(x)$$

- Consider the optimization problem to find

Lagrangian relaxation

- We call the vector $\boldsymbol{u}_* \in \mathbb{R}^m$ a Lagrange multiplier if it is non-negative and if $f_* = \inf_{\mathcal{X}} L(\boldsymbol{x}, \boldsymbol{u}_*)$ holds.

$$(9) \quad \cdot(\boldsymbol{x}) \boldsymbol{g}_{\perp} \boldsymbol{u} + (\boldsymbol{x}) f = (\boldsymbol{x})^i g^i \boldsymbol{u}^i \sum_m^{i=1} + (\boldsymbol{x}) f =: L(\boldsymbol{x}, \boldsymbol{u})$$

- For a vector $\boldsymbol{u} \in \mathbb{R}^m$, we define the Lagrange function

Lagrange multipliers and global optima

- Let μ_* be a Lagrange multiplier. Then, x_* is an optimal solution to (4) if and only if x_* is feasible in

(4) and

$$x_* \in \arg \min_{x \in X} L(x, \mu_*) \text{ and } \mu_*^i g_i(x_*) = 0, \quad i = 1, \dots, m.$$

(7)

$$x_* \in \arg \min_{x \in X} L(x, \mu_*) \text{ and } \mu_*^i g_i(x_*) = 0, \quad i = 1, \dots, m.$$

- Notice the resemblance to the KKT conditions! If $X = \mathbb{R}^n$ and all functions are in C_1 then $x_* \in \arg \min_{x \in X} L(x, \mu_*)$ is the same as the force equilibrium condition, the first row of the KKT conditions. The second item, " $\mu_*^i g_i(x_*) = 0$ for all i " is the complementarity conditions.

- Seems to imply that there is a hidden convexity assumption here. Yes, there is. We show a Strong Duality Theorem later.

The Lagrangian dual problem associated with the

Lagrangian relaxation

is the Lagrangian dual function. The Lagrangian dual problem is to

$$(8) \quad \text{inf}_{x \in X} L(x, \boldsymbol{\mu}) =: L(\boldsymbol{\mu})$$

(q6) subject to $\boldsymbol{\mu} \geq \mathbf{0}_m$.

(9a) maximize $b(\boldsymbol{\mu})$,

For some $\boldsymbol{\mu}$, $b(\boldsymbol{\mu}) = -\infty$ is possible; if this is true for all $\boldsymbol{\mu} \geq \mathbf{0}_m$,

$$\cdot \infty = (\boldsymbol{\mu}) b \text{ supremum } b =: {}^*_b$$

optimal solution.

optimal solution can be used to generate a primal

- But we need still to show how a Lagrangian dual

works!

(we indeed maximize a concave function) is very good

- That the Lagrangian dual problem always is convex

□

concave on D^y .

- The effective domain D^y of y is convex, and y is

$$D^y := \{ \boldsymbol{u} \in \mathbb{R}^m \mid y(\boldsymbol{u}) < \infty \}.$$

- The effective domain of y is

- Let \mathbf{x} and \mathbf{u} be feasible in (4) and (9), respectively.
 - In particular,
 $\cdot(\mathbf{x})_f \geq (\mathbf{u})_b$
 $\cdot(\mathbf{x})_f > (\mathbf{u})_b$
- Then,
- If \mathbf{x} and \mathbf{u} is optimal in its respective problem.

Weak Duality Theorem

- Weak duality is also a consequence of the Relaxation Theorem: For any $\boldsymbol{u} \geq \boldsymbol{0}_m$, let
 - (10a) $S := X \cup \{x \in \mathbb{R}^n \mid \boldsymbol{u} \leq \boldsymbol{b}(x)\}$
 - (10b) $S^R := X$
 - (10c) $f_R := L(\boldsymbol{u}, \cdot)$
 Apply the Relaxation Theorem.
- If $d_* = f_*$, there is no duality gap. If there exists a Lagrange multiplier vector, then by the weak duality theorem, there is no duality gap. There may be cases where no Lagrange multiplier exists even when there is no duality gap; in that case, the Lagrangian dual problem cannot have an optimal solution.

(11d)

$$u_i^* g_i(\mathbf{x}^*, \boldsymbol{\mu}^*) = 0, \quad i = 1, \dots, m. \quad (\text{Complementary slackness})$$

(11c) $\mathbf{0}_m \leq (\mathbf{x}^*, \boldsymbol{\mu}^*)$ (Primal feasibility)

(11b)

(11a) $\mathbf{0}_m \leq \boldsymbol{\mu}^*$ (Dual feasibility) $\mathbf{x}^* \in \arg \min_{\mathbf{x} \in X} L(\mathbf{x}, \boldsymbol{\mu}^*)$ (Lagrangian optimality)

(11a) $\mathbf{0}_m \leq \boldsymbol{\mu}^*$ (Dual feasibility)

and Lagrange multiplier if and only if

• The vector $(\mathbf{x}^*, \boldsymbol{\mu}^*)$ is a pair of optimal primal solution

Global optimality conditions

existence of Lagrange multipliers.

• If $\exists (\mathbf{x}^*, \boldsymbol{\mu}^*)$, equivalent to zero duality gap and

zero duality gap. conditions, the existence of Lagrange multipliers, and a • If $(\mathbf{x}_*, \boldsymbol{\mu}_*)$, equivalent to the global optimality holds.

$$(12) \quad L(\mathbf{x}, \boldsymbol{\mu}) \geq L(\mathbf{x}_*, \boldsymbol{\mu}_*) \quad \text{for all } (\mathbf{x}, \boldsymbol{\mu}) \in X \times \mathbb{R}^m_+,$$

Lagrangian function on $X \times \mathbb{R}^m_+$, that is, $\boldsymbol{\mu}_* \leq \mathbf{0}_m$, and $(\mathbf{x}_*, \boldsymbol{\mu}_*)$ is a saddle point of the and Lagrange multiplier if and only if $\mathbf{x}_* \in X$, • The vector $(\mathbf{x}_*, \boldsymbol{\mu}_*)$ is a pair of optimal primal solution

Saddle points

Strong duality for convex programs, introduction

- Results so far have been rather non-technical to achieve: convexity of the dual problem comes with very few assumptions on the original, primal problem, and the characterization of the primal-dual set of optimal solutions is simple and also quite easily established.
- In order to establish strong duality, that is, to establish sufficient conditions under which there is no duality gap, however takes much more.
- In particular, as is the case with the KKT conditions we need regularity conditions (that is, constraint qualifications), and we also need to utilize separation theorems.

- Consider problem (4), where $f : \mathbb{R}^n \rightarrow \mathbb{R}$ and g_i ($i = 1, \dots, m$) are convex and $X \subseteq \mathbb{R}^n$ is a convex set.
- Introduce the following Slater condition:
- Suppose that (5) and Slater's CQ (13) hold for the (convex) problem (4).
- (a) There is no duality gap and there exists at least one Lagrange multiplier μ_* . Moreover, the set of Lagrange multipliers is bounded and convex.

Strong duality Theorem

- If the infimum in (4) is attained at some \mathbf{x}_* , then the pair $(\mathbf{x}_*, \mathbf{u}_*)$ satisfies the global optimality condition (11).
 (b) If the infimum in (4) is attained at some \mathbf{x}_* , then the pair $(\mathbf{x}_*, \mathbf{u}_*)$ satisfies the global optimality condition (11).
 (c) If the functions f and g_i are in C_1 then the condition (11b) can be written as a variational inequality. If further X is open (for example, $X = \mathbb{R}^n$) then the conditions (11) are the same as the KKT conditions.
 • Similar statements for the case of also having linear equality constraints.
- If all constraints are linear we can remove the Slater condition.

Examples, I: An explicit, differentiable dual problem

- Consider the problem to

$$\underset{\mathbf{x}}{\text{minimize}} \quad f(\mathbf{x}) = x_1 + x_2,$$

subject to $x_1 + x_2 \leq 4$,

$$x_j \leq 0, \quad j = 1, 2.$$

$$\{ (x_1, x_2) \mid x_j \leq 0, j = 1, 2 \}.$$

$$\{ (x_1, x_2) \mid x_j \leq 0, j = 1, 2 \} =: X$$

- Let $g(\mathbf{x}) = -x_1 - x_2 + 4$ and

and $\mathbf{x}(u)$ is unique.

- everywhere (due to the fact that f, g are differentiable everywhere) and it is differentiable

- $b(u) = f(u) - g(u) = f(u) - \frac{u^2}{2}$.

- Substituting this expression into $g(u)$, we obtain that

$$x_1(u) = \frac{u}{2}, x_2(u) = \frac{u}{2}.$$

- For a fixed $u \geq 0$, the minimum is attained at

$$= 4u + \min_{\substack{x_1 \geq 0 \\ x_2 \geq 0}} \{ x_1^2 - ux_1 + \min_{\substack{x_1 \geq 0 \\ x_2 \geq 0}} \{ x_2^2 - ux_2 \}, u \geq 0 \}.$$

$$= 4u + \min_{\mathbb{R}^2} \{ x_1^2 + x_2^2 - ux_1 - ux_2 \}$$

$$(4 - ux_1 - ux_2)u - (x_1^2 + x_2^2) = b(u) = \min_{\mathbb{R}^2} L(\mathbf{x}, u) = f(\mathbf{x}) - g(\mathbf{x})$$

- The Lagrangian dual function is

- $\mathbf{x}^* = \mathbf{x}(u)$ is also unique, and is automatically given by $\mathbf{x}^* = \mathbf{x}_*$ differentiable. In this particular case, the optimum \mathbf{x}_*
- This is an example where the dual function is
- Also: $f(u) = (\mathbf{x}_*)^\top b = 8$.
- $\mathbf{x}_* = (x_1(u_*), x_2(u_*))^\top = (2, 2)$.
- If $u = 4 < 0$, it is the optimum in the dual problem!
- We then have that $d'(u) = 4 - u = 0 \iff u = 4$. As

- Consider Lagrangian relaxing the first constraint,
 - The optimal solution is $\mathbf{x}_* = (3/2, 0)^T$.
- $$0 \leq x_2 \leq 1.$$

$$0 \leq x_1 \leq 2,$$

 subject to $2x_1 + 4x_2 \leq 3,$

$$\underset{\mathbf{x}}{\text{minimize}} \quad f(\mathbf{x}) = -x_1 - x_2,$$
- Consider the linear programming problem to
- problem**
- Examples, II:** An implicit, non-differentiable dual

given in (11).

non-differentiable at u^* . We utilize the characterization

obtain the optimal primal solution from u^* ? d is

linear programs, we have strong duality, but how do we

- We have that $u^* = 1/2$, and hence $d(u^*) = -3/2$. For

$$\left. \begin{aligned} & -3u, \quad 1/2 \leq u. \\ & -2 + u, \quad 1/4 \leq u \leq 1/2, \\ & -3 + 5u, \quad 0 \leq u \leq 1/4, \end{aligned} \right\} =$$

$$L(\mathbf{x}, u) = -x_1 - x_2 + u(2x_1 + 4x_2 - 3);$$

$$d(u) = -3u + \min_{\substack{0 \leq x_1 \leq 2 \\ 0 \leq x_2 \leq 1}} \{ (-1 + 2u)x_1 + (-1 + 4u)x_2 \}$$

obtaining

- First, at u_* , it is clear that $X(u_*)$ is the set $\{ \begin{pmatrix} 0 \\ 2a \end{pmatrix} \mid 0 \leq a \leq 1 \} \cdot$. Among the subproblem solutions, we next have to find one that is primal feasible as well as complementary.
- Primal feasibility means that $2 \cdot 2a + 2 \cdot 0 \leq 3 \iff a \leq 3/4 \cdot$
- Further, complementarity means that $u_* \cdot (2x_1^* + 4x_2^* - 3) = 0 \iff a = 3/4$, since $u_* \neq 0$. We conclude that the only primal vector x that satisfies the system (11) together with the dual optimal solution $u_* = 1/2$ is $x^* = (3/2, 0)^T \cdot$
- Observe finally that $f_* = b_* \cdot$

- Why must $u^* = 1/2$? According to the global optimality conditions, the optimal solution must in this convex case be among the subproblem solutions. Since x_1^* is not in one of the „corners“ (it is between 0 and 2), the value of u^* has to be such that the cost term for x_1 in $L(x, u^*)$ is identically zero! That is, $-1 + u^* \cdot 2 = 0$ implies that $u^* = 1/2$!

- A non-coordinability phenomenon (non-unique subproblem solution means that the optimal solution is not obtained automatically).
- In non-convex cases, the optimal solution may not be among the points in $X(\boldsymbol{\mu}_*)$. What do we do then??

- Let $f : \mathbb{R}^n \rightarrow \mathbb{R}$ be a convex function. We say that a vector $\mathbf{d} \in \mathbb{R}^n$ is a subgradient of f at $\mathbf{x} \in \mathbb{R}^n$ if a such vectors \mathbf{p} defines the subdifferential of f at \mathbf{x} , and is denoted $\partial f(\mathbf{x})$.
 - This set is the collection of "slopes" of the function f at \mathbf{x} .
 - For every $\mathbf{x} \in \mathbb{R}^n$, $\partial f(\mathbf{x})$ is a non-empty, convex, and compact set.
- (14)
- $$(\mathbf{x} - \mathbf{y})^\top \mathbf{d} + f(\mathbf{x}) \leq f(\mathbf{y}).$$

Subgradients of convex functions

Figure 1: Four possible slopes of the convex function f at x .

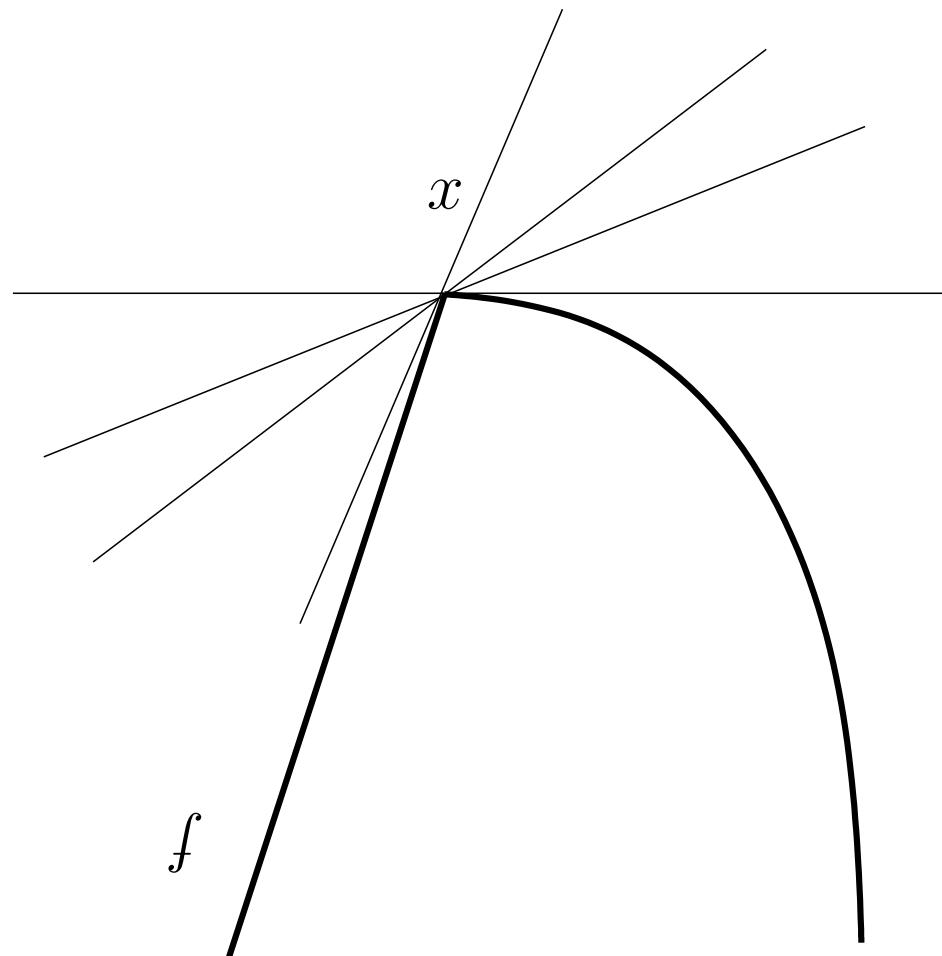
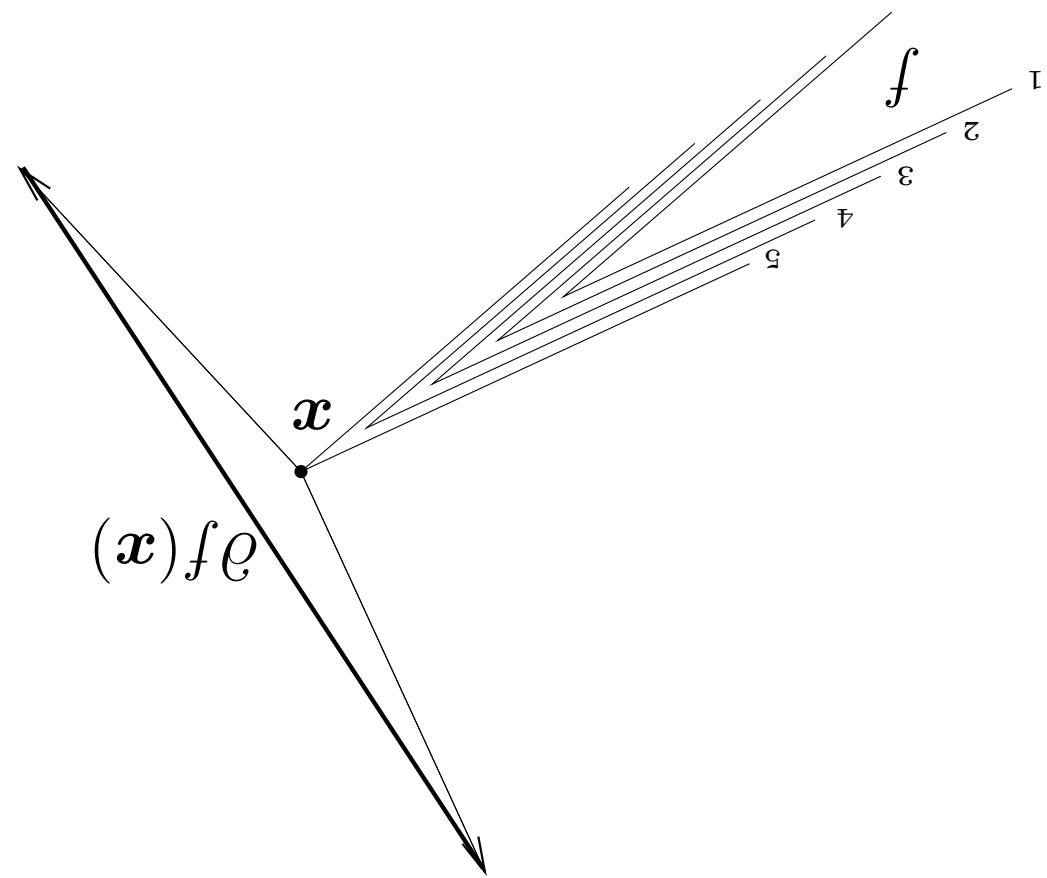


Figure 2: The subdifferential of a convex function f at x .



- The convex function f is differentiable at x exactly when there exists one and only one subgradient of f at x , which then is the gradient of f at x , $\Delta f(x)$.

Differentiability of the Lagrangian dual function:

• Consider the problem (4), under the assumption that $f, g_i (i = 1, \dots, m)$ continuous; X nonempty and compact.

• Then, the set of solutions to the Lagrangian

$$(15) \quad (L(x, \boldsymbol{\mu})) = \arg \min_{\boldsymbol{x} \in X} L(\boldsymbol{x}, \boldsymbol{\mu}), \quad \boldsymbol{\mu} \in \mathbb{R}^m,$$

subproblem,

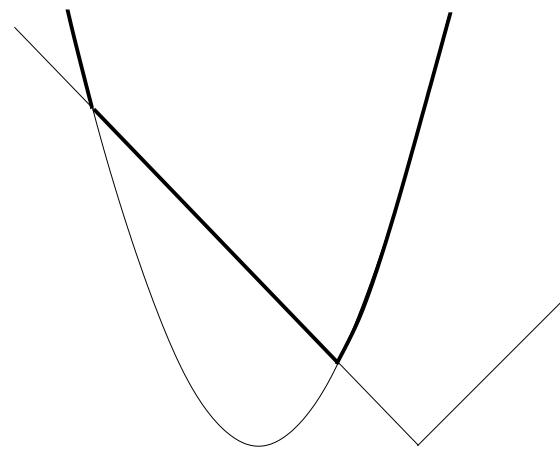
- We develop the sub-differentiability properties of the function ϕ .
is non-empty and compact for every $\boldsymbol{\mu}$.

function ϕ .

$$\begin{aligned}
& \cdot(x)\mathbf{g}_{\perp}(\mathbf{n} - \underline{\mathbf{n}}) + (\mathbf{n})b = \\
& (x)\mathbf{g}_{\perp}\mathbf{n} + (x)\mathbf{g}_{\perp}(\mathbf{n} - \underline{\mathbf{n}}) + (x)f = \\
& (x)\mathbf{g}_{\perp}\underline{\mathbf{n}} + (x)f \stackrel{X \in \mathcal{Y}}{\geq} \inf_{\mathbf{y}} L(\mathbf{y}, \underline{\mathbf{n}}) = (\underline{\mathbf{n}})b
\end{aligned}$$

- Proof. Let $\underline{\mathbf{n}} \in \mathbb{R}^m$ be arbitrary. We have that
 - to b at \mathbf{n} , that is, $\mathbf{g}(x) \in \partial b(\mathbf{n})$
- Let $\mathbf{n} \in \mathbb{R}^m$. If $x \in X(\mathbf{n})$, then $\mathbf{g}(x)$ is a subgradient of its supremum over \mathbb{R}_+^m . If its supremum over \mathbb{R}_+^m is attained, then the optimal solution set therefore is closed and convex.
- The dual function b is finite, continuous and concave on \mathbb{R}^m .
- Suppose that, in the problem (4), (15) holds.

Subgradients and gradients of b



- Let $h(x) = \min\{h_1(x), h_2(x)\}$, where $h_1(x) = |x - 4|$ and $h_2(x) = 4 - (x - 2)^2$.
- Then,
$$h(x) = \begin{cases} 4 - (x - 2)^2, & x \geq 1, x \leq 4 \\ |x - 4|, & 1 \leq x \leq 4 \end{cases}$$

Example

- The function h is non-differentiable at $x = 1$ and $x = 4$, since its graph has non-unique supporting hyperplanes there.
 - The subdifferential is either a singleton (at non-differentiable points) or an interval (at differentiable points) or an interval (at non-differentiable points).
- $$\partial h(x) = \begin{cases} \{-1\}, & 1 > x > 4 \\ \{4 - 2x\}, & x < 1, x > 4 \\ [-1, 2], & x = 1 \\ [-4, -1], & x = 4 \end{cases}$$

- Let $\mathbf{u} \in \mathbb{R}^m$. Then, $\partial q(\mathbf{u}) = \text{conv} \{ (\mathbf{x})\mathbf{b} \mid (\mathbf{x})\mathbf{b} \in X(\mathbf{u}) \}$
- Let $\mathbf{u} \in \mathbb{R}^m$. The dual function q is differentiable at \mathbf{u} if and only if $\{ \mathbf{g}(\mathbf{x}) \mid \mathbf{x} \in X(\mathbf{u}) \}$ is a singleton set [the vector of constraint functions is invariant over $X(\mathbf{u})$].
- Then, for every $\mathbf{x} \in X(\mathbf{u})$.
 - Holds in particular if the Lagrangian subproblem has a unique solution [$X(\mathbf{u})$ is a singleton set]. In particular, if \mathbf{f} is strictly convex on X , and \mathbf{g}_i ($i = 1, \dots, m$) are convex, q then even in C_1 . \square

The Lagrangian dual problem

$$\cdot \left\{ (\underline{x})_{\mathcal{I}} \ni i \mid 0 \leq h_i(\underline{x}) \leq \sum_{j \in \mathcal{I} \setminus \{i\}} h_j(\underline{x}) \right\} = (\underline{x})_{\mathcal{I}}$$

$\{h_i(\underline{x}) \mid i \in \mathcal{I}\}$, that is,

- Then, the subdifferential $Q_h(\underline{x})$ is the convex hull of

segments at \underline{x} .

- Let $\mathcal{I}(\underline{x}) \subseteq \{1, \dots, m\}$ be defined by $i \in \mathcal{I}(\underline{x})$ for $h_i(\underline{x}) > h(\underline{x})$ (the active

is a concave function on \mathbb{R}^n .

function h_i is concave and differentiable on \mathbb{R}^n , then h_i

- Theorem: If $h(\underline{x}) = \min_{i=1, \dots, m} h_i(\underline{x})$, where each

- How do we write the subdifferential of h ?

□

that is, $\mathbf{0}_n \in \partial h(\mathbf{x}^*)$,
 $h(\mathbf{x}) \leq h(\mathbf{x}^*)$ for all $\mathbf{x} \in \mathbb{R}^n$,
Suppose that $\mathbf{x}^* \in \arg \max_{\mathbf{x} \in \mathbb{R}^n} h(\mathbf{x})$

$h(\mathbf{x}^*) \geq h(\mathbf{x})$ for all $\mathbf{x} \in \mathbb{R}^n$,
 $h(\mathbf{x}^*) \geq h(\mathbf{x})$ for all $\mathbf{x} \in \mathbb{R}^n$, that is,
 $\mathbf{x}^* \in \partial h(\mathbf{x}^*)$

$$\cdot (\mathbf{x}^*) \in \partial h(\mathbf{x}^*) \iff \mathbf{x}^* \in \arg \max_{\mathbf{x} \in \mathbb{R}^n} h(\mathbf{x})$$

• Theorem: Assume that h is concave on \mathbb{R}^n . Then,

$$\cdot \mathbf{0}_n = (\mathbf{x}^*) \Delta \iff \mathbf{x}^* \in \arg \max_{\mathbf{x} \in \mathbb{R}^n} h(\mathbf{x})$$

• For a differentiable, concave function h it holds that

Optimality conditions for the dual problem

the conical hull of the active constraints' normals at \mathbf{x}_* .
 where $N^X(\mathbf{x}_*)$ is the normal cone to X at \mathbf{x}_* , that is,

$$\emptyset \neq (\mathbf{x}_*)^X N \cup (\mathbf{x}_*)^h \partial h(\mathbf{x}_*) \iff \mathbf{x}_* \in \arg \max_{\mathbf{x} \in X} h(\mathbf{x})$$

are generalized thus:

- For optimization with constraints the KKT conditions
- The example: $0 \in \partial h(1) \iff x_* = 1$.

$$0 \leq {}_*x \quad ; 0 = ({}_*x)_h \cdot {}_*x \quad ; 0 \geq ({}_*x)_h$$

concave): x_* is optimal if and only if

- Compare with a one-dimensional max-problem (y

$$\text{subgradient } g \in Q(\mathbf{u}_*) \text{ for which the following holds:} \\ \mathbf{u}_*^i g_i = 0, i = 1, \dots, m.$$

- Consider the dual problem (9), and let $\mathbf{u}_* \leq \mathbf{0}_m$. It is then optimal in (9) if and only if there exists a subgradient $g \in Q(\mathbf{u}_*)$ for which the following holds:

- In the case of the dual problem we have only sign conditions.

where $\mathbf{g}_k \in Q(\mathbf{n}_k)$ is arbitrarily chosen.

$$(17c) \quad = (\max_{i=1}^m \{0, \mathbf{n}_k(i) + \alpha_k \mathbf{g}_k(i)\})$$

$$(17b) \quad = [\mathbf{n}_k + \alpha_k \mathbf{g}_k]^+$$

$$(17a) \quad \mathbf{n}_{k+1} = \text{Proj}_{\mathbb{R}^m_+} [\mathbf{n}_k + \alpha_k \mathbf{g}_k]$$

- The simplest type of iteration has the form

using a single subgradient in each iteration.

functions, generating a sequence of dual vectors in \mathbb{R}^m_+

methods from the C_1 to general convex (or, concave)
- Subgradient methods extend gradient projection

A subgradient method for the dual problem

$$(18) \quad \alpha^k \in (0, 2[b_* - \|g^k\|_2] / \|g^k\|_2).$$

holds for every step length α^k in the interval

$$\|\mathbf{n}_* - g^k\| > \|\mathbf{n}_* - g^{k+1}\|$$

every optimal solution $\mathbf{n}_* \in U_*$ in (9),

- Suppose that $\mathbf{n} \in \mathbb{R}_+^m$ is not optimal in (9). Then, for

lengths α^k .

- Cannot do line searches; must use predetermined step

may not be an ascent direction

- Main difference to C_1 case: an arbitrary subgradient g^k

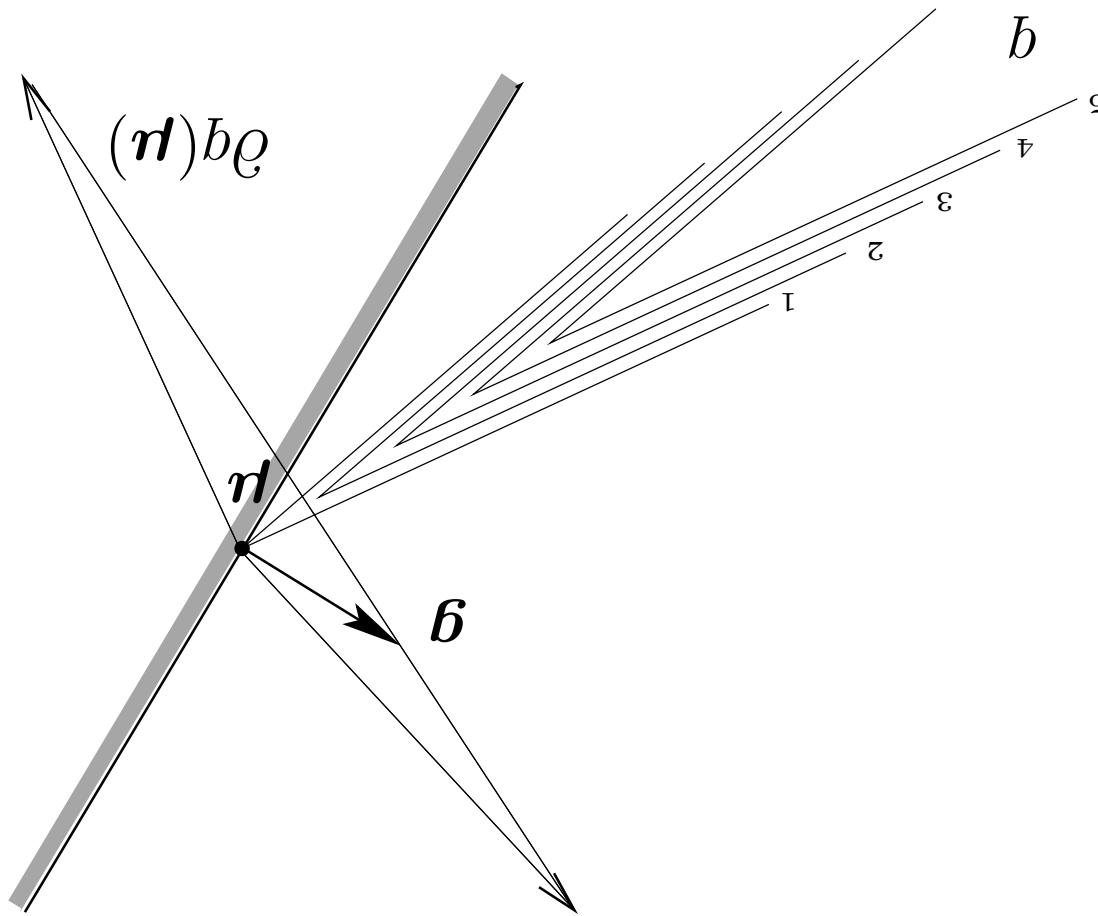
$x^k \in \arg \min_{x \in X} L(x, \mathbf{n}^k)$.

- We often use $g^k = g(x^k)$, where

- Why? Let $\mathbf{g} \in Q(\underline{n})$, and let U_* be the set of optimal solutions to (9). Then,
$$\cdot \{ 0 \leq (\mathbf{n} - \underline{n})_L \mathbf{g} \mid \mathbf{n} \in U_* \} \subseteq U$$
In other words, \mathbf{g} defines a half-space that contains the set of optimal solutions.
- Good news: If the step length is small enough we get closer to the set of optimal solutions!

Figure 3: The half-space defined by the subgradient \mathbf{g} of y at \mathbf{u} . Note

that the subgradient is not an ascent direction.



$$\alpha_k < 0, k = 1, 2, \dots; \quad \lim_{k \rightarrow \infty} \alpha_k = 0; \quad \sum_{\infty}^{l=s} \alpha_k = +\infty. \quad (20)$$

- The divergent series step length rule:

$$y_k = y_* - \frac{\|g_k\|^2}{\alpha_k} \quad (\text{updated.})$$

(Can be replaced by an upper bound on y_* which is updated.)
- Bad news: Utilizes knowledge of the optimal value y_* !
 to converge to zero or a too large value.
- $\alpha < 0$ makes sure that we do not allow the step lengths
- $\alpha \leq \alpha_k \leq 2[y_* - g_k]/\|g_k\|^2 - \alpha, \quad k = 1, 2, \dots. \quad (19)$
- Polyaak step length rule:

(21)

$$\sum_{s=1}^{\infty} a_s^2 < +\infty.$$

- Additional condition often added:

- Suppose that the problem (4) is feasible, and that (15) and (13) hold.
 - (a) Let $\{\mathbf{u}_k\}$ be generated by the method (17), under the Polyak step length rule (19), where α is a small positive number. Then, $\{\mathbf{u}_k\}$ converges to an optimal solution to (9).
 - (b) Let $\{\mathbf{u}_k\}$ be generated by the method (17), under the divergent step length rule (20). Then, $\{\mathbf{u}_k\}$ converges to an optimal solution to (9).
 - (c) Let $\{\mathbf{u}_k\}$ be generated by the method (17), under the divergent step length rule (21). Then, $\{\mathbf{u}_k\}$ converges to an optimal solution to (9).
-

- Given $\mathbf{u}_k \geq \mathbf{0}_m$.
- Solve the Lagrangian subproblem to minimize $L(\mathbf{x}, \mathbf{u}_k)$ over $\mathbf{x} \in X$.
- Let an optimal solution to this problem be \mathbf{x}_k .
- Calculate $\mathbf{g}(\mathbf{x}_k) \in Q(\mathbf{u}_k)$.
- Take a (positive) step in the direction of $\mathbf{g}(\mathbf{x}_k)$ from \mathbf{u}_k , according to a step length rule.
- Set any negative components of this vector to zero.
- We have obtained \mathbf{u}_{k+1} .

Application to the Lagrangian dual problem

- We can choose the subgradient more carefully, such that we will obtain ascent directions. This amounts to gathering several subgradients at nearby points and solving quadratic programming problems to find the best convex combination of them. (*Bundle methods*)
- Pre-multiply the subgradient obtained by some positive definite matrix. We get methods similar to Newton methods. (*Space dilation methods*)
- Pre-project the subgradient vector (onto the tangent cone of \mathbb{R}_m^+) so that the direction taken is a feasible direction. (*Subgradient-projection methods*.)

Additional algorithms

- Discrete Optimization: The size of the duality gap, and the relation to the continuous relaxation.
- Convexification.
- Primal feasibility heuristics.
- Global optimality conditions for discrete optimization (and general problems).

More to come